

Notes on homotopical algebra

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1st May 2013

PREFACE

These notes are intended as a kind of annotated index to the various standard references in homotopical algebra: the focus is on definitions and statements of results, *not* proofs.

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FOUNDATIONS

0.1 Set theory

In category theory it is often convenient to invoke a certain set-theoretic device commonly known as a ‘Grothendieck universe’, but we shall say simply ‘universe’, so as to simplify exposition and proofs by eliminating various circumlocutions involving cardinal bounds, proper classes etc.

Definition 0.1.1. A **pre-universe** is a set \mathbf{U} satisfying these axioms:

1. If $x \in y$ and $y \in \mathbf{U}$, then $x \in \mathbf{U}$.
2. If $x \in \mathbf{U}$ and $y \in \mathbf{U}$ (but not necessarily distinct), then $\{x, y\} \in \mathbf{U}$.
3. If $x \in \mathbf{U}$, then $\mathcal{P}(x) \in \mathbf{U}$, where $\mathcal{P}(x)$ denotes the set of all subsets of x .
4. If $x \in \mathbf{U}$ and $f : x \rightarrow \mathbf{U}$ is a map, then $\bigcup_{i \in x} f(i) \in \mathbf{U}$.

A **universe** is a pre-universe \mathbf{U} with this additional property:

5. $\omega \in \mathbf{U}$, where ω is the set of all finite (von Neumann) ordinals.

Example 0.1.2. The empty set is a pre-universe, and with very mild assumptions, so is the set \mathbf{HF} of all hereditarily finite sets.

¶ 0.1.3. The notion of universe makes sense in any material set theory, but their existence must be postulated. We adopt the following:

- **Grothendieck–Verdier universe axiom.** For each set x , there exists a universe \mathbf{U} with $x \in \mathbf{U}$.

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For definiteness, we may take our base theory to be Mac Lane set theory, which is a weak subsystem of Zermelo–Fraenkel set theory with choice (ZFC). Readers interested in the details of Mac Lane set theory are referred to [Mathias, 2001], but in practice, as long as one is working at all times *inside some universe*, one may as well be working in ZFC. Indeed:

Proposition 0.1.4. *With the assumptions of Mac Lane set theory, any universe is a transitive model of ZFC.*

Proof. Let \mathbf{U} be a universe. By definition, \mathbf{U} is a transitive set containing pairs, power sets, unions, and ω , so the axioms of extensionality, empty set, pairs, power sets, unions, choice, and infinity are all automatically satisfied. We must show that the axiom schemas of separation and replacement are also satisfied, and in fact it is enough to check that replacement is valid; but this is straightforward using axioms 2 and 4. ■

Definition 0.1.5. Let \mathbf{U} be a pre-universe. A **U-set** is a member of \mathbf{U} , a **U-class** is a subset of \mathbf{U} , and a **proper U-class** is a U-class that is not a U-set.

Lemma 0.1.6. *A U-class X is a U-set if and only if there exists a U-class Y such that $X \in Y$.* ■

Proposition 0.1.7. *If \mathbf{U} is a universe in Mac Lane set theory, then the collection of all U-classes is a transitive model of Morse–Kelley class–set theory (MK), and so is a transitive model of von Neumann–Bernays–Gödel class–set theory (NBG) in particular.* ■

Definition 0.1.8. A **U-small category** is a category \mathcal{C} such that $\text{ob } \mathcal{C}$ and $\text{mor } \mathcal{C}$ are U-sets. A **locally U-small category** is a category \mathcal{D} satisfying these conditions:

- $\text{ob } \mathcal{D}$ and $\text{mor } \mathcal{D}$ are U-classes, and
- for all objects x and y in \mathcal{D} , the hom-set $\mathcal{D}(x, y)$ is a U-set.

An **essentially U-small category** is a category \mathcal{D} for which there exist a U-small category \mathcal{C} and a functor $\mathcal{C} \rightarrow \mathcal{D}$ that is fully faithful and essentially surjective on objects.

Proposition 0.1.9. *If \mathcal{D} is a U-small category and \mathcal{C} is a locally U-small category, then the functor category $[\mathcal{D}, \mathcal{C}]$ is locally U-small.*

Proof. Strictly speaking, this depends on the set-theoretic implementation of ordered pairs, categories, functors, etc., but at the very least $[\mathbb{D}, C]$ should be isomorphic to a locally \mathbf{U} -small category.

In the context of $[\mathbb{D}, C]$, we may regard functors $\mathbb{D} \rightarrow C$ as being the pair consisting of the *graph* of the object map $\text{ob } \mathbb{D} \rightarrow \text{ob } C$ and the *graph* of the morphism map $\text{mor } \mathbb{D} \rightarrow \text{mor } C$, and these are \mathbf{U} -sets by the \mathbf{U} -replacement axiom. Similarly, if F and G are objects in $[\mathbb{D}, C]$, then we may regard a natural transformation $\alpha : F \Rightarrow G$ as being the triple (F, G, A) , where A is the set of all pairs (c, α_c) . ■

One complication introduced by having multiple universes concerns the existence of (co)limits.

Theorem 0.I.10 (Freyd). *Let C be a category and let κ be a cardinal such that $|\text{mor } C| \leq \kappa$. If C has products for families of size κ , then any two parallel morphisms in C must be equal.*

Proof. Suppose, for a contradiction, that $f, g : X \rightarrow Y$ are distinct morphisms in C . Let Z be the product of κ -many copies of Y in C . The universal property of products implies there are at least 2^κ -many distinct morphisms $X \rightarrow Z$; but $C(X, Z) \subseteq \text{mor } C$, so this is an absurdity. ■

Definition 0.I.11. Let \mathbf{U} be a pre-universe. A **\mathbf{U} -complete** (resp. **\mathbf{U} -cocomplete**) **category** is a category C with the following property:

- For all \mathbf{U} -small categories \mathbb{D} and all diagrams $A : \mathbb{D} \rightarrow C$, a limit (resp. colimit) of A exists in C .

We may instead say C has all **finite limits** (resp. **finite colimits**) in the special case $\mathbf{U} = \mathbf{HF}$.

Proposition 0.I.12. *Let C be a category and let \mathbf{U} be a non-empty pre-universe. The following are equivalent:*

- (i) C is \mathbf{U} -complete.
- (ii) C has all finite limits and products for all families of objects indexed by a \mathbf{U} -set.

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(iii) For each \mathbf{U} -small category \mathbb{D} , there exists an adjunction

$$\Delta \dashv \lim_{\leftarrow \mathbb{D}} : [\mathbb{D}, \mathcal{C}] \rightarrow \mathcal{C}$$

where ΔX is the constant functor with value X .

Dually, the following are equivalent:

(i') \mathcal{C} is \mathbf{U} -cocomplete.

(ii') \mathcal{C} has all finite colimits and coproducts for all families of objects indexed by a \mathbf{U} -set.

(iii') For each \mathbf{U} -small category \mathbb{D} , there exists an adjunction

$$\lim_{\rightarrow \mathbb{D}} \dashv \Delta : \mathcal{C} \rightarrow [\mathbb{D}, \mathcal{C}]$$

where ΔX is the constant functor with value X .

Proof. This is a standard result; but we remark that we do require a sufficiently powerful form of the axiom of choice to pass from (ii) to (iii). \square

¶ 0.1.13. In the **explicit universe convention**, the words ‘set’, ‘class’, etc. have their usual meanings, and in the **one-universe convention**, these instead abbreviate ‘ \mathbf{U} -set’, ‘ \mathbf{U} -class’, etc. for a fixed (but arbitrary) universe \mathbf{U} . However, the word ‘category’ always refers to a category that is contained in *some* universe, which may or may not be locally \mathbf{U} -small, and we shall use the word ‘ensemble’ to refer to sets which may or may not be in \mathbf{U} . In subsequent chapters, the implicit universe convention should be assumed *unless otherwise stated*.

We now recall some definitions and results about ordinal and cardinal numbers. Readers familiar with axiomatic set theory may wish to skip ahead.

Definition 0.1.14. A **von Neumann ordinal** is a set α with the following properties:

- If $x \in y$ and $y \in \alpha$, then $x \in \alpha$.
- The binary relation \in is strict total ordering of α .
- If S is a subset of α such that

$$- \emptyset \in S,$$

- If $\beta \in S$ and $\beta \cup \{\beta\} \in \alpha$, then $\beta \cup \{\beta\} \in S$.
- If $T \subseteq S$, then $\bigcup T \in S$.

then $S = \alpha$.

We identify 0 with the von Neumann ordinal \emptyset , and by induction, we identify the natural number $n + 1$ with the von Neumann ordinal $\{0, \dots, n\}$.

Proposition 0.I.15.

- (i) *If α is a von Neumann ordinal, then every member of α is an initial segment of α and is in particular a von Neumann ordinal.*
- (ii) *If α is a von Neumann ordinal, so is $\alpha \cup \{\alpha\}$. (This is usually denoted by $\alpha + 1$ and called the **successor** of α .)*
- (iii) *The union of a set S of von Neumann ordinals is another von Neumann ordinal. (This is usually denoted by $\sup S$ and called the **supremum** of S .)*
- (iv) *If \mathbf{U} is a pre-universe and $\kappa(\mathbf{U})$ is the set of von Neumann ordinals in \mathbf{U} , then $\kappa(\mathbf{U})$ a von Neumann ordinal, but $\kappa(\mathbf{U}) \notin \mathbf{U}$.*

Proof. Claims (i) – (iii) are all easy, and claim (iv) is Burali-Forti’s paradox. ◆

Theorem 0.I.16 (Classification of well-orderings).

- (i) *In Zermelo–Fraenkel set theory, every well-ordered set is isomorphic to a unique von Neumann ordinal.*
- (ii) *In Mac Lane set theory, if \mathbf{U} is a pre-universe and X is a well-ordered set in \mathbf{U} , then X is isomorphic to a unique von Neumann ordinal in \mathbf{U} .*

Proof. Claim (i) is a standard result in axiomatic set theory, and claim (ii) is an obvious corollary. □

Definition 0.I.17. A **transitive set** is a set T such that, given $x \in y$, if $y \in T$, then $x \in T$ as well. The **transitive closure** of a set X is a set $\text{tcl}(X)$ such that, for all transitive sets T with $X \subseteq T$, we have $\text{tcl}(X) \subseteq T$ as well.

Lemma 0.I.18. *In Mac Lane set theory, every set has a unique transitive closure.*

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Proof. One of the axioms of Mac Lane set theory states that every set X is a member of some transitive set T , and so $X \subseteq T$. Clearly, the intersection of any family of transitive sets containing X is again a transitive set containing X , so $\text{tcl}(X)$ exists and is unique so long as there is at least one transitive set containing X . ■

Definition 0.1.19. A **partial rank function** from a transitive set T to a well-ordered set W is a partial function $\rho : T \rightarrow W$ with these properties:

- If $\emptyset \in T$, then $\rho(\emptyset)$ is the least element of W .
- If $y \in T$ and $\rho(x)$ is defined for all $x \in y$, then

$$\rho(y) = \min \{w \in W \mid \forall x \in y. \rho(x) < w\}$$

provided the RHS is defined.

- Otherwise $\rho(y)$ is undefined.

A **total rank function** is a partial rank function that is defined on its entire domain. The **rank** of a set X , if it exists, the least von Neumann ordinal $\text{rank}(X)$ for which there exists a total rank function $\text{tcl}(X) \rightarrow \text{rank}(X)$.

Proposition 0.1.20. *In Mac Lane set theory:*

- (i) *If T is a transitive set and W is a well-ordered set, then there is a unique partial rank function $\rho : T \rightarrow W$.*
- (ii) *If \mathbf{U} is a pre-universe and $x \in \mathbf{U}$, then $\text{rank}(x)$ can be defined by a Δ_0 -formula with \mathbf{U} as a parameter, and for each von Neumann ordinal α in \mathbf{U} , the set*

$$\mathbf{V}_\alpha = \{x \in \mathbf{U} \mid \text{rank}(x) < \alpha\}$$

is a \mathbf{U} -set.

- (iii) *Assuming the Grothendieck–Verdier universe axiom, $\text{rank}(x)$ is defined for all x .*

Proof. (i). This is a straightforward application of well-founded induction.

(ii). \mathbf{U} is a transitive set and the set $\kappa(\mathbf{U})$ of all von Neumann ordinals in \mathbf{U} is well-ordered by inclusion, so by claim (i) there is a partial rank function $\rho :$

$\mathbf{U} \rightarrow \kappa(\mathbf{U})$. ZFC proves that every set has a rank, so ρ must in fact be a total rank function; hence, for any $x \in \mathbf{U}$, $\text{rank}(x)$ is defined. It is clear that ρ can be defined by a Δ_0 -formula with only \mathbf{U} as a parameter, and the rest of the claim follows.

(iii). Obvious, assuming claim (ii). □

Definition 0.I.21. Two sets are **equinumerous** if there exists a bijection between them. A **cardinality class** in a pre-universe \mathbf{U} is an equivalence class under the relation of equinumerosity.

Definition 0.I.22. An **\aleph -number** is an infinite von Neumann ordinal κ such that, for any von Neumann ordinal λ such that κ and λ are equinumerous, we have $\kappa \subseteq \lambda$.

Example 0.I.23. The first infinite von Neumann ordinal, i.e. $\omega = \{0, 1, 2, \dots\}$, is the \aleph -number \aleph_0 .

Theorem 0.I.24 (Classification of cardinalities).

- (i) *In Zermelo–Fraenkel set theory, for every well-ordered infinite set X , there exists a unique \aleph -number κ such that X and κ are equinumerous.*
- (ii) *In Zermelo–Fraenkel set theory with the axiom of choice, the same is true for any infinite set whatsoever.*
- (iii) *In Mac Lane set theory, if \mathbf{U} is a universe and X is an infinite set in \mathbf{U} , then there exists a unique \aleph -number κ in the cardinality class of X .*
- (iv) *In Mac Lane set theory with the Grothendieck–Verdier universe axiom, if \mathbf{U} is a pre-universe and κ is an \aleph -number not in \mathbf{U} , then the cardinality of \mathbf{U} is at most κ .*

Proof. Claim (i) is a standard fact, whence claims (ii) and (iii), by the well-ordering theorem. Claim (iv) can be proven using axiom 4 for pre-universes. □

¶ 0.I.25. Henceforth, we identify the cardinality class of a finite set with the unique von Neumann ordinal contained in that class, and similarly we identify the cardinality class of an infinite set with the unique \aleph -number in that class. These are the **cardinal numbers**.

Definition 0.1.26. A **cofinal subset** of a partially-ordered set X is a subset $Y \subseteq X$ such that, for all x in X , there exists some y in Y such that $x \leq y$. A **regular cardinal number** is an \aleph -number κ such that any cofinal subset of κ has cardinality equal to κ . A **singular cardinal number** is an \aleph -number that is not regular.

The following helps to motivate the definition of regular cardinal numbers.

Definition 0.1.27. Let \mathbf{U} be a pre-universe. An **arity class** in \mathbf{U} is a \mathbf{U} -class K of cardinal numbers satisfying the following conditions:

- $1 \in K$.
- If $\kappa \in K$ and $\lambda : \kappa \rightarrow K$ is a function, then the cardinal sum $\sum_{\alpha \in \kappa} \lambda(\alpha)$ is also in K .
- If $\kappa \in K$ and $\lambda : \kappa \rightarrow \mathbf{U}$ is a function such that each $\lambda(\alpha)$ is a cardinal number and $\sum_{\alpha \in \kappa} \lambda(\alpha) \in K$, then $\lambda(\alpha) \in K$ as well.

Theorem 0.1.28 (Classification of arity classes). *In Mac Lane set theory, if K is an arity class in a pre-universe \mathbf{U} , then K must be either*

- $\{1\}$, or
- $\{0, 1\}$, or
- of the form $\{\lambda \in \mathbf{U} \mid \lambda \text{ is a cardinal number and } \lambda < \kappa\}$ for some regular cardinal number κ (possibly not in \mathbf{U}).

Proof. The notion of arity class and this result are due to Shulman [2012]. \square

Definition 0.1.29. Let κ be a regular cardinal number. A **κ -small category** is a category \mathbb{C} such that $\text{mor } \mathbb{C}$ has cardinality $< \kappa$. A **finite category** is an \aleph_0 -small category, i.e. a category \mathbb{C} such that $\text{mor } \mathbb{C}$ is finite. A **finite diagram** (resp. **κ -small diagram**, **\mathbf{U} -small diagram**) in a category \mathcal{C} is a functor $\mathbb{D} \rightarrow \mathcal{C}$ where \mathbb{D} is a finite (resp. κ -small, \mathbf{U} -small) category.

Theorem 0.1.30. *Let \mathbf{U} be a pre-universe, let \mathbf{U}^+ be a universe with $\mathbf{U} \in \mathbf{U}^+$, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathbf{Set}^+ be the category of \mathbf{U}^+ -sets.*

- (i) *If $X : \mathbb{D} \rightarrow \mathbf{Set}$ is a \mathbf{U} -small diagram, then there exist a limit and a colimit for X in \mathbf{Set} .*

- (ii) *The inclusion $\mathbf{Set} \hookrightarrow \mathbf{Set}^+$ is fully faithful and preserves limits and colimits for all \mathbf{U} -small diagrams.*

Proof. One can construct products, equalisers, coproducts, coequalisers, and hom-sets in a completely explicit way, making the preservation properties obvious. ◆

Corollary 0.I.31. *The inclusion $\mathbf{Set} \hookrightarrow \mathbf{Set}^+$ reflects limits and colimits for all \mathbf{U} -small diagrams.* ■

Corollary 0.I.32. *For any \mathbf{U} -small category \mathbb{C} :*

- (i) *The functor category $[\mathbb{C}, \mathbf{Set}]$ is \mathbf{U} -complete and \mathbf{U} -cocomplete, with limits and colimits for \mathbf{U} -small diagrams computed componentwise in \mathbf{Set} .*
- (ii) *The inclusion $[\mathbb{C}, \mathbf{Set}] \hookrightarrow [\mathbb{C}, \mathbf{Set}^+]$ is fully faithful and both preserves and reflects limits and colimits for all \mathbf{U} -small diagrams.* ■

Definition 0.I.33. An **strongly inaccessible cardinal number** is a regular cardinal number κ such that, for all sets X of cardinality less than κ , the power set $\mathcal{P}(X)$ is also of cardinality less than κ .

Example 0.I.34. \aleph_0 is a strongly inaccessible cardinal number and is the only one that can be proven to exist in ZFC. It is more conventional to exclude \aleph_0 from the definition of strongly inaccessible cardinal number by demanding that they be uncountable.

Proposition 0.I.35. *In Mac Lane set theory:*

- (i) *If \mathbf{U} is a non-empty pre-universe, then there exists a strongly inaccessible cardinal number κ such that the members of \mathbf{U} are all the sets of rank less than κ . Moreover, this κ is the rank and the cardinality of \mathbf{U} .*
- (ii) *If \mathbf{U} is a universe and κ is a strongly inaccessible cardinal number such that $\kappa \in \mathbf{U}$, then there exists a \mathbf{U} -set \mathbf{V}_κ whose members are all the sets of rank less than κ , and \mathbf{V}_κ is a pre-universe.*
- (iii) *If \mathbf{U} and \mathbf{U}' are pre-universes, then either $\mathbf{U} \subseteq \mathbf{U}'$ or $\mathbf{U}' \subseteq \mathbf{U}$; and if $\mathbf{U} \not\subseteq \mathbf{U}'$, then $\mathbf{U} \in \mathbf{U}'$.*

Proof. (i). Let κ be the set of all von Neumann ordinals in \mathbf{U} ; this exists by Δ_0 -separation applied to \mathbf{U} . Since \mathbf{U} is closed under power sets and internally-indexed unions, κ must be a strongly inaccessible cardinal.

We can construct the set all of \mathbf{U} -sets of rank less than κ using transfinite recursion on κ as follows: starting with $\mathbf{V}_0 = \emptyset$, for each von Neumann ordinal α less than κ , we set $\mathbf{V}_{\alpha+1} = \mathcal{P}(\mathbf{V}_\alpha)$, and for each ordinal λ that is not a successor, we set $\mathbf{V}_\lambda = \bigcup_{\alpha < \lambda} \mathbf{V}_\alpha$. The well-foundedness of \in (restricted to \mathbf{U}) implies that in fact this must be all of \mathbf{U} .

Clearly, every set of rank less than κ is in fact a \mathbf{U} -set, and \mathbf{U} is itself a set of rank κ . The cardinality of \mathbf{U} is also κ , since κ is a regular cardinal number and any cardinal number less than κ is a member of \mathbf{U} .

(ii). We may construct \mathbf{V}_κ using the same method as in (i). By construction \mathbf{V}_κ satisfies axiom 1; since κ is infinite, \mathbf{V}_κ satisfies axioms 2 and 3; and since κ is strongly inaccessible, \mathbf{V}_κ satisfies axiom 4. Thus \mathbf{V}_κ is a pre-universe.

(iii). Again, let κ be the rank of \mathbf{U} . If $\kappa \in \mathbf{U}'$ then we can show by transfinite induction that $\mathbf{V}_\kappa \in \mathbf{U}'$ and so $\mathbf{U} \subsetneq \mathbf{U}'$; else we must have $\mathbf{U}' \subseteq \mathbf{V}_\kappa = \mathbf{U}$. ■

0.2 Accessibility and ind-completions

Prerequisites. § 0.1.

A classical technology for controlling size problems in category theory, due to Gabriel and Ulmer [1971], Grothendieck and Verdier [SGA 4a, Exposé I, § 9], and Makkai and Paré [1989], is the notion of accessibility. Though we make use of universes, accessibility remains important and is a crucial tool in verifying the stability of various universal constructions when one passes from one universe to a larger one.

Definition 0.2.1. Let κ be a regular cardinal. A κ -**filtered category** is a category \mathcal{J} satisfying these conditions:

- \mathcal{J} is **inhabited**, i.e. there exists an object in \mathcal{J} .
- If λ is a cardinal number strictly less than κ and S is a subset of $\text{ob } \mathcal{J}$ of cardinality λ , then there exist an object j and arrows $f_i : i \rightarrow j$ for each object i in S .

- If $f, g : i \rightarrow j$ are a pair of parallel arrows in \mathcal{J} , then there exist an object k and an arrow $h : j \rightarrow k$ such that $h \circ f = h \circ g$.

A κ -**directed preorder** is a preordered set that is κ -filtered when considered as a category; note that the third condition is then vacuous. A κ -**filtered diagram** (resp. κ -**directed diagram**) in a category \mathcal{C} is a functor $\mathbb{D} \rightarrow \mathcal{C}$ such that \mathbb{D} is a κ -filtered category (resp. κ -directed preorder). It is conventional to omit mention of κ when $\kappa = \aleph_0$.

Example 0.2.2. The category with one object $*$ and only one non-trivial arrow f is filtered if and only if $f = f \circ f$.

Example 0.2.3. Let X be any set. The set of all finite subsets of X , partially ordered by inclusion, is a directed preorder. More generally, if κ is any regular cardinal, then the set of all subsets of X with cardinality strictly less than κ is a κ -directed preorder.

Theorem 0.2.4. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let κ be any regular cardinal. Given a \mathbf{U} -small category \mathbb{D} , the following are equivalent:*

- (i) \mathbb{D} is a κ -filtered category.
- (ii) The functor $\varinjlim_{\mathbb{D}} : [\mathbb{D}, \mathbf{Set}] \rightarrow \mathbf{Set}$ preserves limits for all diagrams that are simultaneously κ -small and \mathbf{U} -small.

Proof. The claim (i) \Rightarrow (ii) is very well known, and the converse is an exercise in using the Yoneda lemma and manipulating limits and colimits for diagrams of representable functors; see Satz 5.2 in [Gabriel and Ulmer, 1971]. \square

Definition 0.2.5. Let κ be a regular cardinal in a universe \mathbf{U}^+ , let \mathbf{U} be a pre-universe with $\mathbf{U} \subseteq \mathbf{U}^+$, and let \mathbf{Set}^+ be the category of \mathbf{U}^+ -sets. A (κ, \mathbf{U}) -**compact object** in a locally \mathbf{U}^+ -small category \mathcal{C} is an object A such that the representable functor $\mathcal{C}(A, -) : \mathcal{C} \rightarrow \mathbf{Set}^+$ preserves colimits for all \mathbf{U} -small κ -filtered diagrams. A κ -**compact object** is one that is (κ, \mathbf{U}) -compact for all pre-universes \mathbf{U} .

Though the above definition is stated using a pre-universe \mathbf{U} contained in a universe \mathbf{U}^+ , the following lemma shows there is no dependence on \mathbf{U}^+ .

Lemma 0.2.6. *Let A be an object in a locally \mathbf{U}^+ -small category \mathcal{C} . The following are equivalent:*

- (i) A is a (κ, \mathbf{U}) -compact object in \mathcal{C} .
- (ii) *For all \mathbf{U} -small κ -filtered diagrams $B : \mathbb{D} \rightarrow \mathcal{C}$, if $\lambda : B \Rightarrow \Delta C$ is a colimiting cocone, then for any morphism $f : A \rightarrow C$, there exist an object i in \mathbb{D} and a morphism $f' : A \rightarrow B_i$ in \mathcal{C} such that $f = \lambda_i \circ f'$; and moreover if $f = \lambda_j \circ f''$ for some morphism $f'' : A \rightarrow B_j$ in \mathcal{C} , then there exists an object k and a pair of arrows $g : i \rightarrow k, h : j \rightarrow k$ in \mathbb{D} such that $B_g \circ f' = B_h \circ f''$.*

Proof. Use the explicit description of $\varinjlim_{\mathbb{D}} C(A, B)$ as a filtered colimit of sets; see Definition 1.1 in [LPAC], or Proposition 5.1.3 in [Borceux, 1994b]. \square

Corollary 0.2.7. *Let $B : \mathbb{D} \rightarrow \mathcal{C}$ be a \mathbf{U} -small κ -filtered diagram, and let $\lambda : B \Rightarrow \Delta C$ be a colimiting cocone in \mathcal{C} . If C is a (κ, \mathbf{U}) -compact object in \mathcal{C} , then for some object i in \mathbb{D} , $\lambda_i : B_i \rightarrow C$ is a split epimorphism. \blacksquare*

Lemma 0.2.8. *Let A be an object in a category \mathcal{C} .*

- (i) *If \mathbf{U} is a pre-universe contained in a universe \mathbf{U}^+ and κ is a regular cardinal such that A is (κ, \mathbf{U}^+) -compact, then A is (κ, \mathbf{U}) -compact as well.*
- (ii) *If κ is a regular cardinal such that A is (κ, \mathbf{U}) -compact and λ is any regular cardinal such that $\kappa \leq \lambda$, then A is also (λ, \mathbf{U}) -compact.*

Proof. Obvious. \blacklozenge

Lemma 0.2.9. *Let λ be a regular cardinal in a universe \mathbf{U}^+ , and let \mathbf{U} be a pre-universe with $\mathbf{U} \subseteq \mathbf{U}^+$. If $B : \mathbb{D} \rightarrow \mathcal{C}$ is a λ -small diagram of (λ, \mathbf{U}) -compact objects in a locally \mathbf{U}^+ -small category, then the colimit $\varinjlim_{\mathbb{D}} B$, if it exists, is a (λ, \mathbf{U}) -compact object in \mathcal{C} .*

Proof. Use [theorem 0.2.4](#) and the fact that $C(-, C) : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}^+$ maps colimits in \mathcal{C} to limits in \mathbf{Set}^+ . \blacksquare

Corollary 0.2.10. *A retract of a (λ, \mathbf{U}) -compact object is also a (λ, \mathbf{U}) -compact object.*

Proof. Suppose $r : A \rightarrow B$ and $s : B \rightarrow A$ are morphisms in C such that $r \circ s = \text{id}_B$. Then $e = s \circ r$ is an idempotent morphism and the diagram below

$$A \begin{array}{c} \xrightarrow{\text{id}_A} \\ \xrightarrow{e} \end{array} A \xrightarrow{r} B$$

is a (split) coequaliser diagram in C , so B is (λ, \mathbf{U}) -compact if A is. ■

Proposition 0.2.11. *Let \mathbf{U} be a pre-universe and let \mathbf{Set} be the category of \mathbf{U} -sets. For any \mathbf{U} -set A , the following are equivalent:*

- (i) A has cardinality less than κ .
- (ii) The representable functor $\mathbf{Set}(A, -) : \mathbf{Set} \rightarrow \mathbf{Set}$ preserves colimits for all \mathbf{U} -small κ -filtered diagrams.
- (iii) The representable functor $\mathbf{Set}(A, -) : \mathbf{Set} \rightarrow \mathbf{Set}$ preserves colimits for all \mathbf{U} -small κ -directed diagrams.

Proof. The claim (i) \Rightarrow (ii) follows from [theorem 0.2.4](#), and (ii) \Rightarrow (iii) is obvious. To see (iii) \Rightarrow (i), we may use [corollary 0.2.7](#) and the fact that every set is the directed union of its subsets of cardinality at most κ . ■

Corollary 0.2.12. *A set is κ -compact if and only if its cardinality is $< \kappa$.* ■

Definition 0.2.13. Let κ be a regular cardinal in a universe \mathbf{U} . A κ -**accessible \mathbf{U} -category** is a locally \mathbf{U} -small category C satisfying the following conditions:

- C has colimits for all \mathbf{U} -small κ -filtered diagrams.
- There exists a \mathbf{U} -set \mathcal{G} such that every object in \mathcal{G} is (κ, \mathbf{U}) -compact and, for every object B in C , there exists a \mathbf{U} -small κ -filtered diagram of objects in \mathcal{G} with B as its colimit in C .

We write $\mathbf{K}_\kappa^\mathbf{U}(C)$ for the full subcategory of C spanned by the (κ, \mathbf{U}) -compact objects.

Example 0.2.14. The category of \mathbf{U} -sets is a κ -accessible \mathbf{U} -category for any regular cardinal κ in \mathbf{U} .

REMARK 0.2.15. [Lemma 0.2.9](#) implies that, for each object A in an accessible \mathbf{U} -category, there exists a regular cardinal λ in \mathbf{U} such that A is (λ, \mathbf{U}) -compact.

Theorem 0.2.16. *Let \mathcal{C} be a locally \mathbf{U} -small category, and let κ be a regular cardinal in \mathbf{U} . There exist a locally \mathbf{U} -small category $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ and a functor $\gamma : \mathcal{C} \rightarrow \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ with the following properties:*

- (i) *The objects of $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ are \mathbf{U} -small κ -filtered diagrams $B : \mathbb{D} \rightarrow \mathcal{C}$, and γ sends an object C in \mathcal{C} to the corresponding trivial diagram $\mathbb{1} \rightarrow \mathcal{C}$ with value C .*
- (ii) *The functor $\gamma : \mathcal{C} \rightarrow \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ is fully faithful, injective on objects, preserves all limits that exist in \mathcal{C} , and preserves all κ -small colimits that exist in \mathcal{C} .*
- (iii) *$\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ has colimits for all \mathbf{U} -small κ -filtered diagrams.*
- (iv) *For every object C in \mathcal{C} , the object γC is (κ, \mathbf{U}) -compact in $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$, and for each \mathbf{U} -small κ -filtered diagram $B : \mathbb{D} \rightarrow \mathcal{C}$, there is a canonical colimiting cocone $\gamma B \Rightarrow \Delta B$ in $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$.*
- (v) *If \mathcal{D} is a category with colimits for all \mathbf{U} -small κ -filtered diagrams, then for each functor $F : \mathcal{C} \rightarrow \mathcal{D}$, there exists a functor $\bar{F} : \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C}) \rightarrow \mathcal{D}$ that preserves colimits for all \mathbf{U} -small κ -filtered diagrams in $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ such that $\gamma \bar{F} = F$, and given any functor $\bar{G} : \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C}) \rightarrow \mathcal{D}$ whatsoever, the induced map $\text{Nat}(\bar{F}, \bar{G}) \rightarrow \text{Nat}(F, \gamma \bar{G})$ is a bijection.*

The category $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ is called the **free (κ, \mathbf{U}) -ind-completion** of \mathcal{C} , or the **category of (κ, \mathbf{U}) -ind-objects** in \mathcal{C} .

Proof. If $B : \mathbb{D} \rightarrow \mathcal{C}$ and $B' : \mathbb{D}' \rightarrow \mathcal{C}$ are two \mathbf{U} -small κ -filtered diagrams, then properties (ii) and (iii) together imply that

$$\text{Hom}(B', B) \cong \lim_{\leftarrow \mathbb{D}'} \lim_{\rightarrow \mathbb{D}} C(B', B)$$

and so, taking the RHS as the *definition* of the LHS, we need only find a suitable notion of composition to make $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C})$ into a locally \mathbf{U} -small category. However, we observe that, if $N : \mathcal{C} \rightarrow [C^{\text{op}}, \mathbf{Set}]$ is the Yoneda embedding, then

$$\text{Hom}\left(\lim_{\leftarrow \mathbb{D}'} N B', \lim_{\rightarrow \mathbb{D}} N B\right) \cong \lim_{\leftarrow \mathbb{D}'} \lim_{\rightarrow \mathbb{D}} C(B', B)$$

and, assuming property (v), the Yoneda embedding $N : \mathcal{C} \rightarrow [C^{\text{op}}, \mathbf{Set}]$ must extend along γ to a functor $\bar{N} : \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{C}) \rightarrow [C^{\text{op}}, \mathbf{Set}]$ that preserves colimits

for \mathbf{U} -small κ -filtered diagram, so, in consideration of properties (i) and (iv), we may as well *define* the composition in $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(C)$ so that $\bar{\mathbf{N}}$ becomes fully faithful. This completes the definition of $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(C)$ as a category.

It remains to be shown that $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(C)$ actually has properties (ii), (iii), (iv), and (v); see Corollary 6.4.14 in [Borceux, 1994a] and Theorem 2.26 in [LPAC]. Note that the fact that γ preserves colimits for κ -small diagrams essentially follows from [theorem 0.2.4](#). \square

Proposition 0.2.17. *Let \mathbb{B} be a \mathbf{U} -small category and let κ be a regular cardinal in \mathbf{U} .*

- (i) $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ is a κ -accessible \mathbf{U} -category.
- (ii) Every (κ, \mathbf{U}) -compact object in $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ is a retract of an object of the form γB , where $\gamma : \mathbb{B} \rightarrow \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ is the canonical embedding.
- (iii) $\mathbf{K}_{\kappa}^{\mathbf{U}}(\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B}))$ is an essentially \mathbf{U} -small category.

Proof. (i). This claim more-or-less follows from the properties of $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ explained in the previous theorem.

(ii). Use [corollary 0.2.10](#).

(iii). Since \mathbb{B} is \mathbf{U} -small and $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ is locally \mathbf{U} -small, claim (ii) implies that $\mathbf{K}_{\kappa}^{\mathbf{U}}(\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B}))$ must be essentially \mathbf{U} -small. \blacksquare

Definition 0.2.18. Let κ be a regular cardinal in a universe \mathbf{U} . A (κ, \mathbf{U}) -**accessible functor** is a functor $F : C \rightarrow D$ such that

- C is a κ -accessible \mathbf{U} -category, and
- F preserves all colimits for \mathbf{U} -small κ -filtered diagrams.

We write $\mathbf{Acc}_{\kappa}^{\mathbf{U}}(C, D)$ for the full subcategory of the functor category $[C, D]$ spanned by the (κ, \mathbf{U}) -accessible functors. An **accessible functor** is a functor that is (κ, \mathbf{U}) -accessible functor for some regular cardinal κ in some universe \mathbf{U} .

Theorem 0.2.19 (Classification of accessible categories). *Let κ be a regular cardinal in a universe \mathbf{U} , and let C be a locally \mathbf{U} -small category. The following are equivalent:*

- (i) C is a κ -accessible \mathbf{U} -category.

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- (ii) *The inclusion $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \hookrightarrow \mathcal{C}$ extends along the embedding $\gamma : \mathcal{C} \rightarrow \mathbf{Ind}_{\mathbf{U}}^\kappa(\mathcal{C})$ to a (κ, \mathbf{U}) -accessible functor $\mathbf{Ind}_{\mathbf{U}}^\kappa(\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})) \rightarrow \mathcal{C}$ that is fully faithful and essentially surjective on objects.*
- (iii) *There exist a \mathbf{U} -small category \mathbb{B} and a functor $\mathbf{Ind}_{\mathbf{U}}^\kappa(\mathbb{B}) \rightarrow \mathcal{C}$ that is fully faithful and essentially surjective on objects.*

Proof. See Theorem 2.26 in [LPAC], or Theorem 5.3.5 in [Borceux, 1994b]. □

Corollary 0.2.20. *If \mathcal{C} is a κ -accessible \mathbf{U} -category and \mathcal{D} is any category, then:*

- (i) *The restriction $\mathbf{Acc}_\kappa^{\mathbf{U}}(\mathcal{C}, \mathcal{D}) \rightarrow [\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}), \mathcal{D}]$ is fully faithful and surjective on objects.*
- (ii) *In particular, if \mathcal{D} is also locally \mathbf{U} -small, then $\mathbf{Acc}_\kappa^{\mathbf{U}}(\mathcal{C}, \mathcal{D})$ is equivalent to a locally \mathbf{U} -small category.*
- (iii) *If \mathcal{D} has colimits for all \mathbf{U} -small κ -filtered diagrams, then the inclusion $\mathbf{Acc}_\kappa^{\mathbf{U}}(\mathcal{C}, \mathcal{D}) \hookrightarrow [\mathcal{C}, \mathcal{D}]$ has a left adjoint.* ■

Proposition 0.2.21. *Let \mathcal{C} be a κ -accessible \mathbf{U} -category and let \mathcal{D} be a locally \mathbf{U} -small category. Given an adjunction $F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$, if G is fully faithful and preserves colimits for all \mathbf{U} -small κ -filtered diagrams, then \mathcal{D} is also a κ -accessible \mathbf{U} -category.*

Proof. Under our hypotheses, given any \mathbf{U} -small κ -filtered diagram $A : \mathbb{J} \rightarrow \mathcal{D}$, we may take $F \varinjlim_{\mathbb{J}} GA$ as its colimit in \mathcal{D} . Our hypotheses also imply that F sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (κ, \mathbf{U}) -compact objects in \mathcal{D} ; thus if \mathcal{G} is a \mathbf{U} -small set of objects that generates \mathcal{C} under \mathbf{U} -small κ -filtered colimits, then $\{FX \mid X \in \mathcal{G}\}$ is a \mathbf{U} -small set of objects that generates \mathcal{D} in the same sense. ■

Definition 0.2.22. Let κ be a regular cardinal in a universe \mathbf{U} . A **locally κ -presentable \mathbf{U} -category** is a κ -accessible \mathbf{U} -category that is also \mathbf{U} -cocomplete. A **locally presentable \mathbf{U} -category** is one that is a locally κ -presentable \mathbf{U} -category for some regular cardinal κ in \mathbf{U} , and we often say ‘locally finitely presentable’ instead of ‘locally \aleph_0 -presentable’.

Example 0.2.23. The category of \mathbf{U} -sets is a locally κ -presentable \mathbf{U} -category for any regular cardinal κ in \mathbf{U} .

Lemma 0.2.24. *Let \mathcal{C} be a locally κ -presentable \mathbf{U} -category.*

- (i) *For any regular cardinal λ in \mathbf{U} , if $\kappa \leq \lambda$, then \mathcal{C} is a locally λ -presentable \mathbf{U} -category.*
- (ii) *With λ as above, if $F : \mathcal{C} \rightarrow \mathcal{D}$ is a (κ, \mathbf{U}) -accessible functor, then it is also a (λ, \mathbf{U}) -accessible functor.*
- (iii) *If \mathbf{U}^+ is any universe with $\mathbf{U} \in \mathbf{U}^+$, and \mathcal{C} is a locally κ -presentable \mathbf{U}^+ -category, then \mathcal{C} must be a preorder.*

Proof. (i). See the remark after Theorem 1.20 in [LPAC], or Propositions 5.3.2 and 5.2.3 in [Borceux, 1994b].

(ii). A λ -filtered diagram is certainly κ -filtered, so if F preserves colimits for all \mathbf{U} -small κ -filtered diagrams in \mathcal{C} , it must also preserve colimits for all \mathbf{U} -small λ -filtered diagrams.

(iii). This is a corollary of [theorem 0.1.10](#). ■

Corollary 0.2.25. *A category \mathcal{C} is a locally presentable \mathbf{U} -category for at most one universe \mathbf{U} , provided \mathcal{C} is not a preorder.*

Proof. Use [proposition 0.1.35](#) together with the above lemma. ■

Theorem 0.2.26 (Classification of locally presentable categories). *Let κ be a regular cardinal in a universe \mathbf{U} , let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{C} be a locally \mathbf{U} -small category. The following are equivalent:*

- (i) *\mathcal{C} is a locally κ -presentable \mathbf{U} -category.*
- (ii) *There exist a \mathbf{U} -small category \mathbb{B} that has colimits for κ -small diagrams and a functor $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B}) \rightarrow \mathcal{C}$ that is fully faithful and essentially surjective on objects.*
- (iii) *The restricted Yoneda embedding $\mathcal{C} \rightarrow [\mathbf{K}_{\kappa}^{\mathbf{U}}(\mathcal{C})^{\text{op}}, \mathbf{Set}]$ is fully faithful, (κ, \mathbf{U}) -accessible, and has a left adjoint.*
- (iv) *There exist a \mathbf{U} -small category \mathbb{A} and a fully faithful (κ, \mathbf{U}) -accessible functor $R : \mathcal{C} \rightarrow [\mathbb{A}, \mathbf{Set}]$ such that \mathbb{A} has limits for all κ -small diagrams, R has a left adjoint, and R is essentially surjective onto the full subcategory of functors $\mathbb{A} \rightarrow \mathbf{Set}$ that preserve limits for all κ -small diagrams.*

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- (v) *There exist a \mathbf{U} -small category \mathbb{A} and a fully faithful (κ, \mathbf{U}) -accessible functor $R : \mathcal{C} \rightarrow [\mathbb{A}, \mathbf{Set}]$ such that R has a left adjoint.*
- (vi) *\mathcal{C} is a κ -accessible \mathbf{U} -category and is \mathbf{U} -complete.*

Proof. See Proposition 1.27, Corollary 1.28, Theorem 1.46, and Corollary 2.47 in [LPAC], or Theorems 5.2.7 and 5.5.8 in [Borceux, 1994b]. \square

REMARK 0.2.27. If \mathcal{C} is equivalent to $\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})$ for some \mathbf{U} -small category \mathbb{B} that has colimits for all κ -small diagrams, then \mathbb{B} must be equivalent to $\mathbf{K}_{\kappa}^{\mathbf{U}}(\mathcal{C})$ by [proposition 0.2.17](#). In other words, every locally κ -presentable \mathbf{U} -category is, up to equivalence, the (κ, \mathbf{U}) -ind-completion of an essentially unique \mathbf{U} -small κ -cocomplete category.

Example 0.2.28. Obviously, for any \mathbf{U} -small category \mathbb{A} , the functor category $[\mathbb{A}, \mathbf{Set}]$ is locally finitely presentable. More generally, one may show that for any κ -ary algebraic theory \mathbf{T} , possibly many-sorted, the category of \mathbf{T} -algebras in \mathbf{U} is a locally κ -presentable \mathbf{U} -category. The above theorem can also be used to show that \mathbf{Cat} , the category of \mathbf{U} -small categories, is a locally finitely presentable \mathbf{U} -small category.

Corollary 0.2.29. *Let \mathcal{C} be a locally κ -presentable \mathbf{U} -category. For any \mathbf{U} -small κ -filtered diagram \mathbb{D} , $\varinjlim_{\mathbb{D}} : [\mathbb{D}, \mathcal{C}] \rightarrow \mathcal{C}$ preserves κ -small limits.*

Proof. The claim is certainly true when $\mathcal{C} = [\mathbb{A}, \mathbf{Set}]$, by [theorem 0.2.4](#). In general, choose a (κ, \mathbf{U}) -accessible fully faithful functor $R : \mathcal{C} \rightarrow [\mathbb{A}, \mathbf{Set}]$ with a left adjoint, and simply note that R creates limits for all \mathbf{U} -small diagrams as well as colimits for all \mathbf{U} -small κ -filtered diagrams. \blacksquare

Proposition 0.2.30. *If \mathcal{C} is a locally κ -presentable \mathbf{U} -category and \mathbb{D} is any \mathbf{U} -small category, then the functor category $[\mathbb{D}, \mathcal{C}]$ is also a locally κ -presentable category.*

Proof. This can be proven using the classification theorem by noting that the 2-functor $[\mathbb{D}, -]$ preserves reflective subcategories, but see also Corollary 1.54 in [LPAC]. \square

It is commonplace to say ‘ λ -presentable object’ instead of ‘ λ -compact object’, especially in algebraic contexts. The following proposition justifies the alternative terminology:

Proposition 0.2.31. *Let \mathcal{C} be a locally κ -presentable \mathbf{U} -category, and let λ be a regular cardinal in \mathbf{U} with $\lambda \geq \kappa$. If \mathcal{H} is a small full subcategory of \mathcal{C} such that*

- *every (κ, \mathbf{U}) -compact object in \mathcal{C} is isomorphic to an object in \mathcal{H} , and*
- *\mathcal{H} is closed in \mathcal{C} under colimits for λ -small diagrams,*

then every (λ, \mathbf{U}) -compact object in \mathcal{C} is isomorphic to an object in \mathcal{H} . In particular, $\mathbf{K}_\lambda^{\mathbf{U}}(\mathcal{C})$ is the smallest replete full subcategory of \mathcal{C} containing $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ and closed in \mathcal{C} under colimits for λ -small diagrams.

TODO: Simplify
this argument.

Proof. Let C be any (λ, \mathbf{U}) -compact object in \mathcal{C} . Clearly, the comma category $(\mathcal{H} \downarrow \mathcal{C})$ is a \mathbf{U} -small λ -filtered category. Let $\mathcal{G} = \mathcal{H} \cap \mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$. One can show that $(\mathcal{G} \downarrow \mathcal{C})$ is a cofinal subcategory in $(\mathcal{H} \downarrow \mathcal{C})$, and the classification theorem (0.2.26) plus proposition A.5.24 implies that the tautological cocone on the diagram $(\mathcal{G} \downarrow \mathcal{C}) \rightarrow \mathcal{C}$ is colimiting, so the tautological cocone on the diagram $(\mathcal{H} \downarrow \mathcal{C}) \rightarrow \mathcal{C}$ is also colimiting. Now, by corollary 0.2.7, C is a retract of an object in \mathcal{H} , and hence C must be isomorphic to an object in \mathcal{H} , because \mathcal{H} is closed under coequalisers.

For the final claim, note that $\mathbf{K}_\lambda^{\mathbf{U}}(\mathcal{C})$ is certainly a replete full subcategory of \mathcal{C} and contained in any replete full subcategory containing $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ and closed in \mathcal{C} under colimits for λ -small diagrams, so we just have to show that $\mathbf{K}_\lambda^{\mathbf{U}}(\mathcal{C})$ is also closed in \mathcal{C} under colimits for λ -small diagrams; for this, we simply appeal to lemma 0.2.9. ■

Proposition 0.2.32. *Let \mathcal{C} be a locally κ -presentable \mathbf{U} -category and let \mathbb{D} be a μ -small category in \mathbf{U} . The (λ, \mathbf{U}) -compact objects in $[\mathbb{D}, \mathcal{C}]$ are precisely the diagrams $\mathbb{D} \rightarrow \mathcal{C}$ that are componentwise (λ, \mathbf{U}) -compact, so long as $\lambda \geq \max\{\kappa, \mu\}$.*

Proof. First, note that Mac Lane's subdivision category^[1] \mathbb{D}^\S is also μ -small, so $[\mathbb{D}, \mathcal{C}](A, B)$ is computed as the limit of a μ -small diagram of hom-sets. More precisely, using end notation,^[2]

$$[\mathbb{D}, \mathcal{C}](A, B) \cong \int_{d: \mathbb{D}} \mathcal{C}(Ad, Bd)$$

and so if A is componentwise (λ, \mathbf{U}) -compact, then $[\mathbb{D}, \mathcal{C}](A, -)$ preserves colimits for \mathbf{U} -small λ -filtered diagrams, hence A is itself (λ, \mathbf{U}) -compact.

[1] See [CWM, Ch. IX, § 5].

[2] See § A.6.

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Now, suppose A is a (λ, \mathbf{U}) -compact object in $[\mathbb{D}, C]$. Let d be an object in \mathbb{D} , let $d^* : [\mathbb{D}, C] \rightarrow C$ be evaluation at d , and let $d_* : C \rightarrow [\mathbb{D}, C]$ be the right adjoint, which is explicitly given by

$$(d_*C)(d') = \mathbb{D}(d', d) \pitchfork C$$

where \pitchfork is defined by following adjunction:

$$\mathbf{Set}(X, C(C, C')) \cong C(C, X \pitchfork C')$$

The unit $\eta_A : A \rightarrow d_*d^*A$ is constructed using the universal property of \pitchfork in the obvious way, and the counit $\varepsilon_C : d^*d_*C \rightarrow C$ is the projection $\mathbb{D}(d, d) \pitchfork C \rightarrow C$ corresponding to $\text{id}_d \in \mathbb{D}(d, d)$. Since C is a locally λ -presentable \mathbf{U} -category, there exist a \mathbf{U} -small λ -filtered diagram $B : \mathbb{J} \rightarrow C$ consisting of (λ, \mathbf{U}) -compact objects in C and a colimiting cocone $\alpha : B \Rightarrow \Delta d^*A$, and since each $\mathbb{D}(d', d)$ has cardinality less than μ , the cocone $d_*\alpha : d_*B \Rightarrow \Delta d_*d^*A$ is also colimiting, by [corollary 0.2.29](#). [Lemma 0.2.6](#) then implies $\eta_A : A \rightarrow d_*d^*A$ factors through $d_*\alpha_j : d_*(Bj) \rightarrow d_*d^*A$ for some j in \mathbb{J} , say

$$\eta_A = d_*\alpha_j \circ \sigma$$

for some $\sigma : A \rightarrow d_*Bj$. But then, by the triangle identity,

$$\text{id}_{Ad} = \varepsilon_{Ad} \circ d^*\eta_A = \varepsilon_{Ad} \circ d^*d_*\alpha_j \circ d^*\sigma = \alpha_j \circ \varepsilon_{Bj} \circ d^*\sigma$$

and so $\alpha_j : Bj \rightarrow Ad$ is a split epimorphism, hence Ad is a (λ, \mathbf{U}) -compact object, by [corollary 0.2.10](#). ■

REMARK 0.2.33. The claim in the above proposition can fail if $\mu > \lambda \geq \kappa$. For example, we could take $C = \mathbf{Set}$, with \mathbb{D} being the set ω considered as a discrete category; then the terminal object in $[\mathbb{D}, \mathbf{Set}]$ is componentwise finite, but is not itself an \aleph_0 -compact object in \mathbf{Set} .

Lemma 0.2.34. *Let κ and λ be regular cardinals in a universe \mathbf{U} , with $\kappa \leq \lambda$.*

- (i) *If \mathcal{D} is a locally λ -presentable \mathbf{U} -category, C is a locally \mathbf{U} -small category, and $G : \mathcal{D} \rightarrow C$ is a (λ, \mathbf{U}) -accessible functor that preserves limits for all \mathbf{U} -small diagrams in C , then, for any (κ, \mathbf{U}) -compact object C in C , the comma category $(C \downarrow G)$ has an initial object.*

- (ii) If \mathcal{C} is a locally κ -presentable \mathbf{U} -category, \mathcal{D} is a locally \mathbf{U} -small category, and $F : \mathcal{C} \rightarrow \mathcal{D}$ is a functor that preserves colimits for all \mathbf{U} -small diagrams in \mathcal{C} , then, for any object D in \mathcal{D} , the comma category $(F \downarrow D)$ has a terminal object.

Proof. (i). Let \mathcal{F} be the full subcategory of $(\mathcal{C} \downarrow G)$ spanned by those (D, g) where D is a (λ, \mathbf{U}) -compact object in \mathcal{D} . G preserves colimits for all \mathbf{U} -small λ -filtered diagrams, so, by lemma 0.2.6, \mathcal{F} must be a weakly initial family in $(\mathcal{C} \downarrow G)$. Proposition 0.2.17 implies \mathcal{F} is an essentially \mathbf{U} -small category, and since \mathcal{D} has limits for all \mathbf{U} -small diagrams and G preserves them, $(\mathcal{C} \downarrow G)$ is also \mathbf{U} -complete. Thus, the inclusion $\mathcal{F} \hookrightarrow (\mathcal{C} \downarrow G)$ has a limit, and it can be shown that this is an initial object in $(\mathcal{C} \downarrow G)$.^[3]

(ii). Let \mathcal{G} be the full subcategory of $(F \downarrow D)$ spanned by those (C, f) where C is a (κ, \mathbf{U}) -compact object in \mathcal{C} ; note that proposition 0.2.17 implies \mathcal{G} is an essentially \mathbf{U} -small category. Since \mathcal{C} has colimits for all \mathbf{U} -small diagrams and F preserves them, $(F \downarrow D)$ is also \mathbf{U} -cocomplete.^[4] Let (C, f) be a colimit for the inclusion $\mathcal{G} \hookrightarrow (F \downarrow D)$. It is not hard to check that (C, f) is a weakly terminal object in $(F \downarrow D)$, so the formal dual of Freyd's initial object lemma^[5] gives us a terminal object in $(F \downarrow D)$; explicitly, it may be constructed as the joint coequaliser of all the endomorphisms of (C, f) . ■

Theorem 0.2.35 (Accessible adjoint functor theorem). *Let κ and λ be regular cardinals in a universe \mathbf{U} , with $\kappa \leq \lambda$, let \mathcal{C} be a locally κ -presentable \mathbf{U} -category, and let \mathcal{D} be a locally λ -presentable \mathbf{U} -category.*

Given a functor $F : \mathcal{C} \rightarrow \mathcal{D}$, the following are equivalent:

- (i) F has a right adjoint $G : \mathcal{D} \rightarrow \mathcal{C}$, and G is a (λ, \mathbf{U}) -accessible functor.
- (ii) F preserves colimits for all \mathbf{U} -small diagrams and sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (λ, \mathbf{U}) -compact objects in \mathcal{D} .
- (iii) F has a right adjoint and sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (λ, \mathbf{U}) -compact objects in \mathcal{D} .

[3] See Theorem 1 in [CWM, Ch. X, § 2].

[4] See the Lemma in [CWM, Ch. V, § 6].

[5] See Theorem 1 in [CWM, Ch. V, § 6].

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On the other hand, given a functor $G : \mathcal{D} \rightarrow \mathcal{C}$, the following are equivalent:

- (iv) G has a left adjoint $F : \mathcal{C} \rightarrow \mathcal{D}$, and F sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (λ, \mathbf{U}) -compact objects in \mathcal{D} .
- (v) G is a (λ, \mathbf{U}) -accessible functor and preserves limits for all \mathbf{U} -small diagrams.
- (vi) G is a (λ, \mathbf{U}) -accessible functor and there exist a functor $F_0 : \mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \rightarrow \mathcal{D}$ and hom-set bijections

$$\mathcal{C}(C, GD) \cong \mathcal{D}(F_0 C, D)$$

natural in D for each (κ, \mathbf{U}) -compact object C in \mathcal{C} , where D varies in \mathcal{D} .

Proof. We will need to refer back to the details of the proof of this theorem later, so here is a sketch of the constructions involved.

(i) \Rightarrow (ii). If F is a left adjoint, then F certainly preserves colimits for all \mathbf{U} -small diagrams. Given a (κ, \mathbf{U}) -compact object C in \mathcal{C} and a \mathbf{U} -small λ -filtered diagram $B : \mathbb{J} \rightarrow \mathcal{D}$, observe that

$$\begin{aligned} \mathcal{D}\left(FC, \lim_{\mathbb{J}} B\right) &\cong \mathcal{C}\left(C, G \lim_{\mathbb{J}} B\right) \cong \mathcal{C}\left(C, \lim_{\mathbb{J}} GB\right) \\ &\cong \lim_{\mathbb{J}} \mathcal{C}(C, GB) \cong \lim_{\mathbb{J}} \mathcal{C}(FC, B) \end{aligned}$$

and thus FC is indeed a (λ, \mathbf{U}) -compact object in \mathcal{D} .

(ii) \Rightarrow (iii). It is enough to show that, for each object D in \mathcal{D} , the comma category $(F \downarrow D)$ has a terminal object (GD, ε_D) ;[6] but this was done in the previous lemma.

(iii) \Rightarrow (i). Given a (κ, \mathbf{U}) -compact object C in \mathcal{C} and a \mathbf{U} -small λ -filtered diagram $B : \mathbb{J} \rightarrow \mathcal{D}$, observe that

$$\begin{aligned} \mathcal{C}\left(C, G \lim_{\mathbb{J}} B\right) &\cong \mathcal{D}\left(FC, \lim_{\mathbb{J}} B\right) \cong \lim_{\mathbb{J}} \mathcal{C}(FC, B) \\ &\cong \lim_{\mathbb{J}} \mathcal{C}(C, GB) \cong \mathcal{C}\left(C, \lim_{\mathbb{J}} GB\right) \end{aligned}$$

[6] See Theorem 2 in [CWM, Ch. IV, § 1].

because FC is a (λ, \mathbf{U}) -compact object in \mathcal{D} ; but [theorem 0.2.26](#) says the restricted Yoneda embedding $\mathcal{C} \rightarrow [\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})^{\text{op}}, \mathbf{Set}]$ is fully faithful, so this is enough to conclude that G preserves colimits for \mathbf{U} -small λ -filtered diagrams.

(iv) \Rightarrow (v). If G is a right adjoint, then G certainly preserves limits for all \mathbf{U} -small diagrams; the rest of this implication is just (iii) \Rightarrow (i).

(v) \Rightarrow (vi). It is enough to show that, for each (κ, \mathbf{U}) -compact object C in \mathcal{C} , the comma category $(\mathcal{C} \downarrow G)$ has an initial object $(F_0 C, \eta_C)$; but this was done in the previous lemma. It is clear how to make F_0 into a functor $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \rightarrow \mathcal{D}$.

(vi) \Rightarrow (iv). We use [theorems 0.2.16](#) and [0.2.26](#) to extend $F_0 : \mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \rightarrow \mathcal{D}$ along the inclusion $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \hookrightarrow \mathcal{C}$ to get (κ, \mathbf{U}) -accessible functor $F : \mathcal{C} \rightarrow \mathcal{D}$. We then observe that, for any \mathbf{U} -small κ -filtered diagram $A : \mathbb{I} \rightarrow \mathcal{C}$ of (κ, \mathbf{U}) -compact objects in \mathcal{C} ,

$$\begin{aligned} \mathcal{C}\left(\lim_{\mathbb{I}} A, GD\right) &\cong \lim_{\mathbb{I}} \mathcal{C}(A, GD) \cong \lim_{\mathbb{I}} \mathcal{C}(F_0 A, D) \\ &\cong \mathcal{C}\left(\lim_{\mathbb{I}} F A, D\right) \cong \mathcal{C}\left(F \lim_{\mathbb{I}} A, D\right) \end{aligned}$$

is a series of bijections natural in D , where D varies in \mathcal{D} ; but \mathcal{C} is a locally κ -presentable \mathbf{U} -category, so this is enough to show that F is a left adjoint of G . The remainder of the claim is a corollary of (i) \Rightarrow (ii). \blacksquare

Corollary 0.2.36. *Let \mathcal{C} and \mathcal{D} be locally presentable \mathbf{U} -categories. If a functor $G : \mathcal{D} \rightarrow \mathcal{C}$ has a left adjoint, then there exists a regular cardinal μ in \mathbf{U} such that G is a (μ, \mathbf{U}) -accessible functor.*

Proof. Suppose \mathcal{C} is a locally κ -presentable \mathbf{U} -category, \mathcal{D} is a locally λ -presentable \mathbf{U} -category, and $F : \mathcal{C} \rightarrow \mathcal{D}$ is a left adjoint for G . Since $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ is an essentially \mathbf{U} -small category, recalling [lemma 0.2.8](#), there certainly exists a regular cardinal μ in \mathbf{U} such that $\mu \geq \lambda$ and F sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (μ, \mathbf{U}) -compact objects in \mathcal{D} . The above theorem, plus [lemma 0.2.24](#), implies G is an (μ, \mathbf{U}) -accessible functor. \blacksquare

0.3 Change of universe

Prerequisites. §§ [0.1](#), [0.2](#), [A.1](#), [A.5](#).

Having introduced universes into our ontology, it becomes necessary to ask whether an object with some universal property retains that property when we enlarge the universe. Though it sounds inconceivable, there do exist examples of badly-behaved constructions that are not stable under change-of-universe; for example, Waterhouse [1975] defined a functor $F : \mathbf{CRing} \rightarrow \mathbf{Set}^+$, where \mathbf{CRing} is the category of commutative rings in a universe \mathbf{U} and \mathbf{Set}^+ is the category of \mathbf{U}^+ -sets for some universe \mathbf{U}^+ with $\mathbf{U} \in \mathbf{U}^+$, such that the value of F at any given commutative ring in \mathbf{U} does not depend on \mathbf{U} , and yet the value of the fpqc sheaf associated with F at the field \mathbb{Q} depends on the size of \mathbf{U} .

Definition 0.3.1. Let κ be a regular cardinal in a universe \mathbf{U} , and let \mathbf{U}^+ be a universe with $\mathbf{U} \subseteq \mathbf{U}^+$. A $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -**accessible extension** is a (κ, \mathbf{U}) -accessible functor $i : \mathcal{C} \rightarrow \mathcal{C}^+$ such that

- \mathcal{C} is a κ -accessible \mathbf{U} -category,
- \mathcal{C}^+ is a κ -accessible \mathbf{U}^+ -category,
- i sends (κ, \mathbf{U}) -compact objects in \mathcal{C} to (κ, \mathbf{U}^+) -compact objects in \mathcal{C}^+ , and
- the functor $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \rightarrow \mathbf{K}_\kappa^{\mathbf{U}^+}(\mathcal{C}^+)$ so induced by i is fully faithful and essentially surjective on objects.

REMARK 0.3.2. Let \mathbb{B} be a \mathbf{U} -small category in which idempotents split. Then the (κ, \mathbf{U}) -accessible functor $\mathbf{Ind}_{\mathbf{U}}^\kappa(\mathbb{B}) \rightarrow \mathbf{Ind}_{\mathbf{U}^+}^\kappa(\mathbb{B})$ obtained by extending the embedding $\gamma^+ : \mathbb{B} \rightarrow \mathbf{Ind}_{\mathbf{U}^+}^\kappa(\mathbb{B})$ along $\gamma : \mathbb{B} \rightarrow \mathbf{Ind}_{\mathbf{U}}^\kappa(\mathbb{B})$ is a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -extension, by [proposition 0.2.17](#). The classification theorem ([0.2.19](#)) implies all examples of $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extensions are essentially of this form.

Proposition 0.3.3. *Let $i : \mathcal{C} \rightarrow \mathcal{C}^+$ be a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension.*

- (i) *\mathcal{C} is a locally κ -presentable \mathbf{U} -category if and only if \mathcal{C}^+ is a locally κ -presentable \mathbf{U}^+ -category.*
- (ii) *The functor $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is fully faithful.*
- (iii) *If $B : \mathcal{J} \rightarrow \mathcal{C}$ is any diagram (not necessarily \mathbf{U} -small) and \mathcal{C} has a limit for B , then i preserves this limit.*

Proof. (i). If \mathcal{C} is a locally κ -presentable \mathbf{U} -category, then $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ has colimits for all κ -small diagrams, so $\mathbf{K}_\kappa^{\mathbf{U}^+}(\mathcal{C}^+)$ also has colimits for all κ -small diagrams.

The classification theorem (0.2.19) then implies C^+ is a locally κ -presentable \mathbf{U}^+ -category. Reversing this argument proves the converse.

(ii). Let $A : \mathbb{I} \rightarrow C$ and $B : \mathbb{J} \rightarrow C$ be two \mathbf{U} -small κ -filtered diagrams of (κ, \mathbf{U}) -compact objects in C . Then,

$$\begin{aligned} C\left(\lim_{\mathbb{I}} A, \lim_{\mathbb{J}} B\right) &\cong \lim_{\mathbb{I}} \lim_{\mathbb{J}} C(A, B) \cong \lim_{\mathbb{I}} \lim_{\mathbb{J}} C^+(iA, iB) \\ &\cong C^+\left(\lim_{\mathbb{I}} iA, \lim_{\mathbb{J}} iB\right) \cong C^+\left(i \lim_{\mathbb{I}} A, i \lim_{\mathbb{J}} B\right) \end{aligned}$$

because i is (κ, \mathbf{U}) -accessible and is fully faithful on the subcategory $\mathbf{K}_\kappa^{\mathbf{U}}(C)$, and therefore $i : C \rightarrow C^+$ itself is fully faithful. Note that this hinges crucially on theorem 0.1.30.

(iii). Let $B : \mathcal{J} \rightarrow C$ be any diagram. We observe that, for any (κ, \mathbf{U}) -compact object C in C ,

$$\begin{aligned} C^+\left(iC, i \lim_{\mathcal{J}} B\right) &\cong C\left(C, \lim_{\mathcal{J}} B\right) && \text{because } i \text{ is fully faithful} \\ &\cong \lim_{\mathcal{J}} C(C, B) && \text{by definition of limit} \\ &\cong \lim_{\mathcal{J}} C^+(iC, iB) && \text{because } i \text{ is fully faithful} \end{aligned}$$

but we know the restricted Yoneda embedding $C^+ \rightarrow [\mathbf{K}_\kappa^{\mathbf{U}}(C)^{\text{op}}, \mathbf{Set}^+]$ is fully faithful, so this is enough to conclude that $i \lim_{\mathcal{J}} B$ is the limit of iB in C^+ . ■

REMARK 0.3.4. Similar methods show that any fully faithful functor $C \rightarrow C^+$ satisfying the four bulleted conditions in the definition above is necessarily (κ, \mathbf{U}) -accessible.

Lemma 0.3.5. *Let \mathbf{U} and \mathbf{U}^+ be universes, with $\mathbf{U} \in \mathbf{U}^+$, and let κ be a regular cardinal in \mathbf{U} . Suppose:*

- C and D are locally κ -presentable \mathbf{U} -categories.
- C^+ and D^+ are locally κ -presentable \mathbf{U}^+ -categories.
- $i : C \rightarrow C^+$ and $j : D \rightarrow D^+$ are $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extensions.

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Given a strictly commutative diagram of the form below,

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{j} & \mathcal{D}^+ \\ G \downarrow & & \downarrow G^+ \\ \mathcal{C} & \xrightarrow{i} & \mathcal{C}^+ \end{array}$$

where G is (κ, \mathbf{U}) -accessible, G^+ is (κ, \mathbf{U}^+) -accessible, if both have left adjoints, then the diagram satisfies the left Beck–Chevalley condition.

Proof. Let C be a (κ, \mathbf{U}) -compact object in \mathcal{C} . Inspecting the proof of [theorem 0.2.35](#), we see that the functor $(\mathcal{C} \downarrow G) \rightarrow (i\mathcal{C} \downarrow G^+)$ induced by j preserves initial objects. [Lemma A.1.7](#) says the component at C of the left Beck–Chevalley natural transformation $F^+i \Rightarrow jF$ is an isomorphism; but C is generated by $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ and the functors F, F^+, i, j all preserve colimits for \mathbf{U} -small κ -filtered diagrams, so in fact $F^+i \Rightarrow jF$ is a natural isomorphism. \blacksquare

Proposition 0.3.6. *If $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension and \mathcal{C} is a locally κ -presentable \mathbf{U} -category, then i preserves colimits for all \mathbf{U} -small diagrams in \mathcal{C} .*

Proof. It is well-known that a functor preserves colimits for all \mathbf{U} -small diagrams if and only if it preserves coequalisers for all parallel pairs and coproducts for all \mathbf{U} -small families, but coproducts for \mathbf{U} -small families can be constructed in a uniform way using coproducts for κ -small families and colimits for \mathbf{U} -small κ -filtered diagrams. It is therefore enough to show that $i : \mathcal{C} \rightarrow \mathcal{C}^+$ preserves all colimits for κ -small diagrams, since i is already (κ, \mathbf{U}) -accessible.

Let \mathbb{D} be a κ -small category. Recalling [proposition 0.1.12](#), our problem amounts to showing that the diagram

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{i} & \mathcal{C}^+ \\ \Delta \downarrow & & \downarrow \Delta^+ \\ [\mathbb{D}, \mathcal{C}] & \xrightarrow{i_*} & [\mathbb{D}, \mathcal{C}^+] \end{array}$$

satisfies the left Beck–Chevalley condition. It is clear that i_* is fully faithful. Colimits for \mathbf{U} -small diagrams in $[\mathbb{D}, \mathcal{C}]$ and in $[\mathbb{D}, \mathcal{C}^+]$ are computed componentwise, so Δ and i_* are certainly (κ, \mathbf{U}) -accessible, and Δ^+ is (κ, \mathbf{U}^+) -accessible. Using [proposition 0.2.32](#), we see that i_* is also a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension,

so we apply the lemma above to conclude that the left Beck–Chevalley condition is satisfied. \blacksquare

Theorem 0.3.7 (Stability of accessible adjoint functors). *Let \mathbf{U} and \mathbf{U}^+ be universes, with $\mathbf{U} \in \mathbf{U}^+$, and let κ and λ be regular cardinals in \mathbf{U} , with $\kappa \leq \lambda$. Suppose:*

- \mathcal{C} is a locally κ -presentable \mathbf{U} -category.
- \mathcal{D} is a locally λ -presentable \mathbf{U} -category.
- \mathcal{C}^+ is a locally κ -presentable \mathbf{U}^+ -category.
- \mathcal{D}^+ is a locally λ -presentable \mathbf{U}^+ -category.

Let $i : \mathcal{C} \rightarrow \mathcal{C}^+$ be a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension and let $j : \mathcal{D} \rightarrow \mathcal{D}^+$ be a fully faithful functor.

- (i) Given a strictly commutative diagram of the form below,

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{j} & \mathcal{D}^+ \\ G \downarrow & & \downarrow G^+ \\ \mathcal{C} & \xrightarrow{i} & \mathcal{C}^+ \end{array}$$

where G is (λ, \mathbf{U}) -accessible and G^+ is (λ, \mathbf{U}^+) -accessible, if both have left adjoints and j is a $(\lambda, \mathbf{U}, \mathbf{U}^+)$ -accessible extension, then the diagram satisfies the left Beck–Chevalley condition.

- (ii) Given a strictly commutative diagram of the form below,

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{i} & \mathcal{C}^+ \\ F \downarrow & & \downarrow F^+ \\ \mathcal{D} & \xrightarrow{j} & \mathcal{D}^+ \end{array}$$

if both F and F^+ have right adjoints, then the diagram satisfies the right Beck–Chevalley condition.

Proof. (i). The proof is essentially the same as [lemma 0.3.5](#), though we have to use [proposition 0.3.6](#) to ensure that j preserves colimits for all \mathbf{U} -small κ -filtered diagrams in \mathcal{C} .

(ii). Let D be any object in \mathcal{D} . Inspecting the proof of [theorem 0.2.35](#), we see that our hypotheses, plus the fact that i preserves colimits for all \mathbf{U} -small diagrams in \mathcal{C} , imply that the functor $(F \downarrow D) \rightarrow (F^+ \downarrow jD)$ induced by i preserves terminal objects. Thus, [lemma A.1.7](#) implies that the diagram satisfies the right Beck–Chevalley condition. \blacksquare

Theorem 0.3.8. *If $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension and \mathcal{C} is a locally κ -presentable \mathbf{U} -category, then:*

- (i) *If λ is a regular cardinal and $\kappa \leq \lambda \in \mathbf{U}$, then $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is also a $(\lambda, \mathbf{U}, \mathbf{U}^+)$ -accessible extension.*
- (ii) *If μ is the cardinality of \mathbf{U} , then $i : \mathcal{C} \rightarrow \mathcal{C}^+$ factors through the inclusion $\mathbf{K}_\mu^{\mathbf{U}^+}(\mathcal{C}^+) \hookrightarrow \mathcal{C}^+$ as functor $\mathcal{C} \rightarrow \mathbf{K}_\mu^{\mathbf{U}^+}(\mathcal{C}^+)$ that is (fully faithful and) essentially surjective on objects.*
- (iii) *The (μ, \mathbf{U}^+) -accessible functor $\mathbf{Ind}_{\mathbf{U}^+}^\mu(\mathcal{C}) \rightarrow \mathcal{C}^+$ induced by $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is fully faithful and essentially surjective on objects.*

Proof. (i). Since $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is a (κ, \mathbf{U}) -accessible functor, it is certainly also (λ, \mathbf{U}) -accessible, by [lemma 0.2.24](#). It is therefore enough to show that i restricts to a functor $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C}) \rightarrow \mathbf{K}_\kappa^{\mathbf{U}^+}(\mathcal{C}^+)$ that is (fully faithful and) essentially surjective on objects.

[Proposition 0.2.31](#) says $\mathbf{K}_\lambda^{\mathbf{U}}(\mathcal{C})$ is the smallest replete full subcategory of \mathcal{C} that contains $\mathbf{K}_\kappa^{\mathbf{U}}(\mathcal{C})$ and is closed in \mathcal{C} under colimits for λ -small diagrams, therefore the replete closure of the image of $\mathbf{K}_\lambda^{\mathbf{U}}(\mathcal{C})$ must be the smallest replete full subcategory of \mathcal{C}^+ that contains $\mathbf{K}_\kappa^{\mathbf{U}^+}(\mathcal{C}^+)$ and is closed in \mathcal{C}^+ under colimits for λ -small diagrams, since i is fully faithful and preserves colimits for all \mathbf{U} -small diagrams. This proves the claim.

(ii). Since every object in \mathcal{C} is (λ, \mathbf{U}) -compact for some regular cardinal $\lambda < \mu$, claim (i) implies that the image of $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is contained in $\mathbf{K}_\mu^{\mathbf{U}^+}(\mathcal{C})$. To show i is essentially surjective onto $\mathbf{K}_\mu^{\mathbf{U}^+}(\mathcal{C})$, we simply have to observe that the inaccessibility of μ ([proposition 0.1.35](#)) and [proposition 0.2.31](#) imply that, for \mathcal{C}' any (μ, \mathbf{U}^+) -compact object in \mathcal{C}^+ , there exists a regular cardinal $\lambda < \mu$ such that \mathcal{C}' is also a (λ, \mathbf{U}^+) -compact object, which reduces the question to claim (i).

(iii). This is an immediate corollary of claim (ii) and the classification theorem ([0.2.19](#)) applied to \mathcal{C}^+ , considered as a (μ, \mathbf{U}^+) -accessible category. \blacksquare

REMARK 0.3.9. Although the fact $i : \mathcal{C} \rightarrow \mathcal{C}^+$ that preserves limits and colimits for all \mathbf{U} -small diagrams in \mathcal{C} is a formal consequence of the theorem above (via e.g. [corollary A.5.29](#)), it is not clear whether the theorem can be proven without already knowing this.

Corollary 0.3.10. *If \mathbb{B} is a \mathbf{U} -small category and has colimits for all κ -small diagrams, and μ is the cardinality of \mathbf{U} , then the canonical (μ, \mathbf{U}^+) -accessible functor $\mathbf{Ind}_{\mathbf{U}^+}^{\mu}(\mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathbb{B})) \rightarrow \mathbf{Ind}_{\mathbf{U}^+}^{\kappa}(\mathbb{B})$ is fully faithful and essentially surjective on objects. \blacksquare*

Theorem 0.3.11 (Stability of pointwise Kan extensions). *Let $F : \mathcal{A} \rightarrow \mathcal{C}$ and $G : \mathcal{A} \rightarrow \mathcal{D}$ be functors, and let $i : \mathcal{C} \rightarrow \mathcal{C}^+$ and $j : \mathcal{D} \rightarrow \mathcal{D}^+$ be fully faithful functors. Consider the following (not necessarily commutative) diagram:*

$$\begin{array}{ccccc}
 \mathcal{A} & \xrightarrow{G} & \mathcal{D} & \xrightarrow{j} & \mathcal{D}^+ \\
 \downarrow F & \nearrow H & & \nearrow H^+ & \\
 \mathcal{C} & & & & \\
 \downarrow i & & & & \\
 \mathcal{C}^+ & & & &
 \end{array}$$

- (i) *If H^+ is a pointwise right Kan extension of jG along iF , and $H^+i \cong jH$, then H is a pointwise right Kan extension of G along F .*
- (ii) *Suppose jH is a pointwise right Kan extension of jG along F . If H^+ is a pointwise right Kan extension of jH along i , then the counit $H^+i \Rightarrow jH$ is a natural isomorphism, and H^+ is also a pointwise right Kan extension of jG along iF ; conversely, if H^+ is a pointwise right Kan extension of jG along iF , then it is also a pointwise right Kan extension of jH along i .*
- (iii) *If \mathbf{U} is a pre-universe such that \mathcal{A} is \mathbf{U} -small and j preserves limits for all \mathbf{U} -small diagrams, and H is a pointwise right Kan extension of G along F , then a pointwise right Kan extension of jG along iF can be computed as a pointwise right Kan extension of jH along i (if either one exists).*

Dually:

- (i') *If H^+ is a pointwise left Kan extension of jG along iF , and $H^+i \cong jH$, then H is a pointwise left Kan extension of G along F .*

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(ii') Suppose jH is a pointwise left Kan extension of jG along F . If H^+ is a pointwise right Kan extension of jH along i , then the unit $jH \Rightarrow H^+i$ is a natural isomorphism, and H^+ is also a pointwise left Kan extension of jG along iF ; conversely, if H^+ is a pointwise left Kan extension of jG along iF , then it is also a pointwise left Kan extension of jH along i .

(iii') If \mathbf{U} is a pre-universe such that \mathcal{A} is \mathbf{U} -small and j preserves colimits for all \mathbf{U} -small diagrams, and H is a pointwise left Kan extension of G along F , then a pointwise left Kan extension of jG along iF can be computed as a pointwise left Kan extension of jH along i (if either one exists).

Proof. (i). [Theorem A.5.15](#) gives an explicit description of $H^+ : C^+ \rightarrow D^+$ as a weighted limit:

$$H^+(C') \cong \{C^+(C', iF), jG\}^{\mathcal{A}}$$

Since i is fully faithful, the weights $C(C, F)$ and $C^+(iC, iF)$ are naturally isomorphic, hence,

$$jH(C) \cong H^+(iC) \cong \{C^+(iC, iF), jG\}^{\mathcal{A}} \cong \{C(C, F), jG\}^{\mathcal{A}}$$

but, since j is fully faithful, j reflects *all* weighted limits, therefore H must be a pointwise right Kan extension of G along F .

(ii). Let \mathbf{U}^+ be a pre-universe such that \mathcal{A} and \mathcal{C} are \mathbf{U}^+ -small categories and \mathcal{D}, C^+, D^+ are locally \mathbf{U}^+ -small categories, and let \mathbf{Set}^+ be the category of \mathbf{U}^+ -sets. Using the interchange law ([theorem A.6.13](#)) and propositions [A.6.7](#) and [A.6.14](#), we obtain the following natural bijections:

$$\begin{aligned} D^+(D', H^+(C')) &\cong D^+(D', \{C^+(C', i), jH\}^{\mathcal{C}}) \\ &\cong \int_{C:\mathcal{C}} \mathbf{Set}^+(C^+(C', iC), D^+(D', jHC)) \\ &\cong \int_{C:\mathcal{C}} \mathbf{Set}^+(C^+(C', iC), D^+(D', \{C(C, F), jG\}^{\mathcal{A}})) \\ &\cong \int_{C:\mathcal{C}} \int_{A:\mathcal{A}} \mathbf{Set}^+(C^+(C', iC), \mathbf{Set}^+(C(C, FA), D^+(D', jGA))) \\ &\cong \int_{C:\mathcal{C}} \int_{A:\mathcal{A}} \mathbf{Set}^+(C(C, FA), \mathbf{Set}^+(C^+(C', iC), D^+(D', jGA))) \\ &\cong \int_{A:\mathcal{A}} \int_{C:\mathcal{C}} \mathbf{Set}^+(C(C, FA), \mathbf{Set}^+(C^+(C', iC), D^+(D', jGA))) \end{aligned}$$

$$\begin{aligned} &\cong \int_{A:A} \mathbf{Set}^+(C^+(C', iFA), D^+(D', jGA)) \\ &\cong D^+(D', \{C^+(C', iF), jG\}^A) \end{aligned}$$

Thus, H^+ is a pointwise right Kan extension of jG along iF if and only if H^+ is a pointwise right Kan extension of jH along i . The fact that the counit $H^+i \Rightarrow jH$ is a natural isomorphism is just [corollary A.5.19](#).

(iii). Apply [corollary A.5.18](#) to claim (ii). ■

Corollary 0.3.12. *Let \mathbf{U} and \mathbf{U}^+ be universes, with $\mathbf{U} \in \mathbf{U}^+$, and let κ and λ be regular cardinals in \mathbf{U} . Suppose:*

- C is a locally κ -presentable \mathbf{U} -category.
- D is a locally λ -presentable \mathbf{U} -category.
- C^+ is a locally κ -presentable \mathbf{U}^+ -category.
- D^+ is a locally λ -presentable \mathbf{U}^+ -category.

Let $F : \mathcal{A} \rightarrow C$ and $G : \mathcal{A} \rightarrow D$ be functors, let $i : C \rightarrow C^+$ be a $(\kappa, \mathbf{U}, \mathbf{U}^+)$ -accessible extension, and let $j : D \rightarrow D^+$ be a $(\lambda, \mathbf{U}, \mathbf{U}^+)$ -accessible extension. Consider the following (not necessarily commutative) diagram:

$$\begin{array}{ccccc} \mathcal{A} & \xrightarrow{G} & D & \xrightarrow{j} & D^+ \\ \downarrow F & \nearrow H & & \nearrow H^+ & \\ C & & & & \\ \downarrow i & & & & \\ C^+ & & & & \end{array}$$

- (i) *If H is a pointwise right Kan extension of G along F , then jH is a pointwise right Kan extension of jG along F , and if H^+ is a pointwise right Kan extension of jH along i , then H^+ is also a pointwise right Kan extension of jG along iF .*
- (ii) *Assuming \mathcal{A} is \mathbf{U} -small, if H is a pointwise left Kan extension of G along F , then jH is a pointwise left Kan extension of jG along F , and if H^+ is a pointwise left Kan extension of jH along i , then H^+ is also a pointwise left Kan extension of jG along iF .*

Proof. Use the theorem and the fact that i and j preserve limits for *all* diagrams and colimits for \mathbf{U} -small diagrams. ■

0.4 Small object arguments

Prerequisites. §§ 0.1, 0.2, 0.3, A.3.

The small object argument is a recurring construction in homotopical algebra, originally due to Quillen [1967, Ch. II, § 3] but refined by many authors since—notably by Garner [2009]. Roughly speaking, the small object argument shows that, under certain hypotheses, starting from a small set \mathcal{I} of morphisms in a cocomplete category \mathcal{C} , one can define the notions of ‘relative \mathcal{I} -cell complex’ and ‘ \mathcal{I} -fibration’ so that every morphism in \mathcal{C} factors as a relative \mathcal{I} -cell complex followed by an \mathcal{I} -fibration.

In this section, we will study the small object argument with a view toward questions of stability under change-of-universe.

Definition 0.4.1. Let \mathcal{C} be a category, and let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$. A **presentation for a relative \mathcal{I} -cell complex** in \mathcal{C} consists of the following data:

- An ordinal α . (We say the presentation is **indexed over α** .)
- A colimit-preserving functor $X_\bullet : [\alpha] \rightarrow \mathcal{C}$, where $[\alpha]$ is the well-ordered set $\{0, \dots, \alpha\}$ considered as a preorder category.
- For each ordinal $\beta < \alpha$, a (possibly empty) indexing set T_β ; and for each element j of T_β , a commutative diagram of the form below,

$$\begin{array}{ccc} U_{\beta,j} & \xrightarrow{u_{\beta,j}} & X_\beta \\ e_{\beta,j} \downarrow & & \downarrow X_{\beta \rightarrow \beta+1} \\ V_{\beta,j} & \xrightarrow{v_{\beta,j}} & X_{\beta+1} \end{array}$$

where $e_{\beta,j} : U_{\beta,j} \rightarrow V_{\beta,j}$ is a morphism in \mathcal{I} .

These data are moreover required to satisfy the following condition:

- For each ordinal $\beta < \gamma$, the coproducts $\coprod_{j \in T_\beta} S_{\beta,j}$ and $\coprod_{j \in T_\beta} D_{\beta,j}$ exist in \mathcal{C} , and the induced diagram

$$\begin{array}{ccc} \coprod_{j \in T_\beta} U_{\beta,j} & \xrightarrow{u_\beta} & X_\beta \\ \coprod_{j \in T_\beta} e_{\beta,j} \downarrow & & \downarrow X_{\beta \rightarrow \beta+1} \\ \coprod_{j \in T_\beta} V_{\beta,j} & \xrightarrow{v_\beta} & X_{\beta+1} \end{array}$$

is a pushout square in \mathcal{C} .

The presentation is said to be **U-small** (resp. κ -**small** for a regular cardinal κ) if α is an ordinal in \mathbf{U} (resp. $|\alpha| < \kappa$) and the disjoint union $\coprod_{\beta < \alpha} T_\beta$ is in \mathbf{U} (resp. has cardinality less than κ). A **sequential presentation** is one where each T_β is a singleton, in which case we suppress the index j in $e_{\beta,j}$, $u_{\beta,j}$, and $v_{\beta,j}$.

A **relative \mathcal{I} -cell complex** in \mathcal{C} is a morphism $f : X \rightarrow Y$ in \mathcal{C} for which there exists a presentation as above with f equal to $X_0 \rightarrow X_\alpha$. Given an initial object 0 in \mathcal{C} , an **\mathcal{I} -cell complex** in \mathcal{C} is an object Y for which the unique morphism $0 \rightarrow Y$ is a relative \mathcal{I} -cell complex.

REMARK 0.4.2. For any object X in \mathcal{C} and any subset $\mathcal{I} \subseteq \text{mor } \mathcal{C}$, the morphism $\text{id} : X \rightarrow X$ is a relative \mathcal{I} -cell complex in \mathcal{C} , with the obvious presentation indexed over 0). More generally, every isomorphism in \mathcal{C} is a relative \mathcal{I} -cell complex, with a presentation indexed over 1 (and $T_0 = \emptyset$); but in order to get a *sequential* presentation, one must assume that there is an isomorphism in \mathcal{I} .

Proposition 0.4.3. *Let \mathcal{C} be a category, let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$, let κ be a regular cardinal, and let $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$ be the set of relative \mathcal{I} -cell complexes in \mathcal{C} that admit a κ -small presentation.*

- (i) *Every morphism in \mathcal{I} is also in $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$.*
- (ii) *For each object X in \mathcal{C} , the morphism $\text{id} : X \rightarrow X$ is in $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$.*
- (iii) *If $f : X \rightarrow Y$ and $g : Y \rightarrow Z$ are both in $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$, then so is $g \circ f$.*
- (iv) *Let α be an ordinal and let $X_\bullet : \alpha \rightarrow \mathcal{C}$ be a colimit-preserving functor. If $|\alpha| < \kappa$ and λ is a colimiting cocone from X_\bullet to Y and, for $\beta \leq \gamma < \alpha$, the morphism $X_{\beta \rightarrow \gamma} : X_\beta \rightarrow X_\gamma$ is in $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$, then each component $\lambda_\beta : X_\beta \rightarrow Y$ is also in $\text{cell}_{\mathcal{I},\kappa} \mathcal{C}$.*

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(v) Given a pushout diagram of the form below in \mathcal{C} ,

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

if g is in $\text{cell}_{\mathcal{L},\kappa} \mathcal{C}$ and \mathcal{C} has colimits for all κ -small diagrams, then f is also in $\text{cell}_{\mathcal{L},\kappa} \mathcal{C}$.

Proof. (i). Given any morphism $e : U \rightarrow V$ in \mathcal{L} , we have the following pushout diagram:

$$\begin{array}{ccc} U & \xrightarrow{\text{id}} & U \\ e \downarrow & & \downarrow e \\ V & \xrightarrow{\text{id}} & V \end{array}$$

Thus $e : U \rightarrow V$ is in $\text{cell}_{\mathcal{L}} \mathcal{C}$.

(ii). See [remark 0.4.2](#).

(iii). It is clear that appending any κ -small presentation for g to any κ -small presentation for f yields a κ -small presentation of $g \circ f$.

(iv). The case $\alpha = 0$ falls under claim (ii). If $\alpha = \gamma + 1$, then the component $\lambda_\gamma : X_\gamma \rightarrow Y$ must be an isomorphism, and thus $\lambda_\beta = \lambda_\gamma \circ X_{\beta \rightarrow \gamma}$ is also in $\text{cell}_{\mathcal{L}} \mathcal{C}$; and if α is a positive limit ordinal, since every terminal segment of α is cofinal in α , it is clear that concatenating κ -small presentations for $X_{\gamma \rightarrow \gamma+1}$ for $\beta \leq \gamma < \alpha$ yields a κ -small presentation for $\lambda_\beta : X_\beta \rightarrow Y$.

(v). Fix a κ -small presentation of $g : Z \rightarrow W$. By the pushout pasting lemma, given a commutative diagram of the form below,

$$\begin{array}{ccccc} \coprod_{j \in T_\beta} U_{\beta,j} & \xrightarrow{u_\beta} & Z_\beta & \longrightarrow & X_\beta \\ \coprod_{j \in T_\beta} e_{\beta,j} \downarrow & & \downarrow Z_{\beta \rightarrow \beta+1} & & \downarrow X_{\beta \rightarrow \beta+1} \\ \coprod_{j \in T_\beta} V_{\beta,j} & \xrightarrow{v_\beta} & Z_{\beta+1} & \longrightarrow & X_{\beta+1} \end{array}$$

if both squares are pushout diagrams, then the outer rectangle is a pushout diagram as well. Since pushout along $z : Z \rightarrow X$ is the left adjoint of the evident

functor $z^* : X/\mathcal{C} \rightarrow Z/\mathcal{C}$, it preserves all colimits, and thus we obtain a κ -small presentation of $f : X \rightarrow Y$. ■

Definition 0.4.4. Let \mathcal{C} be a category and let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$. An **\mathcal{I} -injective morphism** in \mathcal{C} is a morphism that has the right lifting property with respect to every morphism in \mathcal{I} .^[7] An **\mathcal{I} -cofibration** in \mathcal{C} is a morphism that has the left lifting property with respect to every \mathcal{I} -injective morphism.

Proposition 0.4.5. Let \mathcal{C} be a category, let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$, and let $\text{cell}_{\mathcal{I}} \mathcal{C}$, $\text{inj}^{\mathcal{I}} \mathcal{C}$, and $\text{cof}_{\mathcal{I}} \mathcal{C}$ be the set of relative \mathcal{I} -cell complexes, \mathcal{I} -injections, and \mathcal{I} -cofibrations in \mathcal{C} , respectively.

- (i) We have $\mathcal{I} \subseteq \text{cell}_{\mathcal{I}} \mathcal{C} \subseteq \text{cof}_{\mathcal{I}} \mathcal{C}$.
- (ii) A morphism is in $\text{inj}^{\mathcal{I}} \mathcal{C}$ if and only if it has the right lifting property with respect to every \mathcal{I} -cofibration.
- (iii) In particular, a morphism is in $\text{inj}^{\mathcal{I}} \mathcal{C}$ if and only if it has the right lifting property with respect to every relative \mathcal{I} -cell complex.

Proof. (i). Follows immediately from the definition of ‘relative \mathcal{I} -cell complex’ and [proposition A.3.12](#).

(ii) and (iii). See [proposition A.3.3](#). ■

Some authors define ‘relative \mathcal{I} -cell complex’ so that every such morphism admits a *sequential* presentation. The following lemma and its corollary show that there is no loss of generality in doing so.

Lemma 0.4.6. Let κ be a regular cardinal, let \mathcal{C} be a category with colimits for all κ -small diagrams, and let α be an ordinal of cardinality less than κ . For each ordinal $\beta < \alpha$, let $e_{\beta} : U_{\beta} \rightarrow V_{\beta}$ be a morphism in \mathcal{C} , and for each ordinal $\beta \leq \alpha$, let

$$C_{\beta} = \left(\coprod_{\gamma < \beta} V_{\gamma} \right) \amalg \left(\coprod_{\beta \leq \gamma < \alpha} U_{\gamma} \right)$$

be a coproduct in \mathcal{C} with coproduct insertions $u_{\gamma, \beta} : U_{\gamma} \rightarrow C_{\beta}$ (for $\beta \leq \gamma < \alpha$) and $v_{\gamma, \beta} : V_{\gamma} \rightarrow C_{\beta}$ (for $\gamma < \beta$).

[7] Equivalently, it is a morphism $f : X \rightarrow Y$ in \mathcal{C} that is an \mathcal{I} -injective object in the slice category $\mathcal{C}_{/Y}$.

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Given ordinals $\beta < \beta' \leq \alpha$, there is a unique morphism $C_\beta \rightarrow C_{\beta'}$, such that, for $\zeta < \beta \leq \zeta' < \beta' \leq \zeta''$, the following diagrams commute:

$$\begin{array}{ccc}
 V_\zeta \xrightarrow{v_{\zeta,\beta}} C_\beta & U_{\zeta'} \xrightarrow{u_{\zeta',\beta}} C_\beta & U_{\zeta''} \xrightarrow{u_{\zeta'',\beta}} C_\beta \\
 \text{id} \downarrow & e_{\zeta'} \downarrow & \text{id} \downarrow \\
 V_\zeta \xrightarrow{v_{\zeta,\beta'}} C_{\beta'} & V_{\zeta'} \xrightarrow{v_{\zeta',\beta'}} C_{\beta'} & U_{\zeta''} \xrightarrow{u_{\zeta'',\beta'}} C_{\beta'}
 \end{array}$$

This yields a functor $C_\bullet : [\alpha] \rightarrow \mathcal{C}$, and it preserves colimits. Moreover, the diagrams below are pushout squares for all ordinals $\beta < \alpha$:

$$\begin{array}{ccc}
 U_\beta \xrightarrow{u_{\beta,\beta}} C_\beta & & \\
 e_\beta \downarrow & & \downarrow \\
 V_\beta \xrightarrow{v_{\beta,\beta+1}} C_{\beta+1} & &
 \end{array}$$

Proof. This is a straightforward exercise. See Proposition 10.2.7 in [Hirschhorn, 2003]. \square

Corollary 0.4.7. Let κ be a regular cardinal, let \mathcal{C} be a category with colimits for κ -small diagrams, and let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$. If $f : X \rightarrow Y$ is a relative \mathcal{I} -cell complex in \mathcal{C} that admits a κ -small presentation, and either

- $X = Y$ and $f = \text{id}_X$, or
- f is an isomorphism and \mathcal{I} contains an isomorphism, or
- f is not an isomorphism,

then f also admits a κ -small sequential presentation.

Proof. We have already commented on the first two cases in [remark 0.4.2](#). The third case is proven by transfinite induction, where in the induction step we may assume that f is presented by just one pushout diagram:

$$\begin{array}{ccc}
 \coprod_{j \in T} U_j \xrightarrow{u} X & & \\
 \coprod_{j \in T} e_j \downarrow & & \downarrow f \\
 \coprod_{j \in T} V_j \xrightarrow{v} Y & &
 \end{array}$$

By decomposing the morphism $\coprod_{j \in T} e_j : \coprod_{j \in T} U_j \rightarrow \coprod_{j \in T} V_j$ as in the earlier lemma and applying the pushout pasting lemma, we obtain a sequential presentation of f , which is κ -small precisely if $|T| < \kappa$. ■

Definition 0.4.8. Let \mathbf{U} be a universe, let \mathcal{C} be a category, let \mathcal{I} be a subset of $\text{mor } \mathcal{C}$, and let $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$ be the set of relative \mathcal{I} -cell complexes in \mathcal{C} that have a \mathbf{U} -small presentation. We say $(\mathcal{I}, \mathcal{C})$ is **admissible for the \mathbf{U} -small object argument** when the following conditions are satisfied:

- \mathcal{I} is a \mathbf{U} -set.
- \mathcal{C} be a locally \mathbf{U} -small category with colimits for all \mathbf{U} -small diagrams.
- There is a regular cardinal κ in \mathbf{U} such that, for every morphism $e : U \rightarrow V$ in \mathcal{I} , every ordinal α in \mathbf{U} , and every functor $X_\bullet : \alpha \rightarrow \mathcal{C}$, if $|\alpha| \geq \kappa$, and the morphism $X_{\beta \rightarrow \gamma} : X_\beta \rightarrow X_\gamma$ is in $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$ for all ordinals $\beta \leq \gamma < \alpha$, then the canonical comparison map $\lim_{\rightarrow \beta < \alpha} \mathcal{C}(U, X_\beta) \rightarrow \mathcal{C}(U, \lim_{\rightarrow \beta < \alpha} X_\beta)$ is a bijection.

The **sequential \mathbf{U} -rank** of \mathcal{I} in \mathcal{C} is the least cardinal κ with the above property.

REMARK 0.4.9. Notice that, if $|\alpha| \geq \kappa$, then α is a κ -directed preorder. Thus, for any locally presentable \mathbf{U} -category \mathcal{C} and any \mathbf{U} -subset $\mathcal{I} \subseteq \text{mor } \mathcal{C}$ whatsoever, $(\mathcal{I}, \mathcal{C})$ is admissible for the \mathbf{U} -small object argument.

Definition 0.4.10. Let \mathbf{U} be a universe. A **\mathbf{U} -cofibrantly-generated factorisation system** on a category \mathcal{C} on is a weak factorisation system on \mathcal{C} that is cofibrantly generated by some \mathbf{U} -subset of $\text{mor } \mathcal{C}$.

Theorem 0.4.11 (Quillen's small object argument). *Let \mathbf{U} be a universe, let \mathcal{C} be a locally \mathbf{U} -small category with colimits for all \mathbf{U} -small diagrams, and let \mathcal{I} be a \mathbf{U} -subset of $\text{mor } \mathcal{C}$.*

- (i) *There exist a functor $M : [2, \mathcal{C}] \rightarrow \mathcal{C}$ and two natural transformations $i : \text{dom} \Rightarrow M$, $p : M \Rightarrow \text{codom}$ such that, for all morphisms $f : X \rightarrow Y$ in \mathcal{C} , the morphism $i_f : X \rightarrow M(f)$ is in $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$, and we have $f = p_f \circ i_f$.*
- (ii) *If $(\mathcal{I}, \mathcal{C})$ is moreover admissible for the \mathbf{U} -small object argument, then we may choose M , i , and p so that, for all morphisms $f : X \rightarrow Y$ in \mathcal{C} , the morphism $p_f : M(f) \rightarrow Y$ is in $\text{inj}^{\mathcal{I}} \mathcal{C}$.*

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(iii) *In particular, if $(\mathcal{I}, \mathcal{C})$ is admissible, then $(\text{cof}_{\mathcal{I}} \mathcal{C}, \text{inj}^{\mathcal{I}} \mathcal{C})$ is a \mathbf{U} -cofibrantly-generated factorisation system on \mathcal{C} and extends to a functorial weak factorisation system.*

Proof. (i). Let κ be any regular cardinal, and let α be the least ordinal of cardinality κ .^[8] For each morphism $f : X \rightarrow Y$ in \mathcal{C} , we construct by transfinite recursion a colimit-preserving functor $M_{\bullet}(f) : [\alpha] \rightarrow \mathcal{C}$ and a cocone $p_{f,\bullet} : M_{\bullet}(f) \rightarrow Y$ satisfying the following conditions:

- $M_0(f) = X$, $p_{f;0} = p$.
- For each ordinal $\beta < \alpha$, if $T_{\beta}(f)$ is the set of all commutative diagrams in \mathcal{C} of the form below,

$$\begin{array}{ccc} U_{\beta,j} & \xrightarrow{u_{\beta,j}} & M_{\beta}(f) \\ e_{\beta,j} \downarrow & & \downarrow p_{f;\beta} \\ V_{\beta,j} & \xrightarrow{v_{\beta,j}} & Y \end{array}$$

is in \mathcal{I} and $u_{\beta,j} : U_{\beta,j} \rightarrow X_{\beta}$ is in \mathcal{C} , then $T_{\beta}(f)$ is a \mathbf{U} -set (because \mathcal{I} is a \mathbf{U} -set and \mathcal{C} is a locally \mathbf{U} -small category), and we have a pushout square of the following form,

$$\begin{array}{ccc} \coprod_{j \in T_{\beta}(f)} U_{\beta,j} & \xrightarrow{u_{\beta}} & M_{\beta}(f) \\ \coprod_{j \in T_{\beta}(f)} e_{\beta,j} \downarrow & & \downarrow X_{\beta \rightarrow \beta+1} \\ \coprod_{j \in T_{\beta}(f)} V_{\beta,j} & \xrightarrow{\bar{v}_{\beta}} & M_{\beta+1}(f) \end{array}$$

where $u_{\beta} : \coprod_{j \in T_{\beta}(f)} U_{\beta,j} \rightarrow X_{\beta}$ is the evident morphism induced by the universal property of coproducts. Observe that there is then a unique morphism $p_{f;\beta+1} : M_{\beta+1}(f) \rightarrow Y$ such that

$$p_{f;\beta+1} \circ M_{\beta \rightarrow \beta+1}(f) = p_{\beta}$$

and

$$p_{f;\beta+1} \circ \bar{v}_{\beta,j} = v_{\beta,j}$$

for all j in $T_{\beta}(f)$, where $\bar{v}_{\beta,j} : V_{\beta,j} \rightarrow M_{\beta+1}(f)$ is the evident component of $\bar{v}_{\beta} : \coprod_{j \in T_{\beta}(f)} V_{\beta,j} \rightarrow M_{\beta+1}(f)$.

[8] We could also take $\kappa = 0$, but then the factorisation so obtained is trivial.

- For limit ordinals $\gamma \leq \alpha$, $M_\gamma(f) = \varinjlim_{\beta < \gamma} M_\beta(f)$, and $p_\gamma : M_\gamma(f) \rightarrow Y$ is defined by the universal property of X_γ .

It is not hard to see that the functor $M_\bullet(f) : [\alpha] \rightarrow \mathcal{C}$ so defined is itself functorial in f ; in particular, defining $M(f) = M_\alpha(f)$, $i_f = M_{0 \rightarrow \alpha}(f)$, $p_f = p_{f;\alpha}$, we obtain a functor $M : [\mathcal{2}, \mathcal{C}] \rightarrow \mathcal{C}$ with two natural transformations $i : M \Rightarrow \text{dom}$ and $p : M \Rightarrow \text{codom}$; by construction, we have $f = p_f \circ i_f$, and $i_f : X \rightarrow M(f)$ is in $\text{cell}_{\mathcal{I}, \mathcal{U}} \mathcal{C}$.

(ii). Now, take κ to be a regular cardinal as in [definition 0.4.8](#). We wish to show that the morphism p_f constructed above has the right lifting property with respect to all morphisms in \mathcal{I} . Consider a lifting problem of the form below,

$$\begin{array}{ccc} U & \xrightarrow{u} & M(f) \\ e \downarrow & & \downarrow p_f \\ V & \xrightarrow{v} & Y \end{array}$$

where $e : U \rightarrow V$ is in \mathcal{I} . Since \mathcal{I} is admissible, there must exist an ordinal $\beta < \alpha$ and a morphism $u' : U \rightarrow M_\beta(f)$ such that $u = M_{\beta \rightarrow \alpha}(f) \circ u'$. We then obtain the following commutative diagram:

$$\begin{array}{ccc} U & \xrightarrow{u'} & M_\beta(f) \\ e \downarrow & & \downarrow p_{f;\beta} \\ V & \xrightarrow{v} & Y \end{array}$$

Since this is one of the diagrams in the set $T_\beta(f)$, it must embed in a commutative diagram of the form below,

$$\begin{array}{ccccc} U & \xrightarrow{u'} & M_\beta(f) & \longrightarrow & M_\alpha(f) \\ e \downarrow & & \downarrow & \nearrow & \downarrow p_{f;\alpha} \\ V & \longrightarrow & M_{\beta+1}(f) & & \\ \parallel & & & & \\ V & \xrightarrow{\quad\quad\quad} & & \xrightarrow{v} & Y \end{array}$$

and thus we have the required lift $V \rightarrow M(f)$.

(iii). Finally, apply [proposition 0.4.5](#) and [theorem A.3.28](#). ■

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Corollary 0.4.12. *With other notation in the theorem, a morphism $g : Z \rightarrow W$ is in $\text{cof}_I \mathcal{C}$ if and only if there exists a commutative diagram of the following form in \mathcal{C} ,*

$$\begin{array}{ccc} Z & \xrightarrow{i} & W' \\ g \downarrow & \nearrow & \downarrow \\ W & \xrightarrow{\text{id}} & W \end{array}$$

where $i : Z \rightarrow W'$ is in $\text{cell}_{I,U} \mathcal{C}$.

Proof. (i). If $g : Z \rightarrow W$ is in $\text{cof}_I \mathcal{C}$, then g has the left lifting property with respect to $p_g : M(g) \rightarrow W$, and so there exists a commutative diagram of the required form. Conversely, suppose we have $g = p \circ i$, $i = j \circ g$, and $\text{id}_W = p \circ j$ for some $i : Z \rightarrow W'$ in $\text{cell}_{I,U} \mathcal{C}$ and some $j : W \rightarrow W'$ in \mathcal{C} . Then g is a retract of i ,

$$\begin{array}{ccccc} Z & \xrightarrow{\text{id}} & Z & \xrightarrow{\text{id}} & Z \\ g \downarrow & & i \downarrow & & \downarrow g \\ W & \xrightarrow{j} & W' & \xrightarrow{p} & W \\ & \searrow & \text{id} & \nearrow & \\ & & & & \end{array}$$

but [proposition 0.4.5](#) says i is in $\text{cof}_I \mathcal{C}$, so by [proposition A.3.12](#), g is also in $\text{cof}_I \mathcal{C}$. ■

Lemma 0.4.13. *Let \mathcal{C} be a full subcategory of a category \mathcal{C}^+ , let I be a subset of $\text{mor } \mathcal{C}$, and let κ be a regular cardinal. If \mathcal{C} is closed in \mathcal{C}^+ under colimits for all κ -small diagrams, then $\text{cell}_{I,\kappa} \mathcal{C} = \text{cell}_{I,\kappa} \mathcal{C}^+ \cap \text{mor } \mathcal{C}$.*

Proof. Obvious. ◆

Theorem 0.4.14 (Stability of cofibrantly-generated factorisation systems). *Let \mathbf{U} and \mathbf{U}^+ be universes, with $\mathbf{U} \in \mathbf{U}^+$. Suppose:*

- \mathcal{C} is a locally \mathbf{U} -small and \mathbf{U} -cocomplete category.
- \mathcal{C}^+ is a locally \mathbf{U}^+ -small and \mathbf{U}^+ -cocomplete category.
- The inclusion $\mathcal{C} \hookrightarrow \mathcal{C}^+$ preserves colimits for all \mathbf{U} -small diagrams.
- I is a \mathbf{U} -subset of $\text{mor } \mathcal{C}$.

- $(\mathcal{I}, \mathcal{C})$ is admissible for the \mathbf{U} -small object argument, and (L, R) is the functorial factorisation system on \mathcal{C} constructed by Quillen's small object argument.
- $(\mathcal{I}, \mathcal{C}^+)$ is admissible for the \mathbf{U}^+ -small object argument, and (L^+, R^+) is the functorial factorisation system on \mathcal{C}^+ constructed by Quillen's small object argument.

Under these hypotheses, if the sequential \mathbf{U} -rank of \mathcal{I} in \mathcal{C} is equal to the sequential \mathbf{U}^+ -rank of \mathcal{I} in \mathcal{C}^+ , then:

- (i) For each morphism $f : X \rightarrow Y$ in \mathcal{C} , we have a commutative diagram of the following form in \mathcal{C}^+ ,

$$\begin{array}{ccc}
 & M^+(f) & \\
 L^+f \nearrow & \downarrow \cong & \searrow R^+f \\
 X & & Y \\
 Lf \searrow & & \nearrow Rf \\
 & M(f) &
 \end{array}$$

and the isomorphism $M^+(f) \rightarrow M(f)$ is moreover canonical and natural in f .

- (ii) We have $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C} \subseteq \text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}^+ \subseteq \text{cell}_{\mathcal{I}, \mathbf{U}^+} \mathcal{C}^+$.
- (iii) $(\text{cof}_{\mathcal{I}} \mathcal{C}^+, \text{inj}^{\mathcal{I}} \mathcal{C}^+)$ is an extension of $(\text{cof}_{\mathcal{I}} \mathcal{C}, \text{inj}^{\mathcal{I}} \mathcal{C})$.

Proof. (i). This can be seen by examining the explicit construction in the proof of [theorem 0.4.11](#).

(ii). This is implied by the lemma.

(iii). Since $(\text{cof}_{\mathcal{I}} \mathcal{C}, \text{inj}^{\mathcal{I}} \mathcal{C})$ and $(\text{cof}_{\mathcal{I}} \mathcal{C}^+, \text{inj}^{\mathcal{I}} \mathcal{C}^+)$ are both cofibrantly generated by \mathcal{I} , by [proposition A.3.18](#), we have $\text{inj}^{\mathcal{I}} \mathcal{C} \subseteq \text{inj}^{\mathcal{I}} \mathcal{C}^+$ and so $\text{cof}_{\mathcal{I}} \mathcal{C} \supseteq \text{cof}_{\mathcal{I}} \mathcal{C}^+ \cap \text{mor } \mathcal{C}$. It remains to be shown that $\text{cof}_{\mathcal{I}} \mathcal{C} \subseteq \text{cof}_{\mathcal{I}} \mathcal{C}^+$, but this is implied by [corollary 0.4.12](#) applied to claim (ii). ■

REMARK 0.4.15. Let κ be a regular cardinal in \mathbf{U} , let \mathcal{B} be a \mathbf{U} -small category with colimits for all κ -small diagrams, let $\mathcal{C} = \mathbf{Ind}_{\mathbf{U}}^{\kappa}(\mathcal{B})$, and let $\mathcal{C}^+ = \mathbf{Ind}_{\mathbf{U}^+}^{\kappa}(\mathcal{B})$. Then \mathcal{C} is a locally κ -presentable \mathbf{U} -category, the inclusion $\mathcal{C} \hookrightarrow \mathcal{C}^+$ is an accessible

$(\kappa, \mathbf{U}, \mathbf{U}^+)$ extension, and any \mathbf{U} -subset $\mathcal{I} \subseteq \text{mor } \mathcal{C}$ whatsoever will satisfy the hypotheses of the theorem.

Proposition 0.4.16. *Let $F \dashv U : \mathcal{D} \rightarrow \mathcal{C}$ be an adjunction of categories, let $\mathcal{I} \subseteq \text{mor } \mathcal{C}$, and let $\mathcal{J} = \{Ff \mid f \in \mathcal{I}\}$.*

- (i) *F sends relative \mathcal{I} -cell complexes in \mathcal{C} to relative \mathcal{J} -cell complexes in \mathcal{D} .*
- (ii) *U sends \mathcal{J} -injective morphisms in \mathcal{D} to \mathcal{I} -injective morphisms in \mathcal{C} .*
- (iii) *F sends \mathcal{I} -cofibrations in \mathcal{C} to \mathcal{J} -cofibrations in \mathcal{D} .*

Proof. (i). This is a corollary of the fact that F preserves all colimits.

(ii). As in the proof of [proposition A.3.19](#), a morphism $f : X \rightarrow Y$ in \mathcal{D} has the right lifting property with respect to all morphisms in \mathcal{J} if and only if $Uf : UX \rightarrow UY$ has the right lifting property with respect to all morphisms in \mathcal{I} .

(iii). Similarly, a morphism $g : Z \rightarrow W$ in \mathcal{C} has the left lifting property with respect to all morphisms of the form $Uf : UX \rightarrow UY$ where $f : X \rightarrow Y$ is a \mathcal{J} -injective morphism $f : X \rightarrow Y$ in \mathcal{D} if and only if $Fg : FZ \rightarrow FW$ is a \mathcal{J} -cofibration in \mathcal{D} ; but we know that U sends \mathcal{J} -injective morphisms in \mathcal{D} to \mathcal{I} -injective morphisms in \mathcal{C} , so F must send \mathcal{I} -cofibrations in \mathcal{C} to \mathcal{J} -cofibrations in \mathcal{D} . ■

Theorem 0.4.17 (Garner's small object argument). *Let \mathcal{C} be a locally presentable \mathbf{U} -category and let \mathcal{I} be any \mathbf{U} -subset of $\text{mor } \mathcal{C}$. There then exists an algebraic factorisation system (\mathbf{L}, \mathbf{R}) on \mathcal{C} such that the induced weak factorisation system is cofibrantly generated by \mathcal{I} .*

Proof. See Theorem 4.4 in [Garner, 2009]. □

Proposition 0.4.18. *Let \mathbf{U} be a universe, let \mathbf{Set} be the category of \mathbf{U} -sets, let \mathbb{B} be a \mathbf{U} -small category, let $\mathcal{C} = [\mathbb{B}^{\text{op}}, \mathbf{Set}]$, and let \mathcal{I} be the subset of $\text{mor } \mathcal{C}$ consisting of all monomorphisms $e : U \rightarrow V$ in \mathcal{C} where V is a quotient of a representable presheaf.*

- (i) *$(\text{cof}_{\mathcal{I}} \mathcal{C}, \text{inj}^{\mathcal{I}} \mathcal{C})$ is a \mathbf{U} -cofibrantly-generated weak factorisation system.*
- (ii) *$\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$ is precisely the class of all monomorphisms in \mathcal{C} .*
- (iii) *$\text{cof}_{\mathcal{I}} \mathcal{C} = \text{cell}_{\mathcal{I}} \mathcal{C}$.*

Proof. (i). Since \mathbb{B} is small and \mathcal{C} is well-powered and well-copowered, the full subcategory of $[2, \mathcal{C}]$ spanned by \mathcal{I} is essentially \mathbf{U} -small. We know that \mathcal{C} is locally finitely presentable, thus, taking a \mathbf{U} -set of representatives of the isomorphism classes in \mathcal{I} , and recalling [remark 0.4.9](#), Quillen's small object argument ([theorem 0.4.11](#)) implies $(\text{cof}_{\mathcal{I}} \mathcal{C}, \text{inj}^{\mathcal{I}} \mathcal{C})$ is indeed a \mathbf{U} -cofibrantly-generated weak factorisation system.

(ii). It is clear that the class of injective maps is closed under pushout and transfinite composition in \mathbf{Set} , so the same must be true of monomorphisms in \mathcal{C} , since colimits in \mathcal{C} are computed componentwise. Thus every morphism in $\text{cell}_{\mathcal{I}} \mathcal{C}$ is a monomorphism.

Conversely, suppose $f : X \rightarrow Y$ is a monomorphism. Fix an ordinal α and a bijection $y_{\bullet} : \alpha \rightarrow \coprod_{B \in \text{ob } \mathbb{B}} Y(B)$, and write B_{β} for the object in \mathbb{B} such that $y_{\beta} \in Y(B_{\beta})$. We will construct a \mathbf{U} -small presentation for f by transfinite recursion on α .

- To begin, put $X_0 = X$ and $f_0 = f$.
- For each ordinal $\beta < \alpha$, the Yoneda lemma implies there is a unique morphism $a_{\beta} : \mathcal{A}_{B_{\beta}} \rightarrow Y$ in \mathcal{C} such that $a_{\beta}(\text{id}_{B_{\beta}}) = y_{\beta}$; let $\bar{v}_{\beta} : V_{\beta} \rightarrow Y$ be the image of a_{β} , and let $e_{\beta} : U_{\beta} \rightarrow V_{\beta}$ and $u_{\beta} : U_{\beta} \rightarrow X_{\beta}$ be defined by the pullback square shown below:

$$\begin{array}{ccc} U_{\beta} & \xrightarrow{u_{\beta}} & X_{\beta} \\ e_{\beta} \downarrow & & \downarrow f_{\beta} \\ V_{\beta} & \xrightarrow{\bar{v}_{\beta}} & Y \end{array}$$

Since f_{β} is a monomorphism, e_{β} must also be a monomorphism and hence is in \mathcal{I} . There is then a commutative diagram in \mathcal{C} of the following form,

$$\begin{array}{ccccc} U_{\beta} & \xrightarrow{u_{\beta}} & X_{\beta} & & \\ e_{\beta} \downarrow & & \downarrow & \searrow f_{\beta} & \\ V_{\beta} & \xrightarrow{v_{\beta}} & X_{\beta+1} & & \\ & & \downarrow & \searrow f_{\beta+1} & \\ & & & & Y \\ & \searrow \bar{v}_{\beta} & & & \end{array}$$

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where $f_{\beta+1} : X_{\beta+1} \rightarrow Y$ is the union of $f_{\beta} : X_{\beta} \rightarrow Y$ and $\bar{v}_{\beta} : V_{\beta} \rightarrow Y$ considered as subobjects of Y ; note that the inner square of the diagram is then a pushout square.

- Finally, for limit ordinals $\gamma < \alpha$, we take $f_{\gamma} : X_{\gamma} \rightarrow Y$ to be the union $\bigcup_{\beta < \gamma} f_{\beta}$.

This completes the presentation of $f : X \rightarrow Y$ as a relative \mathcal{I} -cell complex in \mathcal{C} , and it is clearly \mathbf{U} -small.

(iii). [Corollary 0.4.12](#) implies that each morphism in $\text{cof}_{\mathcal{I}} \mathcal{C}$ is a retract of some morphism in $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$, but the class of monomorphisms is closed under retracts, so in this case we must have $\text{cof}_{\mathcal{I}} \mathcal{C} = \text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C}$. Since $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C} \subseteq \text{cell}_{\mathcal{I}} \mathcal{C} \subseteq \text{cof}_{\mathcal{I}} \mathcal{C}$, we also deduce that $\text{cell}_{\mathcal{I}, \mathbf{U}} \mathcal{C} = \text{cell}_{\mathcal{I}} \mathcal{C}$. ■

SIMPLICIAL SETS

Simplicial sets, like simplicial complexes, are combinatorial models for spaces built up by gluing standard n -simplices together; unlike simplicial complexes, an n -simplex in a simplicial set need not be uniquely determined by its vertices. It is for this reason that simplicial sets were once known by the unwieldy name ‘complete semi-simplicial (c.s.s.) complex’.

In the 1960s, it was discovered that one can mimic the definitions and constructions of classical homotopy theory by combinatorial means using simplicial sets, and that the resulting theory is moreover equivalent to the classical theory in a natural, functorial way. More recently, it has been shown that the homotopy theory of simplicial sets is *universal* in a precise sense,^[1] so it seems fitting that we begin here.

I.1 Basics

Definition I.1.1. The **simplex category** is the category $\mathbf{\Delta}$ whose objects are the positive finite ordinals and whose morphisms are the monotone maps. We use the geometer’s convention: $[n]$ denotes the ordinal $\{0, 1, \dots, n\}$.

Definition I.1.2. A **simplicial object** in a category \mathcal{C} is a functor $\mathbf{\Delta}^{\text{op}} \rightarrow \mathcal{C}$, and a **morphism of simplicial objects** in \mathcal{C} is a natural transformation of such functors. The **category of simplicial objects** in \mathcal{C} is the functor category $[\mathbf{\Delta}^{\text{op}}, \mathcal{C}]$ and is denoted by $\text{s}\mathcal{C}$.

[1] See [Dugger, 2001a].

Definition 1.1.3. The **coface maps** in $\mathbf{\Delta}$ are the morphisms $\delta_n^i : [n-1] \rightarrow [n]$, where δ_n^i is the unique injective monotone map that misses i ; and the **codegeneracy maps** in $\mathbf{\Delta}$ are the morphisms $\sigma_n^i : [n+1] \rightarrow [n]$, where σ_n^i is the unique surjective monotone map with $\sigma_n^i(i) = \sigma_n^i(i+1) = i$.

Theorem 1.1.4 (Cosimplicial identities). *The following equations hold in $\mathbf{\Delta}$:*

$$\begin{aligned}
 \delta_{n+1}^{j+1} \circ \delta_n^i &= \delta_{n+1}^i \circ \delta_n^j && \text{if } 0 \leq i \leq j \leq n \\
 \sigma_n^j \circ \sigma_{n+1}^i &= \sigma_n^i \circ \sigma_{n+1}^{j+1} && \text{if } 0 \leq i \leq j \leq n \\
 \sigma_{n+1}^{j+1} \circ \delta_{n+1}^i &= \delta_n^i \circ \sigma_n^j && \text{if } 0 \leq i \leq j \leq n \\
 \delta_n^{j+1} \circ \sigma_n^i &= \sigma_{n+1}^i \circ \delta_{n+1}^{j+2} && \text{if } 0 \leq i < j < n \\
 \sigma_n^i \circ \delta_n^i &= \text{id} && \text{if } 0 \leq i \leq n \\
 \sigma_n^{i+1} \circ \delta_n^i &= \text{id} && \text{if } 0 \leq i < n
 \end{aligned}$$

Equivalently, the following diagrams commute:

$$\begin{array}{ccc}
 [n-1] & \xrightarrow{\delta^i} & [n] \\
 \delta^j \downarrow & & \downarrow \delta^{j+1} \\
 [n] & \xrightarrow{\delta^i} & [n+1]
 \end{array} \quad \text{for } 0 \leq i \leq j \leq n$$

$$\begin{array}{ccc}
 [n+1] & \xrightarrow{\sigma^i} & [n] \\
 \sigma^{j+1} \downarrow & & \downarrow \sigma^j \\
 [n] & \xrightarrow{\sigma^i} & [n-1]
 \end{array} \quad \text{for } 0 \leq i \leq j \leq n$$

$$\begin{array}{ccc}
 [n] & \xrightarrow{\delta^i} & [n+1] \\
 \sigma^j \downarrow & & \downarrow \sigma^{j+1} \\
 [n-1] & \xrightarrow{\delta^i} & [n]
 \end{array} \quad \text{for } 0 \leq i \leq j \leq n$$

$$\begin{array}{ccc}
 [n] & \xrightarrow{\sigma^i} & [n-1] \\
 \delta^{j+2} \downarrow & & \downarrow \delta^{j+1} \\
 [n+1] & \xrightarrow{\sigma^i} & [n]
 \end{array} \quad \text{for } 0 \leq i < j < n$$

$$\begin{array}{ccc}
 [n-1] & \xrightarrow{\delta^i} & [n] \\
 \delta^i \downarrow & \searrow \text{id} & \downarrow \sigma^{i+1} \\
 [n] & \xrightarrow{\sigma^i} & [n-1]
 \end{array} \quad \text{for } 0 \leq i \leq n$$

Moreover, every morphism $[n] \rightarrow [m]$ in $\mathbf{\Delta}$ is uniquely a composite of the form

$$\delta_m^{j_1} \circ \dots \circ \delta_k^{j_{m-k}} \circ \sigma_k^{i_{n-k}} \circ \dots \circ \sigma_n^{i_1}$$

where $k \leq \min\{n, m\}$, and

$$\begin{aligned}
 0 &\leq i_{n-k} \leq \dots \leq i_1 \leq n \\
 0 &\leq j_{m-k} \leq \dots \leq j_1 \leq m
 \end{aligned}$$

The category $\mathbf{\Delta}$ is uniquely characterised by these properties.

Proof. See [May, 1967, § 2], [GZ, Ch. II, § 2], or [Weibel, 1994, § 8.1]. \square

Definition 1.1.5. Let A be a simplicial object in a category \mathcal{C} . A **face operator** for A is a morphism of the form $A(\delta_n^i) : A([n]) \rightarrow A([n-1])$, and a **degeneracy operator** for A is a morphism of the form $A(\sigma_n^i) : A([n]) \rightarrow A([n+1])$. For brevity, we will usually write A_n instead of $A([n])$, d_i^n instead of $A(\delta_n^i)$, and s_i^n instead of $A(\sigma_n^i)$.

Corollary 1.1.6 (Simplicial identities). *The face and degeneracy operators of a simplicial object satisfy the formal duals of the equations in theorem 1.1.4.* \blacksquare

Corollary 1.1.7. *A simplicial object A is uniquely determined by the sequence of objects A_0, A_1, A_2, \dots together with the face and degeneracy operators. Conversely, any sequence of objects equipped with face and degeneracy operators satisfying the simplicial identities defined a simplicial object.* \blacksquare

Definition 1.1.8. A **simplicial set** is a simplicial object in \mathbf{Set} , and the **category of simplicial sets** is denoted by \mathbf{sSet} .

Lemma 1.1.9.

- (i) *Limits (resp. colimits) in \mathbf{sSet} are constructed degreewise: a cone (resp. cocone) in \mathbf{sSet} over a diagram is limiting (resp. colimiting) if and only if it is so in every degree.*
- (ii) *A morphism of \mathbf{sSet} is monic (resp. epic) if and only if it is degreewise injective (resp. surjective).*

Proof. These are standard facts about functor categories. \square

Definition 1.1.10. The **standard n -simplex** in \mathbf{sSet} , denoted by Δ^n , is the representable presheaf $\mathbf{\Delta}(-, [n])$.

Theorem 1.1.11. Let $\Delta^\bullet : \mathbf{\Delta} \rightarrow \mathbf{sSet}$ be the functor $[n] \mapsto \Delta^n$.

- (i) For any simplicial set X , the map $\mathbf{sSet}(\Delta^n, X) \rightarrow X_n$ defined by $f \mapsto f_n(\text{id}_{[n]})$ is a bijection and is moreover natural in $[n]$ and X .
- (ii) \mathbf{sSet} has limits and colimits for all small diagrams, every epimorphism is effective, and for all morphisms $f : X \rightarrow Y$ in \mathbf{sSet} , the pullback functor $f^* : \mathbf{sSet}_{/Y} \rightarrow \mathbf{sSet}_{/X}$ preserves colimits.
- (iii) $\Delta^\bullet : \mathbf{\Delta} \rightarrow \mathbf{sSet}$ is a dense functor, i.e. for any simplicial set X , the tautological cocone^[2] from the canonical diagram $(\Delta^\bullet \downarrow X) \rightarrow \mathbf{sSet}$ to X is colimiting.
- (iv) Let \mathcal{E} be a locally small category with colimits for all small diagrams. If $F : \mathbf{sSet} \rightarrow \mathcal{E}$ is a functor that preserves small colimits, then it is left adjoint to the functor $\mathcal{E} \rightarrow \mathbf{sSet}$ defined by $E \mapsto \mathcal{E}(F\Delta^\bullet, E)$.
- (v) With \mathcal{E} as above, the functor $F \mapsto F\Delta^\bullet$ from the category of colimit-preserving functors $\mathbf{sSet} \rightarrow \mathcal{E}$ to the category of all functors $\mathbf{\Delta} \rightarrow \mathcal{E}$ is fully faithful and essentially surjective on objects.

Proof. Claim (i) is just the Yoneda lemma, claim (ii) follows from the lemma above, and claims (iii)–(v) are just facts about dense functors, pointwise left Kan extensions, weighted colimits: see [proposition A.5.24](#), [theorem A.5.15](#), and [proposition A.6.11](#). \blacksquare

Definition 1.1.12. An element of X_n is called an **n -simplex** of X ; in particular, an element of X_0 is a **vertex** of X and an element of X_1 is an **edge** of X . This is justified by statement (i) in the above theorem. Given an edge f of X , the **source** of f is the vertex $d_1(f)$, and the **target** of f is the vertex $d_0(f)$; we write $f : x \rightarrow y$ to mean $d_1(f) = x$ and $d_2(f) = y$.

[2] See [definition A.5.7](#).

Definition I.I.13. The **standard n -simplex** in **Top**, denoted by $|\Delta^n|$, is the topological space

$$|\Delta^n| = \{(x_0, \dots, x_n) \in [0, 1]^{n+1} \mid x_0 + \dots + x_n = 1\}$$

where $[0, 1]$ is the closed unit interval with the standard metric. The functor $|\Delta^\bullet| : \mathbf{\Delta} \rightarrow \mathbf{Top}$ sends $[n]$ to $|\Delta^n|$ and is defined on morphisms by linearly interpolating the obvious map of vertices.

Corollary I.I.14. *There exists an adjunction*

$$|-| \dashv S : \mathbf{Top} \rightarrow \mathbf{sSet}$$

extending the functor $|\Delta^\bullet| : \mathbf{\Delta} \rightarrow \mathbf{Top}$ defined above, and this adjunction is unique up to unique isomorphism. Explicitly, we may take

$$S(Y)_n = \mathbf{Top}(|\Delta^n|, Y)$$

with the evident face and degeneracy operators induced by the coface and codegeneracy maps in $\mathbf{\Delta}$. ■

Definition I.I.15. The **geometric realisation** of a simplicial set X is the topological space $|X|$, and the **singular set** of a topological space Y is the simplicial set $S(Y)$.

REMARK I.I.16. The geometric realisation $|X|$ is stable under universe enlargement, by [theorem 0.3.11](#).

Theorem I.I.17. *Let **CGHaus** be the category of compactly-generated Hausdorff spaces^[3] and continuous maps.*

- (i) *The topological standard n -simplex $|\Delta^n|$ is a compact Hausdorff space.*
- (ii) *For any simplicial set X , the geometric realisation $|X|$ is a compactly-generated Hausdorff space.*
- (iii) *The previously-constructed adjunction $|-| \dashv S : \mathbf{Top} \rightarrow \mathbf{sSet}$ restricts to an adjunction between **CGHaus** and **sSet**, and moreover the functor $|-| : \mathbf{sSet} \rightarrow \mathbf{CGHaus}$ preserves finite limits and reflects isomorphisms.*

Proof. Claim (i) is a standard fact, while claims (ii) and (iii) are proven in [\[GZ, Ch. III, § 3\]](#). □

[3] See [definition A.2.26](#).

1.2 Nerves, skeletons, and coskeletons

Prerequisites. §§ 1.1, A.2.

Proposition 1.2.1. *Let $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$ be the functor defined by the formula*

$$N(\mathbf{C})_n = \text{Fun}([n], \mathbf{C})$$

where $[n]$ here denotes the preorder category $\{0 \rightarrow \dots \rightarrow n\}$.

- (i) $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$ has a left adjoint $\tau_1 : \mathbf{sSet} \rightarrow \mathbf{Cat}$ such that $\tau_1 \Delta^n = [n]$.
- (ii) The functor N is fully faithful and exhibits \mathbf{Cat} as a reflective subcategory of \mathbf{sSet} .
- (iii) $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$ is a cartesian closed functor.
- (iv) The functor τ_1 preserves finite products.

Proof. (i). Apply [theorem 1.1.11](#).

(ii). A functor is entirely determined by its action on objects, arrows, and composable strings of arrows, so N is fully faithful.

(iii). N preserves binary products, so we have the following natural bijections:

$$\begin{aligned} \mathbf{sSet}(\Delta^n, N([\mathbf{C}, \mathbf{D}])) &\cong \text{Fun}([n], [\mathbf{C}, \mathbf{D}]) \\ &\cong \text{Fun}([n] \times \mathbf{C}, \mathbf{D}) \\ &\cong \mathbf{sSet}(N([n] \times \mathbf{C}), N(\mathbf{D})) \\ &\cong \mathbf{sSet}(N([n]) \times N(\mathbf{C}), N(\mathbf{D})) \\ &\cong \mathbf{sSet}(N([n]), [N(\mathbf{C}), N(\mathbf{D})]) \\ &\cong \mathbf{sSet}(\Delta^n, [N(\mathbf{C}), N(\mathbf{D})]) \end{aligned}$$

Thus, by the Yoneda lemma, the canonical morphism $N([\mathbf{C}, \mathbf{D}]) \rightarrow [N(\mathbf{C}), N(\mathbf{D})]$ is an isomorphism.

(iv). It is clear that τ_1 preserves terminal objects. Let X and Y be simplicial sets. We wish to show that the canonical morphism $\tau_1(X \times Y) \rightarrow \tau_1 X \times \tau_1 Y$ is an isomorphism; but since τ_1 is a left adjoint and both \mathbf{sSet} and \mathbf{Cat} are cartesian

closed, it is enough to check the claim for $Y = \Delta^n$, because \mathbf{sSet} is generated under colimits by $\{\Delta^n \mid n \in \mathbb{N}\}$. We have the following natural bijections:

$$\begin{aligned} \mathrm{Fun}(\tau_1(X \times \Delta^n), \mathbb{C}) &\cong \mathbf{sSet}(X \times \Delta^n, \mathbf{N}(\mathbb{C})) \\ &\cong \mathbf{sSet}(X, \mathbf{N}(\mathbb{C})^{\Delta^n}) \\ &\cong \mathbf{sSet}(X, \mathbf{N}([n], \mathbb{C})) \\ &\cong \mathrm{Fun}(\tau_1 X, [[n], \mathbb{C}]) \\ &\cong \mathrm{Fun}(\tau_1 X \times [n], \mathbb{C}) \\ &\cong \mathrm{Fun}(\tau_1 X \times \tau_1 \Delta^n, \mathbb{C}) \end{aligned}$$

The claim follows by the Yoneda lemma. ■

Definition 1.2.2. The **fundamental category** of a simplicial set X is the small category $\tau_1 X$, and the **nerve** of a small category \mathbb{C} is the simplicial set $\mathbf{N}(\mathbb{C})$.

REMARK 1.2.3. Given a simplicial set X , the fundamental category $\tau_1 X$ admits the following presentation by generators and relations: the objects are the vertices of X , and the arrows are generated by the edges of X , modulo the relation $d_0(\alpha) \circ d_2(\alpha) = d_1(\alpha)$ for all 2-simplices α in X . This shows that $\tau_1 X$ is stable under universe enlargement.

Proposition 1.2.4. Let $\mathrm{disc} : \mathbf{Set} \rightarrow \mathbf{sSet}$ be the functor defined by the formula

$$(\mathrm{disc} Y)_n = Y$$

with id_Y for all the face and degeneracy maps.

- (i) $\mathrm{disc} : \mathbf{Set} \rightarrow \mathbf{sSet}$ has a left adjoint $\pi_0 : \mathbf{sSet} \rightarrow \mathbf{Set}$ such that $\pi_0 \Delta^n = 1$.
- (ii) The functor disc is fully faithful and exhibits \mathbf{Set} as a reflective subcategory of \mathbf{sSet} .
- (iii) $\mathbf{N} : \mathbf{Set} \rightarrow \mathbf{sSet}$ is a cartesian closed functor.
- (iv) The functor π_0 preserves products.

Proof. (i). We could apply [theorem 1.1.11](#), but it is also fairly straightforward to check that this explicit construction works: for each simplicial set X , we define $\pi_0 X$ by the coequaliser diagram in \mathbf{Set} shown below,

$$X_1 \begin{array}{c} \xrightarrow{d_0} \\ \xrightarrow{d_1} \end{array} X_0 \longrightarrow \pi_0 X$$

and for each morphism $f : X \rightarrow Y$ in \mathbf{sSet} , we define $\pi_0 f$ to be the unique morphism making the evident diagram commute.

(ii). It is clear that disc is fully faithful.

(iii). By [proposition A.2.15](#), we have an analogous adjunction $\pi_0 \dashv \text{disc} : \mathbf{Set} \rightarrow \mathbf{Cat}$. It is clear that we have a natural isomorphism $N(\text{disc } Y) \cong \text{disc } Y$ for every set Y , and we know $\text{disc} : \mathbf{Set} \rightarrow \mathbf{Cat}$ and $N : \mathbf{Cat} \rightarrow \mathbf{sSet}$ are cartesian closed functors, so $\text{disc} : \mathbf{Set} \rightarrow \mathbf{sSet}$ must also be cartesian closed.

(iv). Similarly, for any simplicial set X , we have a natural isomorphism $\pi_0 X \cong \pi_0 \tau_1 X$; but we know that $\pi_0 : \mathbf{Cat} \rightarrow \mathbf{Set}$ preserves finite products, and $\tau_1 : \mathbf{sSet} \rightarrow \mathbf{Cat}$ preserves finite products by [proposition 1.2.1](#), so $\pi_0 : \mathbf{sSet} \rightarrow \mathbf{Set}$ must also preserve finite products. ■

Definition 1.2.5. The **set of connected components** of a simplicial set X is the set $\pi_0 X$, and a **discrete simplicial set** is one that is isomorphic to $\text{disc } Y$ for some set Y .

¶ 1.2.6. We will usually not distinguish between Y and $\text{disc } Y$ notationally.

Proposition 1.2.7. Let $N : \mathbf{Grpd} \rightarrow \mathbf{sSet}$ be the functor defined by the formula

$$N(\mathbb{G})_n = \text{Fun}(\mathbf{I}[n], \mathbb{G})$$

where $\mathbf{I}[n]$ here denotes the groupoid obtained by freely inverting the arrows in the preorder category $[n]$.

- (i) For any groupoid \mathbb{G} , the nerve $N(\mathbb{G})$ is the same (up to isomorphism) whether computed for \mathbb{G} as a groupoid or \mathbb{G} as a category.
- (ii) $N : \mathbf{Grpd} \rightarrow \mathbf{sSet}$ has a left adjoint $\pi_1 : \mathbf{sSet} \rightarrow \mathbf{Grpd}$ such that $\pi_1 \Delta^n = \mathbf{I}[n]$.
- (iii) The functor N is fully faithful and exhibits \mathbf{Grpd} as a reflective subcategory of \mathbf{sSet} .
- (iv) $N : \mathbf{Grpd} \rightarrow \mathbf{sSet}$ is a cartesian closed functor.
- (v) The functor π_1 preserves finite products.

Proof. (i). By the universal property of $\mathbf{I}[n]$, there is a natural bijection

$$\mathrm{Fun}(\mathbf{I}[n], \mathbb{G}) \cong \mathrm{Fun}([n], \mathbb{G})$$

for all groupoids \mathbb{G} , so the two nerve constructions do indeed agree.

(ii) and (iii). These are proven in exactly the same way as in [proposition 1.2.1](#).

(iv) and (v). These are proven in exactly the same way as in [proposition 1.2.4](#). ■

Definition 1.2.8. The **fundamental groupoid** of a simplicial set X is the small groupoid $\pi_1 X$.

REMARK 1.2.9. Given a simplicial set X , the fundamental groupoid $\pi_1 X$ admits a presentation of the same kind as the fundamental category $\tau_1 X$, and in fact $\pi_1 X$ is isomorphic to the groupoid obtained by freely inverting the arrows in $\tau_1 X$:

$$\mathrm{Fun}(\pi_1 X, \mathbb{G}) \cong \mathbf{sSet}(X, \mathbf{N}(\mathbb{G})) \cong \mathrm{Fun}(\tau_1 X, \mathbb{G})$$

This shows that $\pi_1 X$ is stable under universe enlargement.

Definition 1.2.10. Let n be a natural number, and let $\mathbf{\Delta}_{\leq n}$ be the full subcategory of $\mathbf{\Delta}$ spanned by the objects $[0], \dots, [n]$. An n -**truncated simplicial set** is a functor $\mathbf{\Delta}_{\leq n}^{\mathrm{op}} \rightarrow \mathbf{Set}$, and we write $\mathbf{sSet}_{\leq n}$ for the category of n -truncated simplicial sets. The **brutal n -truncation** of a simplicial set X is the n -truncated simplicial set $X_{\leq n}$ defined by the evident reduct:

$$X_{\leq n}([m]) = X([m])$$

Proposition 1.2.11. Let n be a natural number, and let $j : \mathbf{\Delta}_{\leq n} \rightarrow \mathbf{\Delta}$ be the inclusion.

- (i) The functor $j^* : \mathbf{sSet} \rightarrow \mathbf{sSet}_{\leq n}$ has a left adjoint $\mathrm{Lan}_j : \mathbf{sSet}_{\leq n} \rightarrow \mathbf{sSet}$.
- (ii) The unit $\mathrm{id} \Rightarrow j^* \mathrm{Lan}_j$ is a natural isomorphism.
- (iii) $\mathrm{Lan}_j : \mathbf{sSet}_{\leq n} \rightarrow \mathbf{sSet}$ is a fully faithful functor.
- (i') The functor $j^* : \mathbf{sSet} \rightarrow \mathbf{sSet}_{\leq n}$ has a right adjoint $\mathrm{Ran}_j : \mathbf{sSet}_{\leq n} \rightarrow \mathbf{sSet}$.
- (ii') The counit $j^* \mathrm{Ran}_j \Rightarrow \mathrm{id}$ is a natural isomorphism.

(iii') $\text{Ran}_j : \mathbf{sSet}_{\leq n} \rightarrow \mathbf{sSet}$ is a fully faithful functor.

Proof. (i) and (i'). Use [theorem A.5.15](#).

(ii) and (ii'). The inclusion $j : \Delta_{\leq n} \rightarrow \Delta$ is fully faithful, so the unit $\text{id} \Rightarrow j^* \text{Lan}_j$ and the counit $j^* \text{Ran}_j \Rightarrow \text{id}$ are natural isomorphisms, by [corollary A.5.19](#).

(iii) and (iii'). It is a well-known fact that the unit (resp. counit) of an adjunction is a natural isomorphism if and only if the left (resp. right) adjoint is fully faithful.^[4] ■

Definition 1.2.12. For each natural number n , with notation as above, let $\text{sk}_n : \mathbf{sSet} \rightarrow \mathbf{sSet}$ be the composite $\text{Lan}_j j^*$, and let $\text{cosk}_n : \mathbf{sSet} \rightarrow \mathbf{sSet}$ be the composite $\text{Ran}_j j^*$. The n -**skeleton** of a simplicial set X is the simplicial set $\text{sk}_n(X)$, and the n -**coskeleton** of a simplicial set is the simplicial set $\text{cosk}_n(X)$. A n -**skeletal simplicial set** is one that is isomorphic to the n -skeleton of some simplicial set, and an n -**coskeletal simplicial set** is one that is isomorphic to the n -coskeleton of some simplicial set.

REMARK 1.2.13. In the special case $n = 0$, Lan_j may be identified with the functor $\text{disc} : \mathbf{Set} \rightarrow \mathbf{sSet}$ defined in [proposition 1.2.4](#). Thus, 0-skeletal simplicial sets are precisely the discrete simplicial sets. On the other hand, given a set X , $\text{Ran}_j X$ can be identified with the simplicial set whose m -simplices are $(m + 1)$ -tuples of elements of X , with face and degeneracy maps induced by the appropriate projections.

Proposition 1.2.14. *Let n be a natural number.*

- (i) *The full subcategory of n -skeletal simplicial sets is a coreflective subcategory of \mathbf{sSet} , with coreflector sk_n .*
- (ii) *sk_n is the underlying endofunctor of an idempotent comonad on \mathbf{sSet} .*
- (iii) *A simplicial set X is n -skeletal if and only if the counit $\text{sk}_n(X) \rightarrow X$ is an isomorphism.*
- (iv) *If $m \geq n$, then any n -skeletal simplicial set is also m -skeletal.*
- (i') *The full subcategory of n -coskeletal simplicial sets is a reflective subcategory of \mathbf{sSet} , with reflector cosk_n .*

[4] See e.g. [\[CWM, Ch. IV, § 3\]](#).

- (ii') cosk_n is the underlying endofunctor of an idempotent monad on \mathbf{sSet} .
- (iii') A simplicial set X is n -coskeletal if and only if the unit $X \rightarrow \text{cosk}_n(X)$ is an isomorphism.
- (iv') If $m \geq n$, then any n -coskeletal simplicial set is also m -coskeletal.

Proof. All straightforward from the definitions. ◆

Proposition 1.2.15. *Let n be a natural number, and let X be a simplicial set.*

- (i) *We have the following adjunction:*

$$\text{sk}_n \dashv \text{cosk}_n : \mathbf{sSet} \rightarrow \mathbf{sSet}$$

- (ii) *The counit $\text{sk}_n(X) \rightarrow X$ is a monomorphism, and X is n -skeletal if and only if all m -simplices of X are degenerate for $m > n$.*
- (iii) *X is n -coskeletal if and only if, for all natural numbers m , the map*

$$X_m \cong \mathbf{sSet}(\Delta^m, X) \rightarrow \mathbf{sSet}(\text{sk}_n(\Delta^m), X)$$

induced by the counit $\text{sk}_n(\Delta^m) \rightarrow \Delta^m$ is a bijection.

Proof. (i). Immediate from the definition of sk_n and cosk_n .

(ii). The most straightforward way of seeing this is to construct $\text{sk}_n(X)$ explicitly as the smallest simplicial subset of X containing all of its n -simplices.

(iii). Apply the Yoneda lemma in conjunction with claim (i). ■

Example 1.2.16. For any small category \mathbb{C} , the nerve $N(\mathbb{C})$ is a 2-coskeletal simplicial set: by definition, an m -simplex of $N(\mathbb{C})$ is just a functor $[m] \rightarrow \mathbb{C}$, but the property of being a functor can be detected by only inspecting the vertices, edges, and 2-cells.

Proposition 1.2.17. *The following full subcategories are exponential ideals of \mathbf{sSet} :*

- (i) *Discrete simplicial sets.*
- (ii) *Simplicial sets isomorphic to the nerve of some category.*

(iii) *Simplicial sets isomorphic to the nerve of some groupoid.*

(iv) *n -coskeletal simplicial sets for some natural number n .*

Proof. Apply [proposition A.2.13](#) to [propositions 1.2.4, 1.2.1, 1.2.7, and 1.2.14](#). ■

1.3 The Kan–Quillen model structure

Prerequisites. §§ [0.4](#) [1.1](#), [A.3](#).

In [1967], Quillen constructed an axiomatic framework for doing homotopy theory in abstract categories, which he called ‘closed model categories’, and showed that \mathbf{sSet} can be endowed with a model structure such that the resulting homotopy theory is equivalent in a strong sense to the homotopy theory of topological spaces.

Definition 1.3.1. A **horn** is a simplicial subset of the form $\Lambda_k^n \subseteq \Delta^n$, where Λ_k^n is the union of the images of $\delta_n^0, \dots, \delta_n^{k-1}, \delta_n^{k+1}, \dots, \delta_n^n : \Delta^{n-1} \rightarrow \Delta^n$ in \mathbf{sSet} . In other words, Λ_k^n is the union of all the faces of Δ^n that include the k -th vertex. The **boundary** of Δ^n is the simplicial subset $\partial\Delta^n \subseteq \Delta^n$ generated by the images of $\delta_n^0, \dots, \delta_n^n : \Delta^{n-1} \rightarrow \Delta^n$.

REMARK 1.3.2. The boundary $\partial\Delta^n$ may be identified with $\mathrm{sk}_{n-1}\Delta^n$.

Definition 1.3.3. A **cofibration** in \mathbf{sSet} is a monomorphism. A **Kan fibration** is a morphism $f : X \rightarrow Y$ in \mathbf{sSet} that has the right lifting property with respect to the horn inclusions $\Lambda_k^n \hookrightarrow \Delta^n$, where $n \geq 1$ and $0 \leq k \leq n$. A **Kan complex** is a simplicial set X such that the unique morphism $X \rightarrow 1$ is a Kan fibration.

REMARK 1.3.4. In other words, a Kan complex is a simplicial set X satisfying the **Kan condition**: every horn $\alpha' : \Lambda_k^n \rightarrow X$ has a **filler**, i.e. a morphism $\alpha : \Delta^n \rightarrow X$ (equivalently, an n -simplex of X) such that α' is the restriction along the inclusion $\Lambda_k^n \hookrightarrow \Delta^n$.

Lemma 1.3.5. *If X is a Kan complex, then the fundamental category $\tau_1 X$ is a groupoid, and the unit $\eta_X : X \rightarrow \mathbf{N}(\tau_1 X)$ is an epimorphism.*

Proof. Let x, y , and z be vertices in X , and let $f : x \rightarrow y$ and $g : y \rightarrow z$ be edges in X .^[5] Then the pair (f, g) defines a horn $\Lambda_1^2 \rightarrow X$, and so by the Kan

[5] Recall [definition 1.1.12](#).

condition, there exists a 2-simplex α of X such that $d_2(\alpha) = f$ and $d_0(\alpha) = g$. By remark [remark 1.2.3](#), the composite $g \circ f$ defined in $\tau_1 X$ must correspond to the edge $d_1(\alpha)$. Since the arrows in $\tau_1 X$ are generated by the edges of X , we conclude by induction that $\eta_X : X \rightarrow N(\tau_1 X)$ is a surjection on vertices and edges.

Similarly, given an edge $f : x \rightarrow y$, the Kan condition ensures that there exist two 2-simplices β and γ such that

$$\begin{array}{ll} d_2(\alpha) = f & d_1(\alpha) = \text{id}_x \\ d_0(\alpha) = f & d_1(\alpha) = \text{id}_y \end{array}$$

where $\text{id}_x : x \rightarrow x$ is the edge $s_0(x)$, and $\text{id}_y : y \rightarrow y$ is the edge $s_0(y)$. Together with the argument in the previous paragraph, this shows that $\tau_1 X$ is a groupoid.

Finally, to show that $\eta_X : X \rightarrow N(\tau_1 X)$ is a surjection on n -simplices for $n \geq 2$, we simply observe that an n -simplex of $N(\tau_1 X)$ is just a string of n composable edges of X , so we may appeal to the Kan condition again to obtain the corresponding n -simplex of X . ■

Corollary 1.3.6. *If X is a Kan complex, then the unit $\eta_X : X \rightarrow N(\pi_1 X)$ is an epimorphism.*

Proof. Since $\tau_1 X$ is already a groupoid, the canonical functor $\tau_1 X \rightarrow \pi_1 X$ must be an isomorphism. (See [remark 1.2.9](#).) ■

Proposition 1.3.7. *Let \mathcal{I} and \mathcal{J} be the following subsets of mor sSet :*

$$\begin{aligned} \mathcal{I} &= \{ \Lambda_k^n \hookrightarrow \Delta^n \mid n \geq 1, 0 \leq k \leq n \} \\ \mathcal{J} &= \{ \partial \Delta^n \hookrightarrow \Delta^n \mid n \geq 0 \} \end{aligned}$$

- (i) *There exist a pair of functorial factorisation systems on sSet , one inducing a weak factorisation system cofibrantly generated by \mathcal{I} , and the other inducing a weak factorisation system cofibrantly generated by \mathcal{J} .*
- (ii) *A morphism is \mathcal{I} -injective if and only if it is a Kan fibration, and every \mathcal{I} -cofibration is a monomorphism (but not vice versa).*
- (iii) *A morphism is a \mathcal{J} -cofibration if and only if it is a monomorphism, and every \mathcal{J} -injective morphism is a Kan fibration (but not vice versa).*

Proof. (i). Since \mathbf{sSet} is a locally finitely presentable category, we may apply Quillen’s small object argument (theorem 0.4.11).

(ii). The definition of ‘Kan fibration’ is exactly the definition of ‘ \mathcal{I} -fibration’; on the other hand, the class of monomorphisms is closed under pushout, transfinite composition, and retracts in \mathbf{Set} , so the same is true for \mathbf{sSet} , and thus, by corollary 0.4.12, every \mathcal{I} -cofibration must be a monomorphism.

(iii). To prove that $\text{inj}^{\mathcal{I}} \mathcal{C} \supseteq \text{inj}^{\mathcal{J}} \mathcal{C}$, it is enough to check that $\mathcal{I} \subseteq \text{cof}_{\mathcal{J}} \mathcal{C}$; since every morphism in \mathcal{I} is a monomorphism, it will suffice to show that $\text{cof}_{\mathcal{J}} \mathcal{C}$ is precisely the class of all monomorphisms. For this, see the remarks at the beginning of [Joyal and Tierney, 2008, § 3.1], or Proposition 1 in [Quillen, 1967, Ch. II, § 2]. \square

Definition 1.3.8. An **anodyne extension**, or **trivial cofibration** in \mathbf{sSet} , is a cofibration that has the left lifting property with respect to all Kan fibrations. A **trivial Kan fibration** is a Kan fibration that has the right lifting property with respect to all cofibrations.

Proposition 1.3.9. *Let $i : Z \rightarrow W$ be a cofibration in \mathbf{sSet} and let $p : X \rightarrow Y$ be a Kan fibration. Suppose we have a commutative diagram*

$$\begin{array}{ccc}
 [W, X] & \xrightarrow{[i, X]} & [Z, X] \\
 \downarrow [W, p] & \dashrightarrow & \downarrow [Z, p] \\
 [W, Y] & \xrightarrow{[i, Y]} & [Z, Y]
 \end{array}$$

$L(i, p)$ is positioned between $[Z, X]$ and $[W, Y]$, with a solid arrow from $L(i, p)$ to $[Z, X]$ and a solid arrow from $L(i, p)$ to $[W, Y]$.

where the square in the lower right is a pullback square.

- (i) The unique morphism $[W, X] \rightarrow L(i, p)$ making the diagram commute is a Kan fibration.
- (ii) If $i : Z \rightarrow W$ is an anodyne extension, then $[W, X] \rightarrow L(i, p)$ is a trivial Kan fibration.
- (iii) If $p : X \rightarrow Y$ is a trivial Kan fibration, then so is $[W, X] \rightarrow L(i, p)$.

Proof. (i). See Theorem 3.3.1 in [Hovey, 1999], or Proposition 5.2 in [GJ, Ch. I].

(ii) and (iii). See Proposition 11.5 in [GJ, Ch. I]; for a purely combinatorial proof, see Theorem 3.2.1 in [Joyal and Tierney, 2008]. \square

Corollary 1.3.10.

- (i) *If $p : X \rightarrow Y$ is a Kan fibration (resp. trivial Kan fibration), then for all simplicial sets W , the morphism $[W, p] : [W, X] \rightarrow [W, Y]$ is also a Kan fibration (resp. trivial Kan fibration).*
- (ii) *If $i : Z \rightarrow W$ is a cofibration (resp. anodyne extension) and X is a Kan complex, then the morphism $[i, X] : [W, X] \rightarrow [Z, X]$ is a Kan fibration (resp. trivial Kan fibration).*
- (iii) *If W is any simplicial set and X is a Kan complex, then $[W, X]$ is also a Kan complex.*

Proof. (i). Take $Z = \emptyset$; noting that the canonical morphism $\emptyset \rightarrow W$ is a cofibration, and that $[\emptyset, p] : [\emptyset, X] \rightarrow [\emptyset, Y]$ is an isomorphism, the proposition above then implies $[W, p] : [W, X] \rightarrow [W, Y]$ is a Kan fibration (resp. trivial Kan fibration).

(ii). Take $Y = 1$; since $[W, 1] \rightarrow [Z, 1]$ is an isomorphism, the proposition above implies $[i, X] : [W, X] \rightarrow [Z, X]$ is a Kan fibration (resp. trivial Kan fibration).

(iii). Noting that $[\emptyset, X]$ is a terminal object in \mathbf{sSet} , we apply claim (ii) to the case $Z = \emptyset$ to obtain the desired conclusion. \blacksquare

The following combinatorial definition of weak homotopy equivalence is due to Joyal and Tierney [2008]. Recalling the definition of $\pi_0 : \mathbf{sSet} \rightarrow \mathbf{Set}$ from proposition 1.2.4 as the functor sending a simplicial set X to the set π_0 of its connected components,

Definition 1.3.11. A **weak homotopy equivalence of simplicial sets** is a morphism $f : W \rightarrow Z$ such that, for every Kan complex K , the induced map

$$\pi_0[f, K] : \pi_0[Z, K] \rightarrow \pi_0[W, K]$$

is a bijection of sets.

Proposition 1.3.12.

- (i) A Kan fibration $p : X \rightarrow Y$ is trivial if and only if it is a weak homotopy equivalence.
- (ii) A cofibration $i : Z \rightarrow W$ is an anodyne extension if and only if it is a weak homotopy equivalence.

Proof. See Propositions 3.4.1 and 3.4.2 in [Joyal and Tierney, 2008]. □

In summary, we have:

Theorem 1.3.13. *sSet*, regarded as a *sSet*-enriched category via its cartesian closed structure, is a simplicial model category where

- the cofibrations are the monomorphisms in *sSet*,
- the fibrations are the Kan fibrations, and
- the weak equivalences are the weak homotopy equivalences.

This is the **Kan–Quillen model structure on simplicial sets**.

Proof. We know *sSet* has limits and colimits for all small diagrams and is a cartesian closed category, so it satisfies axioms CM1 and SM0. Using the definition of weak homotopy equivalence given above, the class of weak homotopy equivalences has the 2-out-of-6 property by lemma A.4.13, hence axiom CM2 is satisfied. Proposition 1.3.7 plus theorem 3.1.8 then shows that the announced cofibrations, fibrations, and weak equivalences do indeed constitute a closed model structure on *sSet*.

Finally, we note that proposition 1.3.9 is precisely the condition required by axiom SM7. ■

Proposition 1.3.14. *There exist a functor $R : \mathbf{sSet} \rightarrow \mathbf{sSet}$ and a natural transformation $\eta : \mathrm{id}_{\mathbf{sSet}} \Rightarrow R$ such that, for all simplicial sets X , RX is a Kan complex and $i_X : X \rightarrow RX$ is an anodyne extension. Moreover, any such functor R preserves weak homotopy equivalences.*

Proof. By proposition 1.3.7, for each X , there is a factorisation of the unique morphism $X \rightarrow 1$ as an anodyne extension $i_X : X \rightarrow RX$ followed by a Kan fibration $RX \rightarrow 1$, and this is moreover functorial in X . Finally, if $f : X \rightarrow Y$

is a weak homotopy equivalence in \mathbf{sSet} , then the commutativity of the diagram below

$$\begin{array}{ccc} X & \xrightarrow{i_X} & RX \\ f \downarrow & & \downarrow Rf \\ Y & \xrightarrow{i_Y} & RY \end{array}$$

plus [proposition 1.3.12](#) and the 2-out-of-3 property for weak homotopy equivalences implies Rf is also a weak homotopy equivalence. ■

I.4 Intrinsic homotopy

Prerequisites. §§ 1.3, A.4.

Definition 1.4.1. Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in \mathbf{sSet} . An **intrinsic homotopy** $\alpha : f_0 \Rightarrow f_1$ is an edge of the exponential object $[X, Y]$ such that $d_1(\alpha) = f_0$ and $d_0(\alpha) = f_1$. (Note the subscripts!) We say f_0 and f_1 are **intrinsically homotopic** if there is a zigzag of intrinsic homotopies connecting f_0 and f_1 , or equivalently, if f_0 and f_1 are in the same connected component of $[X, Y]$.

REMARK 1.4.2. By the Yoneda lemma,

$$[X, Y]_1 \cong \mathbf{sSet}(\Delta^1, [X, Y]) \cong \mathbf{sSet}(\Delta^1 \times X, Y)$$

so an intrinsic homotopy $\alpha : f_0 \Rightarrow f_1$ is essentially the same thing as a morphism $\tilde{\alpha} : \Delta^1 \times X \rightarrow Y$ such that $\tilde{\alpha} \circ (\delta^1 \times \text{id}_X) = f_0$ and $\tilde{\alpha} \circ (\delta^0 \times \text{id}_X) = f_1$ (where we have suppressed the canonical isomorphism $X \cong \Delta^0 \times X$), just as in classical homotopy theory. Also,

$$\mathbf{sSet}(\Delta^1 \times X, Y) \cong \mathbf{sSet}(X, [\Delta^1, Y])$$

so intrinsic homotopies $\alpha : f_0 \Rightarrow f_1$ correspond to morphisms $\hat{\alpha} : X \rightarrow [\Delta^1, Y]$ such that $[\delta^1, Y] \circ \hat{\alpha} = f_0$ and $[\delta^0, Y] \circ \hat{\alpha} = f_1$ (where we have suppressed the canonical isomorphism $[\Delta^0, Y] \cong Y$).

The notion of intrinsic homotopy is not well-behaved for general simplicial sets Y . For example, the existence of an intrinsic homotopy $f_0 \Rightarrow f_1$ does not guarantee the existence of an “inverse” intrinsic homotopy $f_1 \Rightarrow f_0$, and even if we have intrinsic homotopies $f_0 \Rightarrow f_1$ and $f_1 \Rightarrow f_2$, there need not be an intrinsic homotopy $f_0 \Rightarrow f_2$. However:

Proposition 1.4.3. *For any simplicial set X and any Kan complex Y , the relation \sim on $\mathbf{sSet}(X, Y)$ defined by*

$$f_0 \sim f_1 \text{ if and only if there exists an intrinsic homotopy } f_0 \Rightarrow f_1$$

is an equivalence relation.

Proof. The relation \sim is certainly reflexive whether or not Y is a Kan complex. Recalling lemma 1.3.5, the transitivity of \sim may be deduced from the fact that the unit $\eta_X : X \rightarrow N(\tau_1 X)$ is an epimorphism, and the symmetry of \sim corresponds to the fact that $\tau_1 X$ is a groupoid. ■

¶ 1.4.4. Let \mathbf{Kan} be the full subcategory of \mathbf{sSet} spanned by the Kan complexes. For each category C with finite products and each functor $F : \mathbf{sSet} \rightarrow C$ that preserves finite products, let $F[\mathbf{Kan}]$ denote the following C -enriched category:

- $\text{ob } F[\mathbf{Kan}] = \text{ob } \mathbf{Kan}$.
- For each pair of Kan complexes X and Y , the hom-object is $F[X, Y]$, where $[X, Y]$ is the exponential object in \mathbf{sSet} .
- Composition and identity morphisms are induced by F from the cartesian closed structure of \mathbf{sSet} .

The next definition is a prime example of the above construction.

Definition 1.4.5. The **homotopy category of Kan complexes** is the category $\mathbf{H} = \pi_0[\mathbf{Kan}]$. A **homotopy type** is an isomorphism class of objects in \mathbf{H} .

Proposition 1.4.6. *For each simplicial set Z , let $\eta_Z : Z_0 \rightarrow \pi_0 Z$ be the map of vertices induced by the adjunction unit $\text{id}_{\mathbf{sSet}} \Rightarrow \text{disc } \pi_0$.*

- (i) *There is a (unique) functor $\pi : \mathbf{Kan} \rightarrow \mathbf{H}$ that acts as the identity on objects and as $\eta_{[X, Y]} : [X, Y]_0 \rightarrow \pi_0[X, Y]$ on morphisms.*
- (ii) *The functor π is full, surjective on objects, and preserves finite products.*
- (iii) *\mathbf{Kan} is closed under products for all small families in \mathbf{sSet} , and \mathbf{H} has products for finite families.*
- (iv) *\mathbf{Kan} and \mathbf{H} are cartesian closed categories, and $\pi : \mathbf{Kan} \rightarrow \mathbf{H}$ is a cartesian closed functor.*

(v) A morphism $f : X \rightarrow Y$ in \mathbf{Kan} is a weak homotopy equivalence if and only if $\pi f : \pi X \rightarrow \pi Y$ is an isomorphism in \mathbf{H} .

Proof. (i). The construction of \mathbf{H} as $\pi_0[\mathbf{Kan}]$ ensures that π is indeed a functor.

(ii). It is clear from the construction of $\pi_0 Z$ as a coequaliser that $\eta_Z : Z_0 \rightarrow \pi_0 Z$ is a surjection; thus π is a full functor. It is obviously surjective on objects, and it preserves finite products because π_0 does.

(iii). By [proposition A.3.12](#), the class of Kan fibrations is closed under products for small families, so \mathbf{Kan} is as well. By claim (ii), \mathbf{H} inherits finite products from \mathbf{Kan} .

(iv). By [proposition 1.3.9](#), $[Y, K]$ is a Kan complex whenever K is, which combined with claim (iii) implies \mathbf{Kan} is cartesian closed. [Proposition A.2.11](#) says we have natural isomorphisms $[X \times Y, K] \cong [X, [Y, K]]$, so it follows that we have natural bijections

$$\pi_0[X \times Y, K] \cong \pi_0[X, [Y, K]]$$

for all X, Y , and K in \mathbf{Kan} , and this descends along π to make \mathbf{H} cartesian closed.

(v). The Joyal–Tierney definition says $f : X \rightarrow Y$ is a weak equivalence if and only if $\pi_0[f, K] : \pi_0[Y, K] \rightarrow \pi_0[X, K]$ is a bijection for all Kan complexes K ; but this is natural in K , so the Yoneda lemma implies this happens if and only if $\pi f : \pi X \rightarrow \pi Y$ is an isomorphism in \mathbf{H} . ■

Proposition 1.4.7.

- (i) For each simplicial set X , there exists a Kan complex RX such that the functors $\pi_0[X, -], \pi_0[RX, -] : \mathbf{Kan} \rightarrow \mathbf{Set}$ are isomorphic.
- (ii) For each simplicial set X , the functor $\pi_0[X, -] : \mathbf{Kan} \rightarrow \mathbf{Set}$ factors through $\pi : \mathbf{Kan} \rightarrow \mathbf{H}$ as a representable functor on \mathbf{H} .
- (iii) The functor $\pi : \mathbf{Kan} \rightarrow \mathbf{H}$ extends to a functor $\pi : \mathbf{sSet} \rightarrow \mathbf{H}$ that sends weak homotopy equivalences to isomorphisms, and this extension is unique up (not necessarily unique) isomorphism.

Proof. (i). By [proposition 1.3.14](#), there is an anodyne extension $i : X \rightarrow RX$ where RX is a Kan complex; but [proposition 1.3.12](#) says that anodyne extensions are weak homotopy equivalences, so $\pi_0[i, K] : \pi_0[RX, K] \rightarrow \pi_0[X, K]$ is a bijection natural in K , as required.

(ii). The claim is certainly true if X were a Kan complex, and by claim (i), $\pi_0[X, -]$ is always isomorphic to $\pi_0[RX, -]$ for some Kan complex RX .

(iii). Formally, what we seek is a functor $F : \mathbf{sSet} \rightarrow \mathbf{H}$ such that, for all Kan complexes Y and K ,

$$\mathbf{H}(FY, \pi K) = \pi_0[Y, K]$$

and, for all weak homotopy equivalences $f : X \rightarrow Y$ in \mathbf{sSet} , the induced hom-set map $\mathbf{H}(Ff, \pi K) : \mathbf{H}(FY, \pi K) \rightarrow \mathbf{H}(FX, \pi K)$ is a bijection for all Kan complexes K . Clearly, for any such F and any simplicial set X , there must be bijections

$$\mathbf{H}(FX, \pi K) \cong \pi_0[X, K]$$

that are natural in K , but by claim (ii), this is representable as a functor $\mathbf{H} \rightarrow \mathbf{Set}$ for each X , so we can certainly construct such a functor F , and it is unique up to isomorphism. ■

Proposition 1.4.8. *Let $F : \mathbf{Kan} \rightarrow C$ be any functor that sends trivial Kan fibrations in \mathbf{Kan} to isomorphisms in C .*

- (i) *If $f_0, f_1 : X \rightarrow Y$ are a parallel pair of morphisms in \mathbf{Kan} and there exists an intrinsic homotopy $f_0 \Rightarrow f_1$, then $Ff_0 = Ff_1$.*
- (ii) *If $f_0, f_1 : X \rightarrow Y$ are an intrinsically homotopic pair of morphisms in \mathbf{Kan} , then $Ff_0 = Ff_1$.*
- (iii) *There exists a unique functor $\bar{F} : \mathbf{H} \rightarrow C$ such that $F = \bar{F}\pi$.*

Proof. (i). By [remark 1.4.2](#), given any intrinsic homotopy $\alpha : f_0 \Rightarrow f_1$, we may construct a morphism $\hat{\alpha} : X \rightarrow [\Delta^1, Y]$ such that $[\delta^1, Y] \circ \hat{\alpha} = f_0$ and $[\delta^0, Y] \circ \hat{\alpha} = f_1$. Clearly, $\delta^1 : \Delta^0 \rightarrow \Delta^1$ is isomorphic to the horn inclusion $\Lambda_0^1 \hookrightarrow \Delta^1$, and $\delta^0 : \Delta^0 \rightarrow \Delta^1$ is isomorphic to the horn inclusion $\Lambda_1^1 \hookrightarrow \Delta^1$, so by [proposition 1.3.9](#), the morphisms $[\delta^1, Y], [\delta^0, Y] : [\Delta^1, Y] \rightarrow Y$ are both trivial Kan fibrations. Thus, we must have $Ff_0 = Ff_1$.

(ii). [Proposition 1.4.3](#) implies that f_0 and f_1 are intrinsically homotopic if and only if there exists an intrinsic homotopy $f_0 \Rightarrow f_1$, so this reduces to claim (i).

(iii). The uniqueness of $\overline{F} : \mathbf{H} \rightarrow \mathcal{C}$ is an immediate corollary of the fact that $\pi : \mathbf{Kan} \rightarrow \mathbf{H}$ is full and surjective on objects; it remains to be shown that such an \overline{F} exists. However, given any parallel pair $f_0, f_1 : X \rightarrow Y$ in \mathbf{Kan} , by the construction of \mathbf{H} , we have $\pi f_0 = \pi f_1$ if and only if f_0 and f_1 are intrinsically homotopic, so F indeed factors through π . ■

Corollary 1.4.9.

- (i) *Any functor $F : \mathbf{Kan} \rightarrow \mathcal{C}$ that sends trivial Kan fibrations in \mathbf{Kan} to isomorphisms in \mathcal{C} must also send weak homotopy equivalences in \mathbf{Kan} to isomorphisms in \mathcal{C} .*
- (ii) *\mathbf{H} is the localisation of \mathbf{Kan} away from weak homotopy equivalences.*
- (iii) *If $\mathbf{Ho sSet}$ is the localisation of \mathbf{sSet} away from weak homotopy equivalences, then the functor $\pi : \mathbf{sSet} \rightarrow \mathbf{H}$ induces a functor $\mathbf{Ho sSet} \rightarrow \mathbf{H}$ that is fully faithful and essentially surjective on objects.*

Proof. (i). The above proposition says $F = \overline{F}\pi$ for some \overline{F} , and we know from [proposition 1.4.6](#) that π inverts weak homotopy equivalences, so F must also invert weak homotopy equivalences.

(ii). This is a restatement of claim (iii) of the above proposition.

(iii). Apply [proposition 1.4.7](#). ■

REMARK 1.4.10. Fixing a fibrant replacement functor $R : \mathbf{sSet} \rightarrow \mathbf{sSet}$ as in [proposition 1.3.14](#), we have the following explicit construction of $\mathbf{Ho sSet}$:

- The objects are simplicial sets.
- For any two simplicial sets X and Y , $\mathbf{Ho sSet}(X, Y) = \pi_0[RX, RY]$.
- Composition and identity morphisms are constructed as in \mathbf{H} .
- The localising functor $\gamma : \mathbf{sSet} \rightarrow \mathbf{Ho sSet}$ inverting weak homotopy equivalences is the one sending $f : X \rightarrow Y$ to the homotopy class of $Rf : RX \rightarrow RY$.

The **homotopy category of simplicial sets** is the category $\mathbf{Ho sSet}$.

Definition 1.4.11. An **intrinsic homotopy equivalence** in \mathbf{sSet} is a pair (f, g) , where $f : X \rightarrow Y$ and $g : Y \rightarrow X$ are morphisms in \mathbf{sSet} such that $g \circ f \sim \text{id}_X$ and $f \circ g \sim \text{id}_Y$. Two morphisms $f : X \rightarrow Y$ and $g : Y \rightarrow X$ are **mutual intrinsic homotopy inverses** when (f, g) constitute an intrinsic homotopy equivalence.

Proposition 1.4.12 (Formal Whitehead theorem).

- (i) *If (f, g) is an intrinsic homotopy equivalence in \mathbf{Kan} , then πf and πg are mutual inverses in \mathbf{H} .*
- (ii) *A morphism in \mathbf{Kan} is a weak homotopy equivalence if and only if it has a intrinsic homotopy inverse.*

Proof. The claims are immediate consequences of propositions 1.4.6 and 1.4.9 applied to the definition of intrinsic homotopy equivalence. ■

HOMOTOPICAL CATEGORIES

2.1 Basics

Prerequisites. § A.4.

Definition 2.1.1. A relative category C is a **category with weak equivalences** if it is semi-saturated and $\text{weq } C$ has the 2-out-of-3 property, and it is a **homotopical category** if $\text{weq } C$ has the 2-out-of-6 property. A **homotopical functor** is a relative functor between homotopical categories.

REMARK 2.1.2. If C is a relative category such that $\text{weq } C$ has the 2-out-of-6 property, then every isomorphism in C is automatically a weak equivalence. Indeed, suppose $f : X \rightarrow Y$ and $g : Y \rightarrow X$ are mutual inverses in C ; then the fact that $g \circ f = \text{id}_X$ and $f \circ g = \text{id}_Y$ are in $\text{weq } C$ implies that f and g must also be in $\text{weq } C$. Recalling [lemma A.4.13](#), it follows that every homotopical category is a category with weak equivalences.

¶ 2.1.3. To simplify notation, we will usually not distinguish between $\text{und } C$ and C . For example, when C and D are relative categories, then by ‘ordinary functor $C \rightarrow D$ ’ we mean a functor $\text{und } C \rightarrow \text{und } D$.

Example 2.1.4. Any saturated relative category is automatically a homotopical category, by [corollary A.4.14](#). In particular, any minimal saturated relative category is a homotopical category. On the other hand, any maximal relative category is obviously a homotopical category.

REMARK 2.1.5. A relative category C is a category with weak equivalences or a homotopical category if and only if the opposite relative category C^{op} is.

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Lemma 2.1.6. *Let A be an object in a homotopical category (resp. category with weak equivalences) C . Then the slice category $C_{/A}$ is also a homotopical category (resp. category with weak equivalences) if we declare a morphism in $C_{/A}$ to be a weak equivalence if and only if it is a weak equivalence in C .*

Proof. Use lemma A.4.13 on the projection functor $C_{/A} \rightarrow C$. ■

Lemma 2.1.7. *Any relative subcategory \mathcal{D} of a homotopical category (resp. category with weak equivalences) C is also a homotopical category (resp. category with weak equivalences).*

Proof. Use lemma A.4.13 on the inclusion $\mathcal{D} \hookrightarrow C$. ■

Lemma 2.1.8. *Let C and \mathcal{D} be two relative categories. If \mathcal{D} is a homotopical category (resp. category with weak equivalences), then the relative functor category $[C, \mathcal{D}]_h$ is also a homotopical category (resp. category with weak equivalences).*

Proof. This is a straightforward check. ◆

Definition 2.1.9. Two objects in a relative category are **weakly equivalent** if they can be connected by a zigzag of weak equivalences; we write $X \stackrel{w}{\simeq} Y$ to mean that X and Y are weakly equivalent.

REMARK 2.1.10. If X and Y are weakly equivalent in a relative category C , then they are isomorphic in $\text{Ho } C$. The converse is certainly true if C is saturated, but is false if C is not semi-saturated.

Definition 2.1.11. A parallel pair of morphisms in a relative category C are **weakly homotopic** if they are equal in $\text{Ho } C$; we write $f \stackrel{w}{\sim} g$ to mean that f and g are weakly homotopic.

Definition 2.1.12. An **equivalence** in a relative category C is a pair (f, g) , where $f : X \rightarrow Y$ and $g : Y \rightarrow X$ are morphisms in C such that $g \circ f \stackrel{w}{\sim} \text{id}_X$ and $f \circ g \stackrel{w}{\sim} \text{id}_Y$. Two morphisms $f : X \rightarrow Y$ and $g : Y \rightarrow X$ in C are **mutual quasi-inverses** when (f, g) constitute an equivalence in C .

REMARK 2.1.13. It follows from the definitions that quasi-inverses are unique up to weak homotopy.

Lemma 2.1.14. *If the localisation functor $\gamma : C \rightarrow \text{Ho } C$ for a relative category C is full, then the following are equivalent for all morphisms $f : X \rightarrow Y$ in C :*

- f is a morphism in C and has a quasi-inverse.
- γf is an isomorphism in C .

Proof. Obvious. ◆

REMARK 2.1.15. Clearly, any isomorphism in any relative category has a quasi-inverse; but this implies that in a relative category that is *not* semi-saturated, a morphism that has a quasi-inverse need *not* be a weak equivalence. On other hand, if f is a morphism in a *saturated* homotopical category and f has a quasi-inverse, then f must be a weak equivalence.

Definition 2.1.16. A relative category C has the **Whitehead property** when the following are equivalent:

- f is a weak equivalence in C .
- f is a morphism in C and has a quasi-inverse.

Theorem 2.1.17. *Let C be a relative category. The following are equivalent:*

- (i) C has the Whitehead property.
- (ii) The localisation functor $\gamma : C \rightarrow \text{Ho } C$ is full, and C is a saturated homotopical category.

Proof. (i) \Rightarrow (ii). By [theorem A.4.21](#), every morphism $\gamma X_0 \rightarrow \gamma X_n$ in $\text{Ho } C$ is of the form

$$(\gamma f_n)^{-1} \circ \dots \circ \gamma h_2 \circ (\gamma f_1)^{-1} \circ \gamma h_1$$

for some morphisms $h_1 : X_0 \rightarrow Y_1$, $f_1 : X_1 \rightarrow Y_1$, $h_2 : X_1 \rightarrow Y_2$, etc. in C , where f_1, \dots, f_n are weak equivalences. By the Whitehead property, each $f_i : X_i \rightarrow Y_i$ has a quasi-inverse in C , say $g_i : Y_i \rightarrow X_i$. Since $\gamma g_i = (\gamma f_i)^{-1}$, it follows that

$$(\gamma f_n)^{-1} \circ \dots \circ h_2 \circ (\gamma f_1)^{-1} \circ \gamma h_1 = \gamma(g_n \circ \dots \circ h_2 \circ g_1 \circ h_1)$$

and therefore $\gamma : C \rightarrow \text{Ho } C$ is indeed full.

In particular, every morphism $f : X \rightarrow Y$ in C such that $\gamma f : \gamma X \rightarrow \gamma Y$ is an isomorphism in $\text{Ho } C$ must have a quasi-inverse, and hence must be a weak equivalence, in view of the Whitehead property. We therefore conclude that C is a saturated homotopical category.

(ii) \Rightarrow (i). The converse follows from the definitions and [lemma 2.1.14](#). ■

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REMARK 2.1.18. The Whitehead property is in general not inherited by slice categories or by functor categories. For example, if $q \circ f = p$ and g is a quasi-inverse for f , it is only guaranteed that $q \overset{w}{\sim} p \circ g$.

Definition 2.1.19. Let $F, G : C \rightarrow D$ be two ordinary functors between relative categories. A **natural weak equivalence** $\alpha : F \Rightarrow G$ is a natural transformation such that $\alpha_C : FC \rightarrow GC$ is a weak equivalence in D for all objects C in C , and we say F and G are **naturally weakly equivalent** if they can be connected by a *zigzag* of natural weak equivalences.

REMARK 2.1.20. This is precisely the notion of weak equivalence in the relative functor category $[\text{min und } C, D]_h$. Although the definition above applies to all functors, if $H : D \rightarrow E$ is an ordinary functor, then the natural transformation $H\alpha : HF \Rightarrow HG$ is only guaranteed to be a natural weak equivalence if we assume H is a relative functor.

Definition 2.1.21. A **relative equivalence** is a relative functor $F : C \rightarrow D$ for which there exists a relative functor $G : D \rightarrow C$ such that GF is naturally weakly equivalent to id_C and FG is naturally weakly equivalent to id_D . Such a G is said to be a **relative inverse** of F . If C and D are homotopical categories then we may say **homotopical equivalence** and **homotopical inverse** instead of ‘relative equivalence’ and ‘relative inverse’.

Proposition 2.1.22. *If $F : C \rightarrow D$ is a relative equivalence of relative categories with relative inverse $G : D \rightarrow C$, then $\text{Ho } F : \text{Ho } C \rightarrow \text{Ho } D$ is an equivalence of categories, with quasi-inverse $\text{Ho } G : \text{Ho } D \rightarrow \text{Ho } C$. ■*

2.2 Homotopical Kan extensions

Prerequisites. §§ 2.1, A.4.

Definition 2.2.1. Let C be a homotopical category. A **homotopically initial object** in C is an object A for which there exists a zigzag of natural transformations of the form

$$\Delta A \overset{\sim}{\sim} F \xrightarrow{\alpha} G \overset{\sim}{\sim} \text{id}_C$$

where $\Delta A : C \rightarrow C$ is the constant functor with value A , $\alpha_A : FA \rightarrow GA$ is a weak equivalence in C , and the squiggles denote (possibly trivial) zigzags of

natural weak equivalences. Dually, a **homotopically terminal object** in C is a homotopically initial object in C^{op} .

Proposition 2.2.2. *Let C be a homotopical category. If A is a homotopically initial (resp. homotopically terminal) object in C , then:*

- (i) *Any object in C weakly equivalent to A is also a homotopically initial (resp. homotopically terminal) object in C .*
- (ii) *A is an initial (resp. terminal) object in $\text{Ho } C$.*
- (iii) *If C is a minimal homotopical category, then A is an initial (resp. terminal) object in C as well.*

Conversely, any initial (resp. terminal) object in C is also homotopically initial (resp. homotopically terminal).

Proof. Obvious. (This is Proposition 38.3 in [DHKS].) ◆

Proposition 2.2.3. *If A is a homotopically initial object in a homotopical category C , then for any object Z in C , the zigzag category $C^{(\mathbf{T})}(A, Z)$ is connected.*

Proof. By theorem A.4.2I, there is a bijection between the connected components of $C^{(\mathbf{T})}(A, Z)$ and the morphisms $A \rightarrow Z$ in $\text{Ho } C$; but we know A is an initial object in $\text{Ho } C$, so $C^{(\mathbf{T})}(A, Z)$ has exactly one connected component. ■

Lemma 2.2.4. *Let $H : C \rightarrow D$ be a relative functor and let $F : C \rightarrow D$ be an ordinary functor. If $\text{weq } D$ has the 2-out-of-3 property and F is naturally weakly equivalent to H , then F is also a relative functor.*

Proof. Apply the 2-out-of-3 property inductively. ◆

Lemma 2.2.5. *If A and A' be homotopically initial objects in a homotopical category C , then $A \stackrel{\text{w}}{\simeq} A'$, and moreover every morphism $A \rightarrow A'$ in C is a weak equivalence.*

Proof. This is paragraph 38.5 in [DHKS].

Suppose, as in the definition, that we have endofunctors F, F', G, G' on C and natural transformations $\alpha : F \Rightarrow G, \alpha' : F' \Rightarrow G'$, such that $F \stackrel{\text{w}}{\simeq} \Delta A, F' \stackrel{\text{w}}{\simeq} \Delta A', G \stackrel{\text{w}}{\simeq} \text{id}_C,$ and $G' \stackrel{\text{w}}{\simeq} \text{id}_C,$ and the morphisms $\alpha_A : FA \rightarrow GA$ and $\alpha'_{A'} : FA' \rightarrow GA'$ are both weak equivalences. Note that the previous lemma

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implies G and G' are both homotopical functors, while a similar argument shows that F and F' sends *all* morphisms to weak equivalences.

Let $f : A \rightarrow A'$ be a morphism in \mathcal{C} . By applying the 2-out-of-3 property repeatedly in the following diagram,

$$\begin{array}{ccccc}
 A & \rightsquigarrow & FA & \xrightarrow{\alpha_A} & GA & \rightsquigarrow & A \\
 \parallel & & \downarrow Ff & & \downarrow Gf & & \downarrow f \\
 A & \rightsquigarrow & FA' & \xrightarrow{\alpha_{A'}} & GA' & \rightsquigarrow & A'
 \end{array}$$

we see that f is a weak equivalence if and only if $\alpha_{A'} : FA' \rightarrow GA'$ is a weak equivalence. Since $\alpha'_{A'} : F'A' \rightarrow G'A'$ is a weak equivalence, and $GA' \stackrel{w}{\simeq} A'$, it follows that $\alpha'_{GA'} : FGA' \rightarrow G'GA'$ is a weak equivalence, and since G is homotopical, so $G\alpha'_{GA'} : GFGA' \rightarrow GG'GA'$ is also a weak equivalence. Similarly, $\alpha_A : FA \rightarrow GA$ is a weak equivalence, and $A \stackrel{w}{\simeq} FA' \stackrel{w}{\simeq} G'FA'$, so $\alpha_{G'FA'} : FG'FA' \rightarrow GG'FA'$ is a weak equivalence as well.

Now, by applying the 2-out-of-6 property to the diagram below,

$$\begin{array}{ccccc}
 FF'FA' & \xrightarrow{\alpha_{F'FA'}} & GF'FA' & \xrightarrow{GF'\alpha_{A'}} & GF'GA' \\
 F\alpha'_{FA'} \downarrow & & G\alpha'_{FA'} \downarrow & & G\alpha'_{GA'} \downarrow \\
 FG'FA' & \xrightarrow{\alpha_{G'FA'}} & GG'FA' & \xrightarrow{GG'\alpha_{A'}} & GG'GA'
 \end{array}$$

we may deduce that $GG'\alpha_{A'} : GG'FA' \rightarrow GG'GA'$ is a weak equivalence, and hence that $\alpha_{A'} : FA' \rightarrow GA'$ is a weak equivalence, as required. \blacksquare

Definition 2.2.6. A **homotopically contractible category** is a homotopical category \mathcal{C} such that the unique (homotopical) functor $\mathcal{C} \rightarrow \mathbb{1}$ is a homotopical equivalence, where $\mathbb{1}$ is the trivial category with only one object.

¶ 2.2.7. We will say that an object in a homotopical category \mathcal{C} characterised by a homotopical universal property is **homotopically unique** if the full subcategory spanned by such objects inside the homotopical category of objects in \mathcal{C} equipped with the relevant additional structure.

Proposition 2.2.8. *Let \mathcal{C} be a homotopical category. The following are equivalent:*

- (i) \mathcal{C} is homotopically contractible.

- (ii) C is inhabited, and for every object A in C , the constant functor ΔA is naturally weakly equivalent to id_C .
- (iii) There exists an object A in C such that ΔA and id_C are naturally weakly equivalent.

Proof. Obvious. (This is paragraph 37.6 in [DHKS].) ◆

Proposition 2.2.9. *Let C be a homotopically contractible category.*

- (i) Every morphism in C is a weak equivalence.
- (ii) The unique functor $\text{Ho } C \rightarrow \mathbb{1}$ is an equivalence of categories.
- (iii) If C is a minimal homotopical category, then $C \rightarrow \mathbb{1}$ is also an equivalence of categories.
- (iv) The opposite homotopical category C^{op} and the homotopical functor category $[D, C]_{\text{h}}$ (for any homotopical category D) are also homotopically contractible.
- (v) Every object in C is both homotopically initial and homotopically terminal.

Proof. Obvious. (This is paragraph 37.6 in [DHKS].) ◆

Proposition 2.2.10. *Let C be a homotopical category. If \mathcal{D} is the full homotopical subcategory of C spanned by the homotopically initial (or homotopically terminal) objects, then \mathcal{D} is homotopically contractible.*

Proof. This follows from lemma 2.2.5. ■

REMARK 2.2.11. Even if C is a saturated homotopical category, an object that is initial in $\text{Ho } C$ need not be homotopically initial in C . Indeed, let C be the maximal homotopical category generated by a graph of the following form:



No object in C is homotopically initial, because the length of the shortest zigzag connecting two objects cannot be bounded above; yet every object in $\text{Ho } C$ is initial. The same argument shows that C is not homotopically contractible, but $\text{Ho } C$ is certainly contractible.

Definition 2.2.12. Let $F : C \rightarrow D$ and $G : C \rightarrow \mathcal{E}$ be two ordinary functors between homotopical categories. A **homotopical left Kan extension** (resp. **homotopical right Kan extension**) of G along F is a homotopically initial (resp. homotopically terminal) object of the homotopical category $(G \downarrow F^*)_h$ (resp. $(F^* \downarrow G)_h$) described below:

- The objects are pairs (H, α) where H is a homotopical functor $D \rightarrow \mathcal{E}$ and α is a natural transformation of type $G \Rightarrow HF$ (resp. $HF \Rightarrow G$).
- The morphisms $(H', \alpha') \rightarrow (H, \alpha)$ are those natural transformations $\beta : H' \Rightarrow H$ such that $\beta F \bullet \alpha' = \alpha$ (resp. $\alpha \bullet \beta F = \alpha'$).
- The weak equivalences are the natural weak equivalences.

REMARK 2.2.13. Note that any homotopical Kan extension of $F : C \rightarrow D$ along $G : C \rightarrow \mathcal{E}$ has, by definition, an underlying *homotopical* functor $D \rightarrow \mathcal{E}$.

Corollary 2.2.14. *Homotopical Kan extensions are homotopically unique, any two homotopical left (resp. right) Kan extensions of G along F are naturally weakly equivalent.* ■

Definition 2.2.15. Let $F : C \rightarrow D$ and $G : C \rightarrow \mathcal{E}$ be two ordinary functors between homotopical categories, and let $L : \mathcal{E} \rightarrow \mathcal{F}$ be a homotopical functor. We say L **preserves** a homotopical left (resp. right) Kan extension (H, α) of G along F if $(LH, L\alpha)$ is a homotopical left (resp. right) Kan extension of LF along G . If a homotopical Kan extension is preserved by *all* homotopical functors, then it is said to be **absolute**.

2.3 Quillen–Verdier derived functors

Prerequisites. §§ 2.1, A.4, A.1, A.5

The fact that $\text{Ho} : \mathfrak{RelCat} \rightarrow \mathfrak{Cat}$ is a 2-functor means that relative functors $F : C \rightarrow D$ descend to functors $\text{Ho } F : \text{Ho } C \rightarrow \text{Ho } D$ in a very well-behaved way. However, what can we say about ordinary (i.e. not necessarily relative) functors $C \rightarrow D$?

In this section, we follow [DHKS, §§ 40–43].

Definition 2.3.1. Let C and D be relative categories. A **left deformation retract** for an ordinary functor $F : C \rightarrow D$ is a triple (C°, Q, p) where

- $Q : C \rightarrow C$ is a relative functor,
- C° is a full subcategory of C with the induced relative subcategory structure, and
- $p : Q \Rightarrow \text{id}_C$ is a natural weak equivalence,

and these data are required to have the following properties:

- For all objects X in C , the object QX is in C° .
- The restriction $F|_{C^\circ} : C^\circ \rightarrow D$ is a relative functor.

An ordinary functor $F : C \rightarrow D$ is **left deformable** if there exists a left deformation retract for F .

Dually, a **right deformation retract** for an ordinary functor $G : D \rightarrow C$ is a triple (D°, R, i) where

- $R : D \rightarrow D$ is a relative functor,
- D° is a full subcategory of D with the induced relative subcategory structure, and
- $i : \text{id}_D \Rightarrow R$ is a natural weak equivalence,

and these data are required to have the following properties:

- For all objects A in D , the object RA is in D° .
- The restriction $G|_{D^\circ} : D^\circ \rightarrow C$ is a relative functor.

An ordinary functor $G : D \rightarrow C$ is **right deformable** if there exists a left right deformation retract for G .

REMARK 2.3.2. Every relative functor is both left deformable and right deformable, with trivial left and right deformation retracts.

Lemma 2.3.3. *Let C and D be relative categories.*

- *If (C°, Q, p) is a left deformation retract for $F : C \rightarrow D$, then the composite $FQ : C \rightarrow D$ is a relative functor.*
- *If (D°, R, i) is a right deformation retract for $G : D \rightarrow C$, then the composite $GR : D \rightarrow C$ is a relative functor.*

Proof. Immediate from the definitions. ■

The following definition is essentially due to Verdier [1963], but we use the formulation of Quillen [1967, Ch. I, § 4].

Definition 2.3.4. Let \mathcal{C} and \mathcal{D} be relative categories, and let $\gamma_{\mathcal{C}} : \mathcal{C} \rightarrow \text{Ho } \mathcal{C}$ and $\gamma_{\mathcal{D}} : \mathcal{D} \rightarrow \text{Ho } \mathcal{D}$ be the localising functors. A **total left derived functor** for an ordinary functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is a right (!) Kan extension of $\gamma_{\mathcal{D}}F : \mathcal{C} \rightarrow \text{Ho } \mathcal{D}$ along $\gamma_{\mathcal{C}} : \mathcal{C} \rightarrow \text{Ho } \mathcal{C}$. Dually, a **total right derived functor** for an ordinary functor $G : \mathcal{D} \rightarrow \mathcal{C}$ is a left (!) Kan extension of $\gamma_{\mathcal{C}}G : \mathcal{D} \rightarrow \text{Ho } \mathcal{C}$ along $\gamma_{\mathcal{D}} : \mathcal{D} \rightarrow \text{Ho } \mathcal{D}$.

REMARK 2.3.5. As with everything with a universal property, total derived functors are unique up to unique isomorphism if they exist.

Theorem 2.3.6. *With other notation as in the definition:*

- (i) *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an ordinary functor. If $(\mathcal{C}^\circ, \mathcal{Q}, p)$ is a left deformation retract for F , then $(\text{Ho}(F\mathcal{Q}), \gamma_{\mathcal{D}}Fp)$ is an absolute right Kan extension of $\gamma_{\mathcal{D}}F : \mathcal{C} \rightarrow \text{Ho } \mathcal{D}$ along $\gamma_{\mathcal{C}} : \mathcal{C} \rightarrow \text{Ho } \mathcal{C}$.*
- (ii) *Let $F, F' : \mathcal{C} \rightarrow \mathcal{D}$ be a parallel pair of ordinary functors. If $(\mathbf{L}F, \delta)$ and $(\mathbf{L}F', \delta')$ are total left derived functors for F and F' (respectively), then for any natural transformation $\varphi : F \Rightarrow F'$, there exists a unique natural transformation $\mathbf{L}\varphi : \mathbf{L}F \Rightarrow \mathbf{L}F'$ such that $\delta' \cdot (\mathbf{L}\varphi)\gamma_{\mathcal{C}} = \gamma_{\mathcal{D}}\varphi \cdot \delta$.*
- (iii) *Moreover, if $(\mathcal{C}^\circ, \mathcal{Q}, p)$ is a left deformation retract for both F and F' , then we may take $\mathbf{L}\varphi = \text{Ho}(\varphi\mathcal{Q})$.*
- (iv) *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{D} \rightarrow \mathcal{E}$ be ordinary functors between relative categories. If $(\mathbf{L}F, \delta^F)$, $(\mathbf{L}G, \delta^G)$, and $(\mathbf{L}(GF), \delta^{GF})$ are total left derived functors for F , G , and GF (respectively), then there is a unique natural transformation $\mu_{G,F} : (\mathbf{L}G)(\mathbf{L}F) \Rightarrow \mathbf{L}(GF)$ such that $\delta^{GF} \cdot \mu_{G,F}\gamma_{\mathcal{C}} = \delta^G F \cdot (\mathbf{L}G)\delta^F$.*
- (v) *If $(\mathcal{C}^\circ, \mathcal{Q}^{\mathcal{C}^\circ}, p^{\mathcal{C}^\circ})$ is a left deformation retract for F , $(\mathcal{D}^\circ, \mathcal{Q}^{\mathcal{D}^\circ}, p^{\mathcal{D}^\circ})$ is a left deformation retract for G , and F maps objects in \mathcal{C}° to objects in \mathcal{D}° , then $(\mathcal{C}^\circ, \mathcal{Q}^{\mathcal{C}^\circ}, p^{\mathcal{C}^\circ})$ is also a left deformation retract for GF , and the canonical comparison $\mu_{G,F} : (\mathbf{L}G)(\mathbf{L}F) \Rightarrow \mathbf{L}(GF)$ is an isomorphism.*

Dually:

- (i') If $G : \mathcal{D} \rightarrow \mathcal{C}$ is an ordinary functor and $(\mathcal{D}^\circ, \mathbf{R}, i)$ is a right deformation retract for F , then $(\mathbf{Ho}(GR), \gamma_C Gi)$ is an absolute left Kan extension of $\gamma_C G : \mathcal{D} \rightarrow \mathbf{Ho} \mathcal{C}$ along $\gamma_D : \mathcal{D} \rightarrow \mathbf{Ho} \mathcal{D}$.
- (ii') Let $G, G' : \mathcal{D} \rightarrow \mathcal{C}$ be a parallel pair of ordinary functors. If $(\mathbf{R}G, \delta)$ and $(\mathbf{R}G', \delta')$ are total right derived functors for G and G' (respectively), then for any natural transformation $\psi : G' \Rightarrow G$, there exists a unique natural transformation $\mathbf{R}\psi : \mathbf{R}G' \Rightarrow \mathbf{R}G$ such that $(\mathbf{R}\psi)\gamma_D \bullet \delta' = \delta \bullet \gamma_C \psi$.
- (iii') Moreover, if $(\mathcal{D}^\circ, \mathbf{R}, i)$ is a right deformation retract for both G and G' , then we may take $\mathbf{R}\psi = \mathbf{Ho}(\psi \mathbf{R})$.
- (iv') Let $F : \mathcal{C} \rightarrow \mathcal{B}$ and $G : \mathcal{D} \rightarrow \mathcal{C}$ be ordinary functors between relative categories. If $(\mathbf{R}F, \delta^F)$, $(\mathbf{R}G, \delta^G)$, and $(\mathbf{R}(FG), \delta^{FG})$ are total right derived functors for F , G , and FG (respectively), then there is a unique natural transformation $\delta_{F,G} : \mathbf{R}(FG) \Rightarrow (\mathbf{R}F)(\mathbf{R}G)$ such that $\delta_{F,G} \gamma_D \bullet \delta^{FG} = (\mathbf{R}F)\delta^G \bullet \delta^F G$.
- (v') If $(\mathcal{C}^\circ, \mathbf{R}^\circ, i^{\mathcal{C}^\circ})$ is a right deformation retract for F , $(\mathcal{D}^\circ, \mathbf{R}^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ is a right deformation retract for G , and G maps objects in \mathcal{D}° to objects in \mathcal{C}° , then $(\mathcal{D}^\circ, \mathbf{Q}^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ is also a right deformation retract for FG , and the canonical comparison $\delta_{F,G} : \mathbf{R}(FG) \Rightarrow (\mathbf{R}F)(\mathbf{R}G)$ is an isomorphism.

Proof. (i). To simplify notation, let $\mathbf{L}F = \mathbf{Ho}(FQ)$. Let $H : \mathbf{Ho} \mathcal{D} \rightarrow \mathcal{E}$ and $K : \mathbf{Ho} \mathcal{C} \rightarrow \mathcal{E}$ be any two ordinary functors, and let $\alpha : K\gamma_C \Rightarrow H\gamma_D F$ be any natural transformation. Observe that the following diagrams commute for every object X in \mathcal{C} :

$$\begin{array}{ccc}
 \gamma_C Q Q X & \xrightarrow{\gamma_C p Q X} & \gamma_C Q X \\
 \gamma_C Q p X \downarrow & & \downarrow \gamma_C p X \\
 \gamma_C Q X & \xrightarrow{\gamma_C p X} & \gamma_C X
 \end{array}
 \qquad
 \begin{array}{ccc}
 \gamma_D F Q Q Q X & \xrightarrow{\gamma_D F p Q Q X} & \gamma_D F Q Q X \\
 \gamma_D F Q p Q X \downarrow & & \downarrow \gamma_D F p Q X \\
 \gamma_D F Q Q X & \xrightarrow{\gamma_D F p Q X} & \gamma_D F Q X
 \end{array}$$

Since $\gamma_C p$ and $\gamma_D F p Q$ are natural isomorphisms, we must have these equalities:

$$\gamma_C Q p = \gamma_C p Q \qquad \gamma_D F Q p Q = \gamma_D F p Q Q$$

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Now, suppose we are given $\bar{\alpha} : K \Rightarrow H(\mathbf{L}F)$ such that $\alpha = H\gamma_D Fp \cdot \bar{\alpha}\gamma_C$. Then:

$$\begin{aligned}\bar{\alpha}\gamma_C &= \bar{\alpha}\gamma_C \cdot K\gamma_C p \cdot K(\gamma_C p)^{-1} \\ &= H(\mathbf{L}F)\gamma_C p \cdot \bar{\alpha}\gamma_C Q \cdot K(\gamma_C p)^{-1} \\ &= H\gamma_D FQp \cdot \bar{\alpha}\gamma_C Q \cdot K(\gamma_C p)^{-1}\end{aligned}$$

and so, recursing once and applying the equations above,

$$\begin{aligned}\bar{\alpha}\gamma_C &= H\gamma_D FQp \cdot \left(H\gamma_D FQp \cdot \bar{\alpha}\gamma_C Q \cdot K(\gamma_C p)^{-1} \right) Q \cdot K(\gamma_C p)^{-1} \\ &= H\gamma_D FQp \cdot H\gamma_D FQpQ \cdot \bar{\alpha}\gamma_C QQ \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= H\gamma_D FQp \cdot H\gamma_D FpQQ \cdot \bar{\alpha}\gamma_C QQ \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= H\gamma_D FQp \cdot (H\gamma_D Fp \cdot \bar{\alpha}\gamma_C) QQ \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= H\gamma_D FQp \cdot \alpha QQ \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= \alpha Q \cdot K\gamma_C Qp \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= \alpha Q \cdot K\gamma_C p Q \cdot K(\gamma_C p)^{-1} Q \cdot K(\gamma_C p)^{-1} \\ &= \alpha Q \cdot K(\gamma_C p)^{-1}\end{aligned}$$

However, the 2-dimensional universal property of $\text{Ho } C$ implies that the map shown below is a bijection,

$$\begin{aligned}[\text{Ho } C, \mathcal{E}](K, H(\mathbf{L}F)) &\rightarrow [C, \mathcal{E}](K\gamma_C, H(\mathbf{L}F)\gamma_C) \\ \bar{\alpha} &\mapsto \bar{\alpha}\gamma_C\end{aligned}$$

and so this calculation determines $\bar{\alpha} : K \Rightarrow H(\mathbf{L}F)$ uniquely. Conversely, *define* $\bar{\alpha}$ to be the unique natural transformation such that $\bar{\alpha}\gamma_C = \alpha Q \cdot K(\gamma_C p)^{-1}$; then,

$$\begin{aligned}H\gamma_D Fp \cdot \bar{\alpha}\gamma_C &= H\gamma_D Fp \cdot \alpha Q \cdot K(\gamma_C p)^{-1} \\ &= \alpha \cdot K\gamma_C p \cdot K(\gamma_C p)^{-1} \\ &= \alpha\end{aligned}$$

and therefore $(\mathbf{L}F, \gamma_D Fp)$ is indeed an absolute right Kan extension of $\gamma_D F : C \rightarrow \text{Ho } D$ along $\gamma_C : C \rightarrow \text{Ho } C$.

(ii). Noting that $\gamma_D \varphi \cdot \delta$ is a natural transformation $(\mathbf{L}F)\gamma_C \Rightarrow \gamma_D F'$, the universal property of $(\mathbf{L}F', \delta')$ yields a unique natural transformation $\mathbf{L}\varphi : \mathbf{L}F \Rightarrow \mathbf{L}F'$ such that $\gamma_D \varphi \cdot \delta = \delta' \cdot (\mathbf{L}\varphi)\gamma_C$, as required.

(iii). $\text{Ho}(\varphi Q)$ is a natural transformation $\text{Ho}(FQ) \Rightarrow \text{Ho}(F'Q)$, and we have

$$\gamma_D Fp \cdot \text{Ho}(\gamma_D \varphi Q) \gamma_C = \gamma_D Fp \cdot \gamma_D \varphi Q = \gamma_D \varphi \cdot \gamma_D F'p$$

as required.

(iv). Since $\delta^G F \cdot (\mathbf{L}F)\delta$ is a natural transformation $(\mathbf{L}G)(\mathbf{L}F)\gamma_C \Rightarrow \gamma_D GF$, the universal property of $(\mathbf{L}(GF), \delta^{GF})$ yields the required natural transformation $\mu_{G,F} : (\mathbf{L}G)(\mathbf{L}F) \Rightarrow \mathbf{L}(GF)$.

(v). Our hypotheses imply that the restriction $GF|_{C^\circ} : C^\circ \rightarrow \mathcal{E}$ is a relative functor, so $(C^\circ, Q^{C^\circ}, p^{C^\circ})$ is indeed a left deformation retract for $GF : C \rightarrow \mathcal{E}$. It follows that $GQ^{D^\circ} FQ^{C^\circ}$ and GFQ^{C° are both relative functors $C \rightarrow \mathcal{E}$; moreover,

$$\begin{aligned} \gamma_\varepsilon GFp^{C^\circ} \cdot \text{Ho}(Gp^{D^\circ} FQ^{C^\circ}) \gamma_C &= \gamma_\varepsilon GFp^{C^\circ} \cdot \gamma_\varepsilon Gp^{D^\circ} FQ^{C^\circ} \\ &= \gamma_\varepsilon Gp^{D^\circ} F \cdot \gamma_\varepsilon GQ^{D^\circ} Fp^{C^\circ} \\ &= \gamma_\varepsilon Gp^{D^\circ} F \cdot \text{Ho}(GQ^{D^\circ})(\gamma_D Fp^{C^\circ}) \end{aligned}$$

so we must have $\mu_{G,F} = \text{Ho}(Gp^{D^\circ} FQ^{C^\circ})$. However, because $FQ^{C^\circ} X$ is in D° for all objects X in C , $Gp^{D^\circ} FQ^{C^\circ} : GQ^{D^\circ} FQ^{C^\circ} \Rightarrow GFQ^{C^\circ}$ must be a natural weak equivalence, and so $\mu_{G,F} : \text{Ho}(GQ^{D^\circ}) \text{Ho}(FQ^{C^\circ}) \Rightarrow \text{Ho}(GFQ^{C^\circ})$ is indeed a natural isomorphism. \blacksquare

Definition 2.3.7. The **2-category of small left deformation retracts** is defined as follows:

- The objects are pairs $(C, C^\circ, Q^{C^\circ}, p^{C^\circ})$ where C is a small relative category and $(C^\circ, Q^{C^\circ}, p^{C^\circ})$ is a left deformation retract for $\text{id} : C \rightarrow C$.
- A 1-morphism $F : (C, C^\circ, Q^{C^\circ}, p^{C^\circ}) \rightarrow (D, D^\circ, Q^{D^\circ}, p^{D^\circ})$ is an ordinary functor $F : C \rightarrow D$, such that $(C^\circ, Q^{C^\circ}, p^{C^\circ})$ is a left deformation retract for F , and F sends objects in C° to objects in D° .
- The 2-morphisms are ordinary natural transformations.
- All compositions and identities are inherited from 2-category of small categories.

We write \mathfrak{LDef} for this 2-category, and we write LDefFun for its hom-sets.

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The **2-category of small right deformation retracts** is defined dually:

- The objects are pairs $(\mathcal{D}, \mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ where \mathcal{D} is a small relative category and $(\mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ is a right deformation retract for $\text{id} : \mathcal{D} \rightarrow \mathcal{D}$.
- A 1-morphism $G : (\mathcal{D}, \mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ}) \rightarrow (\mathcal{C}, \mathcal{C}^\circ, R^{\mathcal{C}^\circ}, i^{\mathcal{C}^\circ})$ is an ordinary functor $G : \mathcal{D} \rightarrow \mathcal{C}$, such that $(\mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ is a right deformation retract for G , and G sends objects in \mathcal{D}° to objects in \mathcal{C}° .
- The 2-morphisms are ordinary natural transformations.
- All compositions and identities are inherited from 2-category of small categories.

We write \mathfrak{RDef} for this 2-category, and we write $\mathbf{RDefFun}$ for its hom-sets.

REMARK 2.3.8. The duality principle for deformation retracts can be formalised as follows: there is a 2-functor $\mathfrak{LDef}^{\text{co}} \rightarrow \mathfrak{RDef}$ that sends $(\mathcal{C}, \mathcal{C}^\circ, Q^{\mathcal{C}^\circ}, p^{\mathcal{C}^\circ})$ to its opposite $(\mathcal{C}^{\text{op}}, (\mathcal{C}^\circ)^{\text{op}}, (Q^{\mathcal{C}^\circ})^{\text{op}}, (p^{\mathcal{C}^\circ})^{\text{op}})$, and it has an evident strict inverse $\mathfrak{RDef}^{\text{co}} \rightarrow \mathfrak{LDef}$. Note that these two 2-functors reverse the direction of 2-morphisms but preserve the direction of 1-morphisms!

Corollary 2.3.9. *There are two pseudofunctors, \mathbf{L} and \mathbf{R} , where:*

- \mathbf{L} is a pseudofunctor $\mathfrak{LDef} \rightarrow \mathfrak{Cat}$ that sends an object $(\mathcal{C}, \mathcal{C}^\circ, Q^{\mathcal{C}^\circ}, p^{\mathcal{C}^\circ})$ to the homotopy category $\text{Ho } \mathcal{C}$, a 1-morphism $F : (\mathcal{C}, \mathcal{C}^\circ, Q^{\mathcal{C}^\circ}, p^{\mathcal{C}^\circ}) \rightarrow (\mathcal{D}, \mathcal{D}^\circ, Q^{\mathcal{D}^\circ}, p^{\mathcal{D}^\circ})$ to its total left derived functor $\mathbf{L}F : \text{Ho } \mathcal{C} \rightarrow \text{Ho } \mathcal{D}$, and a 2-morphism $\varphi : F \Rightarrow F'$ to the derived natural transformation $\mathbf{L}\varphi : \mathbf{L}F \Rightarrow \mathbf{L}F'$, and \mathbf{L} preserves identity 1-morphisms strictly.
- \mathbf{R} is a pseudofunctor $\mathfrak{RDef} \rightarrow \mathfrak{Cat}$ that sends an object $(\mathcal{D}, \mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ})$ to the homotopy category $\text{Ho } \mathcal{C}$, a 1-morphism $G : (\mathcal{D}, \mathcal{D}^\circ, R^{\mathcal{D}^\circ}, i^{\mathcal{D}^\circ}) \rightarrow (\mathcal{C}, \mathcal{C}^\circ, R^{\mathcal{C}^\circ}, i^{\mathcal{C}^\circ})$ to its total right derived functor $\mathbf{R}G : \text{Ho } \mathcal{D} \rightarrow \text{Ho } \mathcal{C}$, and a 2-morphism $\psi : G' \Rightarrow G$ to the derived natural transformation $\mathbf{R}\psi : \mathbf{R}G' \Rightarrow \mathbf{R}G$, and \mathbf{R} preserves identity 1-morphisms strictly.

- \mathbf{L} and \mathbf{R} are compatible with the duality principle, in the sense that the following diagrams commute (strictly):

$$\begin{array}{ccc}
 \mathfrak{L}\mathfrak{D}\mathfrak{e}\mathfrak{f}^{\text{co}} & \xrightarrow{\mathbf{L}^{\text{co}}} & \mathfrak{C}\mathfrak{a}\mathfrak{t}^{\text{co}} \\
 \downarrow (-)^{\text{op}} & & \downarrow (-)^{\text{op}} \\
 \mathfrak{R}\mathfrak{D}\mathfrak{e}\mathfrak{f} & \xrightarrow{\mathbf{R}} & \mathfrak{C}\mathfrak{a}\mathfrak{t}
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathfrak{R}\mathfrak{D}\mathfrak{e}\mathfrak{f}^{\text{co}} & \xrightarrow{\mathbf{R}^{\text{co}}} & \mathfrak{C}\mathfrak{a}\mathfrak{t}^{\text{co}} \\
 \downarrow (-)^{\text{op}} & & \downarrow (-)^{\text{op}} \\
 \mathfrak{L}\mathfrak{D}\mathfrak{e}\mathfrak{f} & \xrightarrow{\mathbf{L}} & \mathfrak{C}\mathfrak{a}\mathfrak{t}
 \end{array}$$

Proof. The main claims follow from [theorem 2.3.6](#); the only thing left to check is that the collection of 2-isomorphisms $\boldsymbol{\mu}$ and $\boldsymbol{\delta}$ satisfy the coherence laws for pseudofunctors; that is, we should show that the following diagrams commute:

$$\begin{array}{ccc}
 (\mathbf{L}H)(\mathbf{L}G)(\mathbf{L}F) & \xrightarrow{(\mathbf{L}H)\boldsymbol{\mu}_{G,F}} & (\mathbf{L}H)\mathbf{L}(GF) \\
 \downarrow \boldsymbol{\mu}_{H,G}(\mathbf{L}F) & & \downarrow \boldsymbol{\mu}_{H,GF} \\
 \mathbf{L}(HG)(\mathbf{L}F) & \xrightarrow{\boldsymbol{\mu}_{HG,F}} & \mathbf{L}(HGF)
 \end{array}$$

$$\begin{array}{ccc}
 \mathbf{R}(FGH) & \xrightarrow{\boldsymbol{\delta}_{F,GH}} & (\mathbf{R}F)\mathbf{R}(GH) \\
 \downarrow \boldsymbol{\delta}_{FG,H} & & \downarrow (\mathbf{R}F)\boldsymbol{\delta}_{G,H} \\
 \mathbf{R}(FG)(\mathbf{R}H) & \xrightarrow{\boldsymbol{\delta}_{F,G}(\mathbf{R}H)} & (\mathbf{R}F)(\mathbf{R}G)(\mathbf{R}H)
 \end{array}$$

However, using the explicit formulae for $\boldsymbol{\mu}$ and $\boldsymbol{\delta}$ in the proof of the theorem, it is easy to see that these diagrams do indeed commute. \blacksquare

Definition 2.3.10. A **deformable adjunction** between two relative categories is an ordinary adjunction where the left adjoint is left deformable and the right adjoint is right deformable.

Theorem 2.3.11. Let \mathcal{C} and \mathcal{D} be two relative categories, and let

$$F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$$

be an adjunction of ordinary categories, with

$$\begin{aligned}
 \eta : \text{id}_{\mathcal{C}} &\Rightarrow GF \\
 \varepsilon : FG &\Rightarrow \text{id}_{\mathcal{D}}
 \end{aligned}$$

as the unit and counit (respectively).

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- (i) If (C°, Q, p) is a left deformation retract for F , and (D°, R, i) is a right deformation retract for G , then for $\mathbf{L}F = \text{Ho}(FQ)$ and $\mathbf{R}G = \text{Ho}(GR)$,

$$\mathbf{L}F \dashv \mathbf{R}G : \text{Ho } \mathcal{D} \rightarrow \text{Ho } \mathcal{C}$$

is an adjunction with the following unit and counit:

$$\begin{aligned} \bar{\eta} &= \text{Ho}(GiFQ \cdot \eta Q) \cdot (\text{Ho } p)^{-1} : \text{id}_{\text{Ho } \mathcal{C}} \Rightarrow (\mathbf{R}G)(\mathbf{L}F) \\ \bar{\varepsilon} &= (\text{Ho } i)^{-1} \cdot \text{Ho}(\varepsilon R \cdot FpGR) : (\mathbf{L}F)(\mathbf{R}G) \Rightarrow \text{id}_{\text{Ho } \mathcal{D}} \end{aligned}$$

- (ii) Let $F' \dashv G' : \mathcal{D}' \rightarrow \mathcal{C}'$ be another adjunction, with unit η' and counit ε' , and let $H : \mathcal{C}' \rightarrow \mathcal{C}$ and $K : \mathcal{D}' \rightarrow \mathcal{D}$ be homotopical functors. If

- (C°, Q, p) is a left deformation retract for F ,
- (C'°, Q', p') is a left deformation retract for F' ,
- H sends objects in C'° to objects in C° ,
- (D°, R, i) is a right deformation retract for G ,
- (D'°, R', i') is a right deformation retract for G' , and
- K sends objects in D'° to objects in D° ,

then for any conjugate pair of natural transformations $\varphi : FH \Rightarrow KF'$, $\psi : HG' \Rightarrow GK$,^[1] i.e. a pair (φ, ψ) satisfying the equations below,

$$\varepsilon K \cdot F\psi = K\varepsilon' \cdot \varphi G' \qquad G\varphi \cdot \eta H = \psi F' \cdot H\eta'$$

the derived natural transformations $\mathbf{L}\varphi : (\mathbf{L}F)(\text{Ho } H) \Rightarrow (\text{Ho } K)(\mathbf{L}F')$ and $\mathbf{R}\psi : (\text{Ho } K)(\mathbf{R}G') \Rightarrow (\mathbf{R}G)(\text{Ho } K)$ also constitute a conjugate pair.

- (iii) Let $F' \dashv G' : \mathcal{D}' \rightarrow \mathcal{D}$ be another adjunction, with unit η' and counit ε' . If

- (C°, Q, p) is a left deformation retract for F ,
- (C'°, Q', p') is a left deformation retract for F' ,
- F sends objects in C° to objects in C'° ,
- (D°, R, i) is a right deformation retract for G ,

[1] Recall [definition A.1.3](#).

- $(\mathcal{D}'^\circ, R', i')$ is a right deformation retract for G' , and
- G' sends objects in \mathcal{D}'° to objects in \mathcal{D}° ,

then the three derived adjunctions

$$\begin{aligned} \mathbf{L}F \dashv \mathbf{R}G &: \text{Ho } \mathcal{D} \rightarrow \text{Ho } \mathcal{C} \\ \mathbf{L}F' \dashv \mathbf{R}G' &: \text{Ho } \mathcal{D}' \rightarrow \text{Ho } \mathcal{D} \\ \mathbf{L}(F'F) \dashv \mathbf{R}(GG') &: \text{Ho } \mathcal{D}' \rightarrow \text{Ho } \mathcal{C} \end{aligned}$$

are compatible with the comparison isomorphisms $\mu_{F',F}$ and $\delta_{G,G'}$, i.e.

$$\begin{aligned} \bar{\eta}'' &= (\delta_{G,G'}^{-1} \circ \mu_{F',F}) \bullet (\mathbf{R}G)\bar{\eta}'(\mathbf{L}F) \bullet \bar{\eta} \\ \bar{\varepsilon}'' &= \bar{\varepsilon}' \bullet (\mathbf{L}F')\bar{\varepsilon}(\mathbf{R}G') \bullet (\mu_{F',F}^{-1} \circ \delta_{G,G'}) \end{aligned}$$

where $\bar{\eta}''$ and $\bar{\varepsilon}''$ are the unit and counit for $\mathbf{L}(F'F) \dashv \mathbf{R}(GG')$ as constructed in claim (i), and \bullet denotes the horizontal composition^[2] of natural transformations.

Proof. (i). We will check that the triangle identities hold for the announced unit and counit. First, observe that these diagrams commute:

$$\begin{array}{ccccc} FQ & \xrightarrow{FQ\eta} & FQGF & \xrightarrow{FQG iF} & FQGRF \\ Fp \downarrow & & FpGF \downarrow & & FpGRF \downarrow \\ F & \xrightarrow{F\eta} & FGF & \xrightarrow{FG iF} & FGRF \\ & \searrow \text{id} & \varepsilon F \downarrow & & \varepsilon RF \downarrow \\ & & F & \xrightarrow{iF} & RF \end{array}$$

$$\begin{array}{ccccc} QG & \xrightarrow{\eta QG} & GFQG & \xrightarrow{GiFQG} & GRFQG \\ pG \downarrow & & GFpG \downarrow & & GRFpG \downarrow \\ G & \xrightarrow{\eta G} & GFG & \xrightarrow{GiFG} & GRFG \\ & \searrow \text{id} & G\varepsilon \downarrow & & GR\varepsilon \downarrow \\ & & G & \xrightarrow{Gi} & GR \end{array}$$

[2] — also known as the Godement product.

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Thus, we have the following equations:

$$\begin{aligned}
\bar{\varepsilon}(\mathbf{L}F) \cdot (\mathbf{L}F)\bar{\eta} &= \text{Ho}(iFQ)^{-1} \cdot \text{Ho}(\varepsilon RFQ \cdot FpGRFQ) \\
&\quad \cdot \text{Ho}(FQGiFQ \cdot FQ\eta Q) \cdot \text{Ho}(FQp)^{-1} \\
&= \text{Ho}(iFQ)^{-1} \cdot \text{Ho}(iFQ \cdot FpQ) \cdot \text{Ho}(FQp)^{-1} \\
&= \text{Ho}(FpQ) \cdot \text{Ho}(FQp)^{-1}
\end{aligned}$$

$$\begin{aligned}
(\mathbf{R}F)\bar{\varepsilon} \cdot \bar{\eta}(\mathbf{R}G) &= \text{Ho}(GRi)^{-1} \cdot \text{Ho}(GR\varepsilon R \cdot GRFpGR) \\
&\quad \cdot \text{Ho}(GiFQGR \cdot \eta QGR) \cdot \text{Ho}(pGR)^{-1} \\
&= \text{Ho}(GRi)^{-1} \cdot \text{Ho}(GiR \cdot pGR) \cdot \text{Ho}(pGR)^{-1} \\
&= \text{Ho}(GRi)^{-1} \cdot \text{Ho}(GiR)
\end{aligned}$$

We must now show that

$$\text{Ho}(FpQ) \cdot \text{Ho}(FQp)^{-1} = \text{id}_{\text{Ho}(FQ)}$$

and

$$\text{Ho}(GRi)^{-1} \cdot \text{Ho}(GiR) = \text{id}_{\text{Ho}(GR)}$$

but those equations hold because the diagrams of natural *weak equivalences* below commute:

$$\begin{array}{ccccccc}
FQ & \xlongequal{\quad} & FQ & \xleftarrow{FQp} & FQQ & \xrightarrow{FpQ} & FQ & \xlongequal{\quad} & FQ \\
\parallel & & \uparrow FQp & & \uparrow FQQp & & \uparrow FQp & & \parallel \\
FQ & \xleftarrow{FQp} & FQQ & \xleftarrow{FQpQ} & FQQQ & \xrightarrow{FpQQ} & FQQ & \xrightarrow{FQp} & FQ \\
\parallel & & \parallel & & \parallel & & \parallel & & \parallel \\
FQ & \xleftarrow{FQp} & FQQ & \xleftarrow{FQpQ} & FQQQ & \xrightarrow{FQpQ} & FQQ & \xrightarrow{FQp} & FQ
\end{array}$$

$$\begin{array}{ccccccc}
GR & \xlongequal{\quad} & GR & \xrightarrow{GiR} & GRR & \xleftarrow{GRi} & GR & \xlongequal{\quad} & GR \\
\parallel & & \downarrow GRi & & \downarrow GRRi & & \downarrow GRi & & \parallel \\
GR & \xrightarrow{GRi} & GRR & \xrightarrow{GiRR} & GRRR & \xleftarrow{GRiR} & GRR & \xleftarrow{GRi} & GR \\
\parallel & & \parallel & & \parallel & & \parallel & & \parallel \\
GR & \xrightarrow{GRi} & GRR & \xrightarrow{GRiR} & GRRR & \xleftarrow{GRiR} & GRR & \xleftarrow{GRi} & GR
\end{array}$$

(Recall [theorem A.4.21](#).)

(ii). We use the following explicit formulae for $\mathbf{L}\varphi$ and $\mathbf{R}\psi$:

$$\begin{aligned}\mathbf{L}\varphi &= \text{Ho}(\varphi Q') \cdot \text{Ho}(FpHQ') \cdot \text{Ho}(FQH p')^{-1} \\ \mathbf{R}\psi &= \text{Ho}(GRKi')^{-1} \cdot \text{Ho}(GiKR') \cdot \text{Ho}(\psi R')\end{aligned}$$

We wish to show that these equations hold:

$$\begin{aligned}(1) \quad & \bar{\varepsilon}(\text{Ho } K) \cdot (\mathbf{L}F)(\mathbf{R}\psi) = (\text{Ho } K)\bar{\varepsilon}' \cdot (\mathbf{L}\varphi)(\mathbf{R}G') \\ (2) \quad & (\mathbf{R}G)(\mathbf{L}\varphi) \cdot \bar{\eta}(\text{Ho } H) = (\mathbf{R}\psi)(\mathbf{L}F') \cdot (\text{Ho } H)\bar{\eta}'\end{aligned}$$

Observe that the following diagrams commute,

$$\begin{array}{ccccccc} & & FQGRK & \xrightarrow{FpGRK} & FGRK & \xrightarrow{\varepsilon RK} & RK \xleftarrow{iK} K \\ & & \downarrow FQGRKi' & & & & \downarrow RKi' \quad \downarrow Ki' \\ & & FQGRKR' & \xrightarrow{FpGRKR'} & FGRKR' & \xrightarrow{\varepsilon RKR'} & RKR' \xleftarrow{iKR'} KR' \\ & & \uparrow FQGiKR' & & & & \uparrow iKR' \quad \parallel \\ FQHG'R' & \xrightarrow{FQ\psi R'} & FQGKR' & \xrightarrow{FpGKR'} & FGKR' & \xrightarrow{\varepsilon KR'} & KR' = KR' \\ \parallel & & & & \parallel & & \\ FQHG'R' & \xrightarrow{FpHG'R'} & FHG'R' & \xrightarrow{F\psi R'} & FGKR' & & \\ \\ & & FHQ'G'R' & \xrightarrow{\varphi Q'G'R'} & KF'Q'G'R' & \xrightarrow{KF'p'G'R'} & KF'G'R' \\ & & \parallel & & & & \parallel \\ & & FHQ'G'R' & \xrightarrow{FHp'G'R'} & FHG'R' & \xrightarrow{\varphi G'R'} & KF'G'R' \xrightarrow{K\varepsilon'R'} KR' \xleftarrow{Ki'} K \\ & & \uparrow FpHQ'G'R' & & \uparrow FpHG'R' & & \\ & & FQHQ'G'R' & \xrightarrow{FQH p'G'R'} & FQHG'R' & & \\ & & \downarrow FQH p'G'R' & & \parallel & & \\ & & FQHG'R' & = & FQHG'R' & & \end{array}$$

and so we have the identities shown below:

$$\begin{aligned}\bar{\varepsilon}(\text{Ho } K) \cdot (\mathbf{L}F)(\mathbf{R}\psi) &= \text{Ho}(Ki')^{-1} \cdot \text{Ho}(\varepsilon KR' \cdot F\psi R' \cdot FpHG'R') \\ (\text{Ho } K)\bar{\varepsilon}' \cdot (\mathbf{L}\varphi)(\mathbf{R}G') &= \text{Ho}(Ki')^{-1} \cdot \text{Ho}(K\varepsilon'R' \cdot \varphi G'R' \cdot FpHG'R')\end{aligned}$$

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Since $\varepsilon K \bullet F\psi = K\varepsilon' \bullet \varphi G'$, we conclude that equation (1) holds. The dual calculation proves equation (2).

(iii). Recall that the comparison isomorphisms have the following explicit forms:

$$\mu_{F',F} = \text{Ho}(F' p' FQ) \qquad \delta_{G,G'} = \text{Ho}(GiG' R')$$

Thus, $(\delta_{G,G'}^{-1} \circ \mu_{F',F}) \bullet (\mathbf{R}G)\bar{\eta}'(\mathbf{L}F) \bullet \bar{\eta}$ expands to

$$\begin{aligned} & \text{Ho}(GiG' R' F' FQ)^{-1} \bullet \text{Ho}(GRG' R' F' p' FQ) \\ & \bullet \text{Ho}(GRG' i' F' Q' FQ \bullet GR\eta Q' FQ) \bullet \text{Ho}(GRp' FQ)^{-1} \\ & \bullet \text{Ho}(GiFQ \bullet \eta Q) \bullet (\text{Ho } p)^{-1} \end{aligned}$$

and a straightforward calculation then shows

$$(\delta_{G,G'}^{-1} \circ \mu_{F',F}) \bullet (\mathbf{R}G)\bar{\eta}'(\mathbf{L}F) \bullet \bar{\eta} = \text{Ho}(GG' i' F' FQ \bullet G\eta FQ \bullet \eta Q) \bullet (\text{Ho } p)^{-1}$$

but the RHS is precisely the definition of $\bar{\eta}''$. The dual calculation proves the other equation. \blacksquare

Proposition 2.3.12. *Let \mathcal{C} and \mathcal{D} be two relative categories, let $F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$ be an adjunction of ordinary categories with unit η and counit ε , let $(\mathcal{C}^\circ, Q, p)$ be a left deformation retract for F , and let $(\mathcal{D}^\circ, R, i)$ be a right deformation retract for G . Consider the following statements:*

- (i) *For all objects \tilde{X} in \mathcal{C}° and all objects \hat{B} in \mathcal{D}° , if $F\tilde{X} \rightarrow \hat{B}$ is a weak equivalence in \mathcal{D} , then its right adjoint transpose $\tilde{X} \rightarrow G\hat{B}$ is a weak equivalence in \mathcal{C} .*
- (ii) *The natural transformation $GiFQ \bullet \eta Q : Q \Rightarrow GRFQ$ is a natural weak equivalence.*
- (iii) *The derived adjunction unit $\bar{\eta} : \text{id}_{\text{Ho } \mathcal{C}} \Rightarrow (\mathbf{R}G)(\mathbf{L}F)$ is a natural isomorphism.*
- (i') *For all objects \tilde{X} in \mathcal{C}° and all objects \hat{B} in \mathcal{D}° , if $\tilde{X} \rightarrow G\hat{B}$ is a weak equivalence in \mathcal{C} , then its left adjoint transpose $F\tilde{X} \rightarrow \hat{B}$ is a weak equivalence in \mathcal{D} .*
- (ii') *The natural transformation $\varepsilon R \bullet FpGR : FQGR \Rightarrow R$ is a natural weak equivalence.*

(iii') *The derived adjunction counit $\bar{\epsilon} : (\mathbf{LF})(\mathbf{RG}) \Rightarrow \text{id}_{\text{Ho } \mathcal{D}}$ is a natural isomorphism.*

We have the implications (i) \Rightarrow (ii) \Rightarrow (iii); if $\text{weq } \mathcal{C}$ has the 2-out-of-3 property, then (ii) \Rightarrow (i); and if \mathcal{C} is a saturated homotopical category, then (iii) \Rightarrow (ii). Dually, (i') \Rightarrow (ii') \Rightarrow (iii'); if $\text{weq } \mathcal{D}$ has the 2-out-of-3 property, then (ii') \Rightarrow (i'); and if \mathcal{D} is a saturated homotopical category, then (iii') \Rightarrow (ii').

Proof. (i) \Rightarrow (ii). We have a natural weak equivalence $iFQ : FQ \Rightarrow RFQ$, so, by the hypothesis, its right adjoint transpose $GiFQ \bullet \eta Q$ is also a natural weak equivalence.

(ii) \Rightarrow (iii). The derived adjunction unit is given by $\bar{\eta} = \text{Ho}(GiFQ \bullet \eta Q) \bullet (\text{Ho } p)^{-1}$, which is certainly a natural isomorphism if $GiFQ \bullet \eta Q$ is a natural weak equivalence.

(ii) \Rightarrow (i). Assume $\text{weq } \mathcal{C}$ has the 2-out-of-3 property. Given \tilde{X} in \mathcal{C}° , the diagram below commutes,

$$\begin{array}{ccccc} Q\tilde{X} & \xrightarrow{\eta_{Q\tilde{X}}} & GFQ\tilde{X} & \xrightarrow{Gi_{FQ\tilde{X}}} & GRFQ\tilde{X} \\ p_{\tilde{X}} \downarrow & & & & \downarrow GRFp_{\tilde{X}} \\ \tilde{X} & \xrightarrow{\eta_{\tilde{X}}} & GF\tilde{X} & \xrightarrow{Gi_{F\tilde{X}}} & GRF\tilde{X} \end{array}$$

but the top row and the two vertical arrows are weak equivalences in \mathcal{C} , so the bottom row must be a weak equivalence as well, by the 2-out-of-3 property.

Let $g : F\tilde{X} \rightarrow \hat{B}$ be a weak equivalence in \mathcal{D} , and let $f = Gg \circ \eta_{\tilde{X}}$ be its right adjoint transpose in \mathcal{C} . We know $GR : \mathcal{D} \rightarrow \mathcal{C}$ is a relative functor, so $GRg : GRF\tilde{X} \rightarrow GR\hat{B}$ is a weak equivalence in \mathcal{C} ; but

$$Gi_{\hat{B}} \circ f = Gi_{\hat{B}} \circ Gg \circ \eta_{\tilde{X}} = GRg \circ (Gi_{F\tilde{X}} \circ \eta_{\tilde{X}})$$

and we know $Gi_{\hat{B}} : G\hat{B} \rightarrow GR\hat{B}$ is a weak equivalence in \mathcal{C} , so by the 2-out-of-3 property again, f must be a weak equivalence in \mathcal{C} .

(iii) \Rightarrow (ii). Now assume \mathcal{C} is a saturated homotopical category. If $\bar{\eta}$ is a natural isomorphism, then $\text{Ho}(GiFQ \bullet \eta Q)$ must also be a natural isomorphism, and so $GiFQ \bullet \eta Q$ is a natural weak equivalence, by the saturation hypothesis. \blacksquare

Corollary 2.3.13. *With notation as above, suppose the **Quillen equivalence condition** is satisfied:*

- *For all objects \tilde{X} in \mathcal{C}° and all objects \hat{B} in \mathcal{D}° , a morphism $F\tilde{X} \rightarrow \hat{B}$ is a weak equivalence in \mathcal{D} if and only if its right adjoint transpose $\tilde{X} \rightarrow G\hat{B}$ is a weak equivalence in \mathcal{C} .*

Then the derived adjunction is an adjoint equivalence of categories. ■

2.4 DHKS derived functors

Prerequisites. §§ 2.1, 2.2, 2.3.

Notice that in [theorem 2.3.6](#), we constructed derived functors by applying the 2-functor $\text{Ho} : \mathfrak{RelCat} \rightarrow \mathfrak{Cat}$ to a relative functor that is naturally weakly equivalent to the original functor, and then we showed that its homotopy version has a universal property. This suggests that the intermediate relative functor might itself have a homotopical universal property.

In this section we follow [[DHKS](#), Ch. VII].

Definition 2.4.1. Let \mathcal{C} and \mathcal{D} be homotopical categories. A **homotopical left approximation** for an ordinary functor $F : \mathcal{C} \rightarrow \mathcal{D}$ is a homotopical right (!) Kan extension of F along $\text{id}_{\mathcal{C}}$. Dually, a **homotopical right approximation** for an ordinary functor $G : \mathcal{D} \rightarrow \mathcal{C}$ is a homotopical left (!) Kan extension of G along $\text{id}_{\mathcal{D}}$.

REMARK 2.4.2. More explicitly, a homotopical left approximation for $F : \mathcal{C} \rightarrow \mathcal{D}$ is a homotopically terminal object in the homotopical category $([\mathcal{C}, \mathcal{D}]_h \downarrow F)_h$ described below:

- The objects are pairs (K, α) where K is a homotopical functor $\mathcal{C} \rightarrow \mathcal{D}$ and α is a natural transformation of type $K \Rightarrow F$.
- The morphisms $(K', \alpha') \rightarrow (K, \alpha)$ are those natural transformations $\psi : K' \Rightarrow K$ such that $\alpha \bullet \psi = \alpha'$.
- The weak equivalences are the natural weak equivalences.

Dually, a homotopical right approximation for $G : \mathcal{D} \rightarrow \mathcal{C}$ is a homotopically initial object in the homotopical category $(F \downarrow [\mathcal{D}, \mathcal{C}]_h)_h$. By [corollary 2.2.14](#), homotopical approximations are homotopically unique.

We have the following special case:

Proposition 2.4.3. *Let \mathcal{Q} be a homotopical endofunctor on a homotopical category \mathcal{C} and let $p : \mathcal{Q} \Rightarrow \text{id}_{\mathcal{C}}$ be a natural transformation. The following are equivalent:*

- (i) (\mathcal{Q}, p) is a homotopical left approximation for $\text{id}_{\mathcal{C}}$.
- (ii) $(\mathcal{C}, \mathcal{C}, \mathcal{Q}, p)$ is a left deformation retract for $\text{id}_{\mathcal{C}}$.

Dually, let \mathcal{R} be a homotopical endofunctor on a homotopical category \mathcal{D} , and let $i : \text{id}_{\mathcal{D}} \Rightarrow \mathcal{R}$ be a natural transformation. The following are equivalent:

- (i') (\mathcal{R}, i) is a homotopical right approximation for $\text{id}_{\mathcal{D}}$.
- (ii') $(\mathcal{D}, \mathcal{D}, \mathcal{R}, i)$ is a right deformation retract for $\text{id}_{\mathcal{D}}$.

Proof. (i) \Rightarrow (ii). If (\mathcal{Q}, p) is a homotopical left approximation for $\text{id}_{\mathcal{C}}$, then there must exist a commutative diagram of the form below,

$$\begin{array}{ccccccc}
 \text{id}_{\mathcal{C}} & \longrightarrow & \mathcal{Q}_1 & \longleftarrow & \mathcal{Q}_2 & \longrightarrow & \cdots & \longleftarrow & \mathcal{Q} \\
 \text{id} \downarrow & & p_1 \downarrow & & p_2 \downarrow & & & & p \downarrow \\
 \text{id}_{\mathcal{C}} & \longlongequal{\quad} & \text{id}_{\mathcal{C}} & \longlongequal{\quad} & \text{id}_{\mathcal{C}} & \longlongequal{\quad} & \cdots & \longlongequal{\quad} & \text{id}_{\mathcal{C}}
 \end{array}$$

where all the arrows in the top row are natural weak equivalences. Using 2-out-of-3 property, we deduce (by induction) that p_1, p_2, \dots, p are also natural weak equivalences; thus $(\mathcal{C}, \mathcal{C}, \mathcal{Q}, p)$ is indeed a left deformation retract for $\text{id}_{\mathcal{C}}$.

(ii) \Rightarrow (i). If $(\mathcal{C}, \mathcal{C}, \mathcal{Q}, p)$ is a left deformation retract for $\text{id}_{\mathcal{C}}$, then $p : \mathcal{Q} \Rightarrow \text{id}_{\mathcal{C}}$ is a natural weak equivalence; but $(\text{id}_{\mathcal{C}}, \text{id}_{\text{id}_{\mathcal{C}}})$ is a terminal object in $([\mathcal{C}, \mathcal{C}]_{\text{h}} \downarrow \text{id}_{\mathcal{C}})_{\text{h}}$, so by [proposition 2.2.2](#), (\mathcal{Q}, p) must be a homotopically terminal object. \blacksquare

Definition 2.4.4. Let $F, F' : \mathcal{C} \rightarrow \mathcal{D}$ be ordinary functors between homotopical categories, and let $\varphi : F \Rightarrow F'$ be a natural transformation. We define the homotopical category $([\min 2, [\mathcal{C}, \mathcal{D}]_{\text{h}}]_{\text{h}} \downarrow \varphi)_{\text{h}}$ as follows:

- The objects are tuples $(H, H', \alpha, \alpha', \theta)$ where H and H' are homotopical functors $\mathcal{C} \rightarrow \mathcal{D}$, α and α' are natural transformations of type $H \Rightarrow F$ and $H' \Rightarrow F'$ (respectively), and $\theta : H \Rightarrow H'$ is a natural transformation such that $\varphi \bullet \alpha = \alpha' \bullet \theta$.

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- The morphisms $(H, H', \alpha, \alpha', \theta) \rightarrow (K, K', \beta, \beta', \chi)$ are pairs (ζ, ζ') of natural transformations, where $\zeta : H \Rightarrow K$ and $\zeta' : H' \Rightarrow K'$, such that $\chi \cdot \zeta = \zeta' \cdot \theta$, $\beta \cdot \zeta = \alpha$, and $\beta' \cdot \zeta' = \alpha'$.
- The weak equivalences are those (ζ, ζ') where both ζ and ζ' are natural weak equivalences.

A **homotopical left approximation** for φ is a homotopically terminal object $(\mathbb{L}F, \mathbb{L}F', \delta, \delta', \mathbb{L}\varphi)$ in $([\min \mathbb{2}, [C, \mathcal{D}]_h]_h \downarrow \varphi)_h$ such that $(\mathbb{L}F, \delta)$ is a homotopical left approximation for F and $(\mathbb{L}F', \delta')$ is a homotopical left approximation for F' .

Dually, let $G, G' : \mathcal{D} \rightarrow \mathcal{C}$ be ordinary functors between homotopical categories, and let $\psi : G' \Rightarrow G$ be a natural transformation. We define the homotopical category $(\psi \downarrow [\min \mathbb{2}, [D, C]_h]_h)$ as follows:

- The objects are tuples $(H, H', \alpha, \alpha', \theta)$ where H and H' are homotopical functors $\mathcal{D} \rightarrow \mathcal{C}$, α and α' are natural transformations of type $G \Rightarrow H$ and $G' \Rightarrow H'$ (respectively), and $\theta : H' \Rightarrow H$ is a natural transformation such that $\alpha \cdot \psi = \theta \cdot \alpha'$.
- The morphisms $(K, K', \beta, \beta', \chi) \rightarrow (H, H', \alpha, \alpha', \theta)$ are pairs (ζ, ζ') of natural transformations, where $\zeta : K \Rightarrow H$ and $\zeta' : K' \Rightarrow H'$, such that $\zeta \cdot \chi = \theta \cdot \zeta'$, $\zeta \cdot \beta = \alpha$, and $\zeta' \cdot \beta' = \alpha'$.
- The weak equivalences are those (ζ, ζ') where both ζ and ζ' are natural weak equivalences.

A **homotopical right approximation** for ψ is a homotopically initial object $(\mathbb{R}G, \mathbb{R}G', \delta, \delta', \mathbb{R}\psi)$ in $(\psi \downarrow [\min \mathbb{2}, [D, C]_h]_h)_h$ such that $(\mathbb{R}G, \delta)$ is a homotopical right approximation for G and $(\mathbb{R}G', \delta')$ is a homotopical right approximation for G' .

Theorem 2.4.5. *Let \mathcal{C} and \mathcal{D} be homotopical categories.*

- (i) *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be an ordinary functor. If (C°, Q, p) is a left deformation retract for F , then (FQ, Fp) is a homotopical absolute right Kan extension of F along $\text{id}_{\mathcal{C}}$.*

- (ii) Let $F, F' : \mathcal{C} \rightarrow \mathcal{D}$ be a parallel pair of ordinary functors. If $(\mathcal{C}^\circ, \mathcal{Q}, p)$ is a left deformation retract for both F and F' , then for any natural transformation $\varphi : F \Rightarrow F'$, $(F\mathcal{Q}, F'\mathcal{Q}, Fp, F'p, \varphi\mathcal{Q})$ is a homotopical left approximation for φ .
- (iii) Let $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{D} \rightarrow \mathcal{E}$ be ordinary functors between homotopical categories. If $(\mathcal{C}^\circ, \mathcal{Q}^{\mathcal{C}}, p^{\mathcal{C}})$ is a left deformation retract for F , $(\mathcal{D}^\circ, \mathcal{Q}^{\mathcal{D}}, p^{\mathcal{D}})$ is a left deformation retract for G , and F maps objects in \mathcal{C}° to objects in \mathcal{D}° , then, for any homotopical left approximation $((\mathbb{L}F), \delta^F)$ for F and any homotopical left approximation $((\mathbb{L}G), \delta^G)$ for G , $((\mathbb{L}G)(\mathbb{L}F), \delta^G \circ \delta^F)$ is a homotopical left approximation for GF .

Dually:

- (i') Let $G : \mathcal{D} \rightarrow \mathcal{C}$ be an ordinary functor. If $(\mathcal{D}^\circ, \mathcal{R}, i)$ is a right deformation retract for F , then (GR, Gi) is a homotopical absolute left Kan extension of G along $\text{id}_{\mathcal{D}}$.
- (ii') Let $G, G' : \mathcal{D} \rightarrow \mathcal{C}$ be a parallel pair of ordinary functors. If $(\mathcal{D}^\circ, \mathcal{R}, i)$ is a right deformation retract for both G and G' , then for any natural transformation $\psi : G' \Rightarrow G$, $(GR, G'R, Gi, G'i, \psi\mathcal{R})$ is a homotopical right approximation for ψ .
- (iii') Let $F : \mathcal{C} \rightarrow \mathcal{B}$ and $G : \mathcal{D} \rightarrow \mathcal{C}$ be ordinary functors between homotopical categories. If $(\mathcal{C}^\circ, \mathcal{R}^{\mathcal{C}}, i^{\mathcal{C}})$ is a right deformation retract for F , $(\mathcal{D}^\circ, \mathcal{R}^{\mathcal{D}}, i^{\mathcal{D}})$ is a right deformation retract for G , and G maps objects in \mathcal{D}° to objects in \mathcal{C}° , then, for any homotopical right approximation $((\mathbb{R}F), \delta^F)$ for F and any homotopical right approximation $((\mathbb{R}G), \delta^G)$ for G , $((\mathbb{R}F)(\mathbb{R}G), \delta^F \circ \delta^G)$ is a homotopical right approximation for FG .

Proof. (i). Let $H : \mathcal{D} \rightarrow \mathcal{E}$ and $K : \mathcal{C} \rightarrow \mathcal{E}$ be any two homotopical functors, and let $\alpha : K \Rightarrow HF$ be any natural transformation. Then, we have the following commutative diagram of natural transformations,

$$\begin{array}{ccccc}
 K & \xleftarrow{Kp} & KQ & \xrightarrow{\alpha Q} & HFQ \\
 & \searrow \alpha & \downarrow & \swarrow HFp & \\
 & & HF & &
 \end{array}$$

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and, for any other homotopical functor $K' : \mathcal{C} \rightarrow \mathcal{E}$ and natural transformation $\psi : K' \Rightarrow K$, for $\alpha' = \alpha \bullet \psi$, the diagram

$$\begin{array}{ccccc}
 K' & \xleftarrow{K'p} & K'Q & \xrightarrow{\alpha'Q} & HFQ \\
 \psi \downarrow & & \downarrow \psi Q & & \parallel \\
 K & \xleftarrow{Kp} & KQ & \xrightarrow{\alpha Q} & HFQ \\
 & \searrow \alpha & \downarrow & \swarrow HFp & \\
 & & HF & &
 \end{array}$$

also commutes; thus, (HFQ, HFp) is indeed a homotopically terminal object in $([C, \mathcal{E}]_h \downarrow HF)_h$.

(ii). Suppose $(H, H', \alpha, \alpha', \theta)$ is an object in $([\min 2, [C, D]_h]_h \downarrow \varphi)_h$. The diagram below commutes,

$$\begin{array}{ccccc}
 & & HQ & \xrightarrow{\theta Q} & H'Q \\
 & \swarrow H_p & \downarrow & \swarrow H'_p & \downarrow \alpha'Q \\
 H & \xrightarrow{\theta} & H' & & \\
 \alpha \downarrow & & \downarrow \alpha Q & & \downarrow \alpha' \\
 & & FQ & \xrightarrow{\varphi Q} & F'Q \\
 & \swarrow F_p & \downarrow & \swarrow F'_p & \\
 F & \xrightarrow{\varphi} & F' & &
 \end{array}$$

and (H_p, H'_p) is a weak equivalence, so $(FQ, F'Q, F_p, F'_p, \varphi Q)$ is indeed a homotopically terminal object in $([\min 2, [C, D]_h]_h \downarrow \varphi)_h$.

(iii). To begin, observe that $Gp^{D^\circ} FQ^{C^\circ} : GQ^{D^\circ} FQ^{C^\circ} \Rightarrow GFQ^{C^\circ}$ is a natural weak equivalence; and, as established above, both $\delta^F Q^{C^\circ} : (\mathbb{L}F)Q^{C^\circ} \Rightarrow FQ^{C^\circ}$ and $\delta^G Q^{D^\circ} : (\mathbb{L}G)Q^{D^\circ} \Rightarrow GQ^{D^\circ}$ are natural weak equivalences, so their horizontal composite $(\delta^G Q^{D^\circ}) \circ (\delta^F Q^{C^\circ})$ is also a natural weak equivalence. We also know that $(C^\circ, Q^{C^\circ}, p^{C^\circ})$ is a left deformation retract for GF , so $(GFQ^{C^\circ}, GFp^{C^\circ})$ is a homotopical left approximation for GF . Now, noting that the following diagram commutes,

$$\begin{array}{ccccc}
 (\mathbb{L}G)Q^{D^\circ} (\mathbb{L}F)Q^{C^\circ} & \xrightarrow{(\delta^G Q^{D^\circ}) \circ (\delta^F Q^{C^\circ})} & GQ^{D^\circ} FQ^{C^\circ} & \xrightarrow{Gp^{D^\circ} FQ^{C^\circ}} & GFQ^{C^\circ} \\
 \downarrow ((\mathbb{L}G)p^{D^\circ}) \circ ((\mathbb{L}F)p^{C^\circ}) & & \downarrow (Gp^{D^\circ}) \circ (Fp^{C^\circ}) & & \downarrow GFp^{C^\circ} \\
 (\mathbb{L}G)(\mathbb{L}F) & \xrightarrow{\delta^G \circ \delta^F} & GF & \xlongequal{\quad} & GF
 \end{array}$$

we conclude that $((\mathbb{L}G)(\mathbb{L}F), \delta^G \circ \delta^F)$ and (GFQ^{c^o}, GFp^{c^o}) are weakly equivalent in $([C, \mathcal{E}]_h \downarrow GF)_h$, and so $((\mathbb{L}G)(\mathbb{L}F), \delta^G \circ \delta^F)$ is also a homotopical left approximation for GF , by [proposition 2.2.2](#). \blacksquare

REMARK 2.4.6. Unfortunately, the assignment $F \mapsto FQ$ (resp. $G \mapsto GR$) does not extend to lax (resp. oplax) 2-functors, because we do not have a natural transformation $\text{id}_C \Rightarrow Q$ (resp. $R \Rightarrow \text{id}_D$).

Corollary 2.4.7. *Let C and D be homotopical categories, and let $\gamma_C : C \rightarrow \text{Ho } C$ and $\gamma_D : D \rightarrow \text{Ho } D$ be the respective localising functors.*

- *If $F : C \rightarrow D$ is a left deformable functor and $(\mathbb{L}F, \delta)$ is any homotopical left approximation for F , then $(\text{Ho}(\mathbb{L}F), \gamma_D \delta)$ is a total left derived functor for F .*
- *If $G : D \rightarrow C$ is a right deformable functor and $(\mathbb{R}G, \delta)$ is any homotopical right approximation for G , then $(\text{Ho}(\mathbb{R}G), \gamma_C \delta)$ is a total right derived functor for G .*

Proof. Combine theorems [2.3.6](#) and [2.4.5](#). \blacksquare

2.5 Three-arrow calculi

Prerequisites. §§ [2.1](#), [A.4](#).

In this section, we follow [\[DHKS, § 36\]](#) and [\[Thomas, 2011\]](#).

Definition 2.5.1. Let C be a relative category, let $\mathcal{W} = \text{weq } C$ be the subcategory of weak equivalences in C , and let \mathcal{U} and \mathcal{V} be subcategories of \mathcal{W} . We say C **admits a three-arrow calculus** for C with respect to $(\mathcal{U}, \mathcal{V})$ if the following conditions are satisfied:

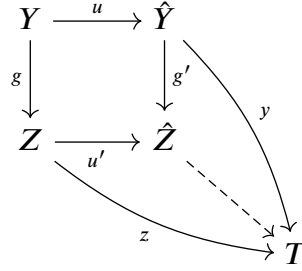
A1. For each weak equivalence w in C , there exist u in \mathcal{U} and v in \mathcal{V} such that $w = v \circ u$.

A2. Given a diagram of the form $\hat{Y} \xleftarrow{u} Y \xrightarrow{g} Z$ in C with u in \mathcal{U} , there exists a diagram of the form $\hat{Y} \xrightarrow{g'} \hat{Z} \xleftarrow{u'} Z$ such that

- $g' \circ u = u' \circ g$,
- u' is in \mathcal{U} , and

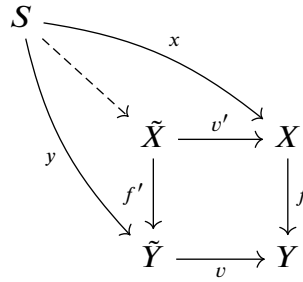
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- given any diagram of the form $\hat{Y} \xrightarrow{y} T \xleftarrow{z} Z$ such that $y \circ u = z \circ g$, there exists a (not necessarily unique) morphism $T \rightarrow \hat{Z}$ making the diagram below commute:



- A3.** Given a diagram of the form $X \xrightarrow{f} Y \xleftarrow{v} \tilde{Y}$ in \mathcal{C} with v in \mathcal{V} , there exists a diagram of the form $X \xleftarrow{v'} \tilde{X} \xrightarrow{f'} \tilde{Y}$ such that

- $f \circ v' = v \circ g'$,
- v' is in \mathcal{V} , and
- given any diagram of the form $X \xleftarrow{x} S \xrightarrow{y} Y$ such that $f \circ x = v \circ y$, there exists a (not necessarily unique) morphism $S \rightarrow \tilde{X}$ making the diagram below commute:



A **uni-fractionable category** is a relative category \mathcal{C} together with a pair of subcategories $(\mathcal{U}, \mathcal{V})$ such that $\text{weq } \mathcal{C}$ has the 2-out-of-3 property in \mathcal{C} and \mathcal{C} admits a three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$.

REMARK 2.5.2. Note that axiom A1 implies that $\text{ob } \mathcal{U} = \text{ob } \mathcal{V} = \text{ob } \mathcal{C}$; in particular, every identity morphism in \mathcal{C} is also in \mathcal{U} and \mathcal{V} .

REMARK 2.5.3. Consider diagrams of the following forms,



where u, u' are in \mathcal{U} and v, v' are in \mathcal{V} . Under the assumption that \mathcal{W} has the 2-out-of-3 property in \mathcal{C} , the morphism g is in \mathcal{W} if and only if g' is in \mathcal{W} , and the morphism f is in \mathcal{W} if and only if f' is in \mathcal{W} .

Definition 2.5.4. Let \mathcal{C} be a relative category, let $\mathcal{W} = \text{weq } \mathcal{C}$ be the subcategory of weak equivalences in \mathcal{C} , and let \mathcal{U} and \mathcal{V} be subcategories of \mathcal{W} . A **functorial three-arrow calculus** for \mathcal{C} with respect to $(\mathcal{U}, \mathcal{V})$ consists of the following data:

- FA1.** A functorial factorisation system on \mathcal{W} with left class contained in $\text{mor } \mathcal{U}$ and right class contained in $\text{mor } \mathcal{V}$.
- FA2.** A functor from the full subcategory of $[\{\bullet \leftarrow \bullet \rightarrow \bullet\}, \mathcal{C}]$ spanned by those diagrams of the form $\hat{Y} \xleftarrow{u} Y \xrightarrow{g} Z$, where u is in \mathcal{U} , to the category $[\{\bullet \rightarrow \bullet \leftarrow \bullet\}, \mathcal{C}]$, such that each diagram $\hat{Y} \xleftarrow{u} Y \xrightarrow{g} Z$ is sent to a diagram of the form $\hat{Y} \xrightarrow{g'} \hat{Z} \xleftarrow{u'} Z$, where $g' \circ u = u' \circ g$, u' is in \mathcal{U} , and u' is an isomorphism if u is.
- FA3.** A functor from the full subcategory of $[\{\bullet \rightarrow \bullet \leftarrow \bullet\}, \mathcal{C}]$ spanned by those diagrams of the form $X \xrightarrow{f} Y \xleftarrow{v} \tilde{Y}$, where v is in \mathcal{V} , to the category $[\{\bullet \leftarrow \bullet \rightarrow \bullet\}, \mathcal{C}]$, such that each diagram $X \xrightarrow{f} Y \xleftarrow{v} \tilde{Y}$ is sent to a diagram of the form $X \xleftarrow{v'} \tilde{X} \xrightarrow{f'} \tilde{Y}$, where $f \circ v' = v \circ g'$, v' is in \mathcal{V} , and v' is an isomorphism if v is.

If such data exist, then we say \mathcal{C} **admits a functorial three-arrow calculus** with respect to $(\mathcal{U}, \mathcal{V})$.

REMARK 2.5.5. If $\text{mor } \mathcal{U}$ is closed under pushout in \mathcal{C} , then we may take pushouts to construct datum FA2; similarly, if $\text{mor } \mathcal{V}$ is closed under pullback in \mathcal{C} , then we may take pullbacks to construct datum FA3.

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REMARK 2.5.6. A relative category C admits a (functorial) three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$ if and only if the opposite relative category C^{op} admits a (functorial) three-arrow calculus with respect to $(\mathcal{V}, \mathcal{U})$.

Proposition 2.5.7. *Let C be a relative category and let \mathcal{U} and \mathcal{V} be subcategories of $\mathcal{W} = \text{weq } C$ (itself considered as a subcategory of C). If C admits a functorial three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$, then C admits a three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$.*

Proof. Obviously, having datum FA1 implies axiom A1 is satisfied. Now suppose we have a commutative square of the form below in C ,

$$\begin{array}{ccc} Y & \xrightarrow{u} & \hat{Y} \\ g \downarrow & & \downarrow y \\ Z & \xrightarrow{z} & T \end{array}$$

where u is in \mathcal{U} . The datum FA2 then gives us the following commutative diagram,

$$\begin{array}{ccccc} Y & \xrightarrow{u} & \hat{Y} & & \\ g \downarrow & \searrow & \downarrow y & & \\ & & T & \xrightarrow{\quad} & T \\ & & \parallel & \downarrow g' & \downarrow w \\ Z & \xrightarrow{\quad} & \hat{Z} & \xrightarrow{u'} & \hat{T} \\ z \searrow & & \downarrow & \searrow & \\ & & T & \xrightarrow{w} & \hat{T} \end{array}$$

and $w : T \rightarrow \hat{T}$ is an isomorphism, thus, there exists a morphism $\hat{Z} \rightarrow \hat{Y}$ making the diagram below commute:

$$\begin{array}{ccc} Y & \xrightarrow{u} & \hat{Y} \\ g \downarrow & & \downarrow g' \\ Z & \xrightarrow{u'} & \hat{Z} \\ & \searrow z & \searrow y \\ & & T \end{array}$$

This shows that axiom A2 is satisfied, and the dual argument proves axiom A3. ■

Proposition 2.5.8. *Let \mathcal{A} and \mathcal{C} be relative categories. If \mathcal{C} admits a functorial three-arrow calculus, and either*

- *weq \mathcal{C} has the 2-out-of-3 property in \mathcal{C} , or*
- *\mathcal{A} is a minimal relative category,*

then the relative functor category $[\mathcal{A}, \mathcal{C}]_{\text{h}}$ admits a functorial three-arrow calculus constructed componentwise from \mathcal{C} .

Proof. Let $(\mathcal{U}, \mathcal{V})$ be a functorial three-arrow calculus for \mathcal{C} . It is clear that, when \mathcal{A} is a minimal relative category, all the data constituting a three-arrow calculus for \mathcal{C} may be lifted componentwise to define a three-arrow calculus for $[\mathcal{A}, \mathcal{C}]_{\text{h}}$.

In general, we must check that $[\mathcal{A}, \mathcal{C}]_{\text{h}}$ is closed under the various componentwise constructions. However, if $f : A \rightarrow B$ is a weak equivalence in \mathcal{A} and $\theta : X \Rightarrow Y$ is a natural weak equivalence of relative functors $X, Y : \mathcal{A} \rightarrow \mathcal{M}$, and $\psi \bullet \varphi$ is the componentwise $(\mathcal{U}, \mathcal{V})$ -factorisation of θ , then the diagram below commutes,

$$\begin{array}{ccccc}
 XA & \xrightarrow{\varphi_A} & ZA & \xrightarrow{\psi_A} & YA \\
 Xf \downarrow & & Zf \downarrow & & Yf \downarrow \\
 XB & \xrightarrow{\varphi_B} & ZB & \xrightarrow{\psi_B} & YB
 \end{array}$$

and so by the 2-out-of-3 property of weq \mathcal{C} , Zf is also a weak equivalence in \mathcal{C} , thus $Z : \mathcal{A} \rightarrow \mathcal{M}$ is a relative functor. Similarly, one uses the 2-out-of-3 property of weq \mathcal{C} to ensure that the componentwise constructions satisfy the conditions to be data FA2 and FA3 for a functorial three-arrow calculus. ■

Theorem 2.5.9 (Fundamental theorem of three-arrow calculi). *Let \mathcal{C} be a relative category such that weq \mathcal{C} has the 2-out-of-3 property in \mathcal{C} . If \mathcal{C} admits a three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$, then:*

- (i) *Every morphism in $\text{Ho } \mathcal{C}$ can be represented by a zigzag in \mathcal{C} of the form below,*

$$X \xleftarrow{v} \tilde{X} \xrightarrow{f} \hat{Y} \xleftarrow{u} Y$$

where u is in \mathcal{U} and v is in \mathcal{V} .

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- (ii) Two such zigzags represent the same morphism in $\text{Ho } C$ if and only if there exists a commutative diagram in C of the form

$$\begin{array}{ccccccc}
 X & \xleftarrow{v} & \tilde{X} & \xrightarrow{f} & \hat{Y} & \xleftarrow{u} & Y \\
 \parallel & & \uparrow v_1 & & \uparrow v_2 & & \parallel \\
 X & \xleftarrow{v_3} & \bullet & \xrightarrow{f_3} & \bullet & \xleftarrow{u_3} & Y \\
 \parallel & & \downarrow w_1 & & \downarrow w_2 & & \parallel \\
 X & \xleftarrow{v_4} & \bullet & \xrightarrow{f_4} & \bullet & \xleftarrow{u_4} & Y \\
 \parallel & & \uparrow u_1 & & \uparrow u_2 & & \parallel \\
 X & \xleftarrow{v'} & \tilde{X}' & \xrightarrow{f'} & \hat{Y}' & \xleftarrow{u'} & Y
 \end{array}$$

where u_1, u_2, u_3, u_4 are in \mathcal{U} , v_1, v_2, v_3, v_4 are in \mathcal{V} , and w_1, w_2 are weak equivalences in C .

Proof. For the functorial case, see paragraph 36.3 in [DHKS]; for the general case, see Lemma 4.9 and Theorem 5.13 in [Thomas, 2011]. \square

Proposition 2.5.10. *If C is a homotopical category that admits a three-arrow calculus, then C is a saturated homotopical category.*

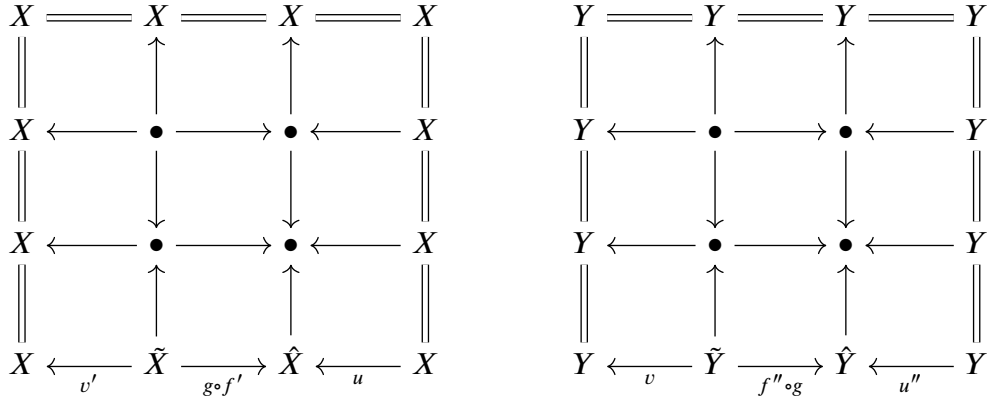
Proof. Suppose C admits a three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$. Let $f : X \rightarrow Y$ be a morphism in C whose image in $\text{Ho } C$ is an isomorphism, with inverse represented by the following zigzag,

$$Y \xleftarrow{v} \tilde{Y} \xrightarrow{g} \hat{X} \xleftarrow{u} X$$

where u is in \mathcal{U} and v is in \mathcal{V} . Then, by axioms A2 and A3, there exist v' in \mathcal{V} , f' in C , u'' in \mathcal{U} , and f'' in C such that the diagrams below commute,

$$\begin{array}{ccc}
 \tilde{X} & \xrightarrow{f'} & \tilde{Y} \\
 v' \downarrow & & \downarrow v \\
 X & \xrightarrow{f} & Y
 \end{array}
 \qquad
 \begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 u \downarrow & & \downarrow u'' \\
 \hat{X} & \xrightarrow{f''} & \hat{Y}
 \end{array}$$

and by theorem 2.5.9, we have commutative diagrams in C of the following form,



where all leftward- and upward-pointing arrows are weak equivalences in \mathcal{C} . We may then deduce that *every* arrow appearing in the above diagrams are in $\text{weq } \mathcal{C}$ by iteratively applying the 2-out-of-3 property of $\text{weq } \mathcal{C}$. In particular, $g \circ f'$ and $f'' \circ g$ are weak equivalences in \mathcal{C} , so the 2-out-of-6 property of $\text{weq } \mathcal{C}$ implies that f' , f'' , g are all in $\text{weq } \mathcal{C}$. We then conclude that f is in $\text{weq } \mathcal{C}$, by using the 2-out-of-3 property again. ■

MODEL CATEGORIES

3.1 Basics

Prerequisites. §§ 2.1, A.3.

In [1967], Quillen introduced the notion of a ‘closed model category’ (but we shall say simply ‘model category’) for homotopy theory, so as to formalise the similarities between the homotopy theory of spaces and homological algebra. The idea was that, to do homotopy theory, one really only needed to know which morphisms are cofibrations, which are weak equivalences, and which are fibrations.

Definition 3.1.1. A **model category** is a locally small category \mathcal{M} equipped with three subclasses $\mathcal{C}, \mathcal{W}, \mathcal{F}$ of $\text{mor } \mathcal{M}$ satisfying the following axioms:^[1]

CM1. \mathcal{M} has finite limits and finite colimits.

CM2. \mathcal{W} has the 2-out-of-3 property.

CM3. $\mathcal{C}, \mathcal{W},$ and \mathcal{F} are closed under retracts.

CM4. Given a commutative diagram

$$\begin{array}{ccc} A & \longrightarrow & X \\ \downarrow i & & \downarrow p \\ B & \longrightarrow & Y \end{array}$$

[1] This presentation is due to Quillen [1969].

where i is in C and p is in \mathcal{F} , if at least one of i or p is also in \mathcal{W} , then there exists a morphism $B \rightarrow X$ making both of the evident triangles commute.

CM5. Any morphism f in \mathcal{M} may be factored in two ways:

- $f = p \circ i$ where i is in $C \cap \mathcal{W}$ and p is in \mathcal{F} , and
- $f = q \circ j$, where j is in C and q is in $\mathcal{W} \cap \mathcal{F}$.

The triple $(C, \mathcal{W}, \mathcal{F})$ is said to be a **model structure** on \mathcal{M} . Given such a model structure on \mathcal{M} ,

- a **cofibration** is a morphism in C ,
- a **weak equivalence** is a morphism in \mathcal{W} ,
- a **fibration** is a morphism in \mathcal{F} ,
- a **trivial cofibration** (or **acyclic cofibration**) is a morphism in $C \cap \mathcal{W}$, and
- a **trivial fibration** (or **acyclic fibration**) is a morphism in $\mathcal{W} \cap \mathcal{F}$;
- a **cofibrant object** is an object X such that the unique morphism $0 \rightarrow X$ is a cofibration, and
- a **fibrant object** is an object X such that the unique morphism $X \rightarrow 1$ is a fibration.
- a **cofibrant–fibrant object** is an object that is both cofibrant and fibrant.

Definition 3.1.2. A **DHK model category** is a model category satisfying the following variants of CM1 and CM5:

CM1*. \mathcal{M} is complete and cocomplete.

CM5*. The $(C \cap \mathcal{W}, \mathcal{F})$ and $(C, \mathcal{W} \cap \mathcal{F})$ -factorisations can be chosen *functorially* in the sense of [definition A.3.21](#).

REMARK 3.1.3. Hovey [1999] and Hirschhorn [2003] attribute the stronger definition of ‘model category’ to Dwyer, Hirschhorn and Kan [DHK], hence the name ‘DHK model category’; of course, this is the definition used in the cited works, as well as in [DHKS]. Note also that the definition in [Hovey, 1999] includes the functorial factorisations as a *structure* instead of a property. On the other hand, [DS] and [GJ] use Quillen’s 1969 definition essentially verbatim.

Example 3.1.4. Let \mathcal{M} be any category with finite limits and finite colimits. The **trivial model structure** on \mathcal{M} is defined by the following data:

- The weak equivalences are the isomorphisms.
- Every morphism is both a cofibration and a fibration.

It is straightforward to directly verify that the axioms are satisfied in this case. Notice that if \mathcal{M} is complete and cocomplete, then the trivial model structure even makes \mathcal{M} into a DHK model category.

Example 3.1.5. The **epi–mono model structure** on **Set** is defined by the following data:

- Every morphism is a weak equivalence.
- The cofibrations are the injective maps.
- The fibrations are the surjective maps.

The key observation is that **Set** admits a mono–epi weak factorisation system;^[2] in fact we may even choose the mono–epi factorisations functorially: for example, given a map $f : X \rightarrow Y$, we may take the cograph factorisation $X \rightarrow X \amalg Y \rightarrow Y$, where $X \rightarrow X \amalg Y$ is the coproduct insertion and $X \amalg Y \rightarrow Y$ is the map $[f, \text{id}_Y]$.

REMARK 3.1.6. Let \mathcal{M} be a category with finite limits and finite colimits. Then, $(\mathcal{C}, \mathcal{W}, \mathcal{F})$ is a model structure on \mathcal{M} if and only if $(\mathcal{F}^{\text{op}}, \mathcal{W}^{\text{op}}, \mathcal{C}^{\text{op}})$ is a model structure on \mathcal{M}^{op} .

Lemma 3.1.7. *Let \mathcal{M} be a category with a pair of weak factorisation systems $(\mathcal{C}', \mathcal{F})$ and $(\mathcal{C}, \mathcal{F}')$. Define the following subensemble of $\text{mor } \mathcal{C}$:*

$$\mathcal{W} = \{q \circ j \mid j \in \mathcal{C}', q \in \mathcal{F}'\}$$

- (i) $\mathcal{C} \cap \mathcal{W} \subseteq \mathcal{C}' \subseteq \mathcal{W}$.
- (ii) *If $\mathcal{C}' \subseteq \mathcal{C}$, then $\mathcal{F}' \subseteq \mathcal{F}$ and $\mathcal{C} \cap \mathcal{W} = \mathcal{C}'$.*

[2] — not to be confused with the epi–mono *orthogonal* factorisation system!

Dually:

(i') $\mathcal{W} \cap \mathcal{F} \subseteq \mathcal{F}' \subseteq \mathcal{W}$.

(ii') If $\mathcal{F}' \subseteq \mathcal{F}$, then $\mathcal{C}' \subseteq \mathcal{C}$ and $\mathcal{W} \cap \mathcal{F} = \mathcal{F}'$.

Proof. (i). If $j : X \rightarrow Y$ is in \mathcal{C}' , then j is also in \mathcal{W} , because id_Y is in \mathcal{F} ; thus $\mathcal{C}' \subseteq \mathcal{W}$. Now, suppose $i : X \rightarrow Z$ is in $\mathcal{C} \cap \mathcal{W}$; then there must be $j : X \rightarrow Y$ in \mathcal{C}' and $q : Y \rightarrow Z$ in \mathcal{F}' such that $i = q \circ j$, and so we have the commutative diagram shown below:

$$\begin{array}{ccc} X & \xrightarrow{j} & Y \\ i \downarrow & & \downarrow q \\ Z & \xrightarrow{\text{id}} & Z \end{array}$$

Since $i \sqsupseteq q$, i must be a retract of j ; hence, by [proposition A.3.12](#), i is in \mathcal{C}' , and therefore $\mathcal{C} \cap \mathcal{W} \subseteq \mathcal{C}'$.

(ii). If we know $\mathcal{C}' \subseteq \mathcal{C}$, then $\mathcal{F}' \subseteq \mathcal{F}$ by [proposition A.3.3](#), and $\mathcal{C}' \subseteq \mathcal{C} \cap \mathcal{W}$, so from claim (i) it follows that $\mathcal{C}' = \mathcal{C} \cap \mathcal{W}$. ■

Theorem 3.1.8. *Let \mathcal{M} be a locally small category and let $\mathcal{C}, \mathcal{W}, \mathcal{F}$ be subclasses of $\text{mor } \mathcal{M}$. Assuming \mathcal{M} has finite limits and finite colimits, the following are equivalent:*

- (i) $(\mathcal{C}, \mathcal{W}, \mathcal{F})$ is a model structure for \mathcal{M} .
- (ii) \mathcal{W} has the 2-out-of-3 property in \mathcal{M} , and both $(\mathcal{C} \cap \mathcal{W}, \mathcal{F})$ and $(\mathcal{C}, \mathcal{W} \cap \mathcal{F})$ are weak factorisation systems for \mathcal{M} .

Proof. (i) \Rightarrow (ii). The fact that we have two weak factorisation systems follows from Lemma 1.1 in [GJ, Ch. II] or Proposition 7.2.3 in [Hirschhorn, 2003].

(ii) \Rightarrow (i). We may deduce from [proposition A.3.12](#) that \mathcal{C} and \mathcal{F} are closed under retracts, and it remains to be shown that \mathcal{W} is closed under retracts.

Let $w : X \rightarrow Z$ be a morphism in \mathcal{W} , and consider a commutative diagram of the form below:

$$\begin{array}{ccccc} & & \text{id} & & \\ & \xrightarrow{\quad} & \text{---} & \xrightarrow{\quad} & \\ X' & \xrightarrow{s_X} & X & \xrightarrow{r_X} & X' \\ w' \downarrow & & w \downarrow & & \downarrow w' \\ Z' & \xrightarrow{s_Z} & Z & \xrightarrow{r_Z} & Z' \\ & & \text{id} & & \end{array}$$

Choose a $(C \cap \mathcal{W}, \mathcal{F})$ factorisation for w' , say $w' = p' \circ j'$, with $j' : X' \rightarrow Y'$ in $C \cap \mathcal{W}$ and $p' : Y' \rightarrow Z'$ in \mathcal{F} . Construct the following commutative diagram,

$$\begin{array}{ccccc}
 X' & \xrightarrow{s_X} & X & \xrightarrow{r_X} & X' \\
 j' \downarrow & & u \downarrow & & \downarrow j' \\
 Y' & \xrightarrow{s_Y} & Y & \xrightarrow{r_Y} & Y' \\
 p' \downarrow & & v \downarrow & & \downarrow p' \\
 Z' & \xrightarrow{s_Z} & Z & \xrightarrow{r_Z} & Z'
 \end{array}$$

where the top left square is a pushout square, $v \circ u = w$, and $r_Y \circ s_Y = \text{id}_Y$. Since $C \cap \mathcal{W}$ is closed under pushouts, u is also in $C \cap \mathcal{W}$, and by the 2-out-of-3 property, v is in \mathcal{W} . Thus, p' is in \mathcal{F} and is a retract of v :

$$\begin{array}{ccccc}
 & & \text{id} & & \\
 & \curvearrowright & & \curvearrowleft & \\
 Y' & \xrightarrow{s_Y} & Y & \xrightarrow{r_Y} & Y' \\
 p' \downarrow & & v \downarrow & & \downarrow p' \\
 Z' & \xrightarrow{s_Z} & Z & \xrightarrow{r_Z} & Z' \\
 & \curvearrowleft & & \curvearrowright & \\
 & & \text{id} & &
 \end{array}$$

Using the 2-out-of-3 property again, choose a $(C \cap \mathcal{W}, \mathcal{W} \cap \mathcal{F})$ -factorisation of v , say $v = q \circ j$. Since $j \square p'$, there exists a morphism r such that $r \circ j = r_Y$ and $p' \circ r = r_Z \circ q$; putting $s = j \circ s_Y$, we obtain $r \circ s = r_Y \circ s_Y = \text{id}_Y$, thus p' is a retract of q and must therefore be in $\mathcal{F} \cap \mathcal{W}$. Hence, $w' = p' \circ j'$ is in \mathcal{W} . \square

REMARK 3.1.9. May and Ponto [2012, Ch. 14] define ‘model category’ to mean a complete and cocomplete locally small category \mathcal{M} equipped with a triple of classes $(C, \mathcal{W}, \mathcal{F})$ satisfying condition (ii) of the above proposition; if the two weak factorisation systems can be extended to a pair of functorial factorisation systems, then this is a DHK model category.

Lemma 3.1.10. *Let A be an object in a model category \mathcal{M} . Then the slice category $\mathcal{M}_{/A}$ has the **slice model structure**, where a morphism in $\mathcal{M}_{/A}$ is a cofibration, weak equivalence, or fibration if it is so in \mathcal{M} .*

Proof. Use lemmas 2.1.6 and A.3.II, plus the fact that $\mathcal{M}_{/A}$ has finite limits and finite colimits if \mathcal{M} does. \blacksquare

Definition 3.1.11. Let X be an object in a model category \mathcal{M} .

- A **cofibrant replacement** for X is a pair (\tilde{X}, p) where \tilde{X} is a cofibrant object in \mathcal{M} and p is a weak equivalence $\tilde{X} \rightarrow X$.
- A **fibrant replacement** for X is a pair (\hat{X}, i) where \hat{X} is a fibrant object in \mathcal{M} and i is a weak equivalence $X \rightarrow \hat{X}$.
- A **fibrant cofibrant replacement** for X is a cofibrant replacement (\tilde{X}, p) where $p : \tilde{X} \rightarrow X$ is a trivial fibration.
- A **cofibrant fibrant replacement** for X is a fibrant replacement (\hat{X}, i) where $i : X \rightarrow \hat{X}$ is a trivial cofibration.

Definition 3.1.12. Let \mathcal{M} be a model category.

- A **cofibrant replacement functor** for \mathcal{M} is a pair (Q, p) , where Q is an endofunctor on \mathcal{M} and p is a natural transformation $Q \Rightarrow \text{id}_{\mathcal{M}}$ such that, for every object X in \mathcal{M} , (QX, p_X) is a cofibrant replacement for X .
- A **fibrant replacement functor** for \mathcal{M} is a pair (R, i) , where R is an endofunctor on \mathcal{M} and i is a natural transformation $\text{id}_{\mathcal{M}} \Rightarrow R$ such that, for every object X in \mathcal{M} , (RX, i_X) is a fibrant replacement for X .
- A **fibrant cofibrant replacement functor** for \mathcal{M} is a pair (Q, p) , where Q is an endofunctor on \mathcal{M} and p is a natural transformation $Q \Rightarrow \text{id}_{\mathcal{M}}$ such that, for every object X in \mathcal{M} , (QX, p_X) is a fibrant cofibrant replacement for X .
- A **cofibrant fibrant replacement functor** for \mathcal{M} is a pair (R, i) , where R is an endofunctor on \mathcal{M} and i is a natural transformation $\text{id}_{\mathcal{M}} \Rightarrow R$ such that, for every object X in \mathcal{M} , (RX, i_X) is a cofibrant fibrant replacement for X .

REMARK 3.1.13. Note that a fibrant cofibrant replacement for X is precisely a cofibrant replacement for X that is fibrant as an object in \mathcal{M}/X , and a cofibrant fibrant replacement for X is precisely a fibrant replacement for X that is cofibrant as an object in ${}^X/\mathcal{M}$.

Moreover, if X is fibrant and (\tilde{X}, p) is a fibrant cofibrant replacement for X , then \tilde{X} is both fibrant and cofibrant in \mathcal{M} , and if X is cofibrant and (\hat{X}, i) is a cofibrant fibrant replacement for X , then \hat{X} is both cofibrant and fibrant in \mathcal{M} .

Proposition 3.1.14.

- (i) Any object in a model category has both a fibrant cofibrant replacement and a cofibrant fibrant replacement.
- (ii) Any DHK model category has both a fibrant cofibrant replacement functor and a cofibrant fibrant replacement functor.

Proof. (i). Use axiom CM5.

(ii). Use axiom CM5*. ■

Proposition 3.1.15. Let \mathcal{M} be a model category, let \mathcal{U} be the class of trivial cofibrations in \mathcal{M} , and let \mathcal{V} be the class of trivial fibrations in \mathcal{M} .

1. \mathcal{U} and \mathcal{V} are closed under composition.
2. \mathcal{U} is closed under pushouts, and \mathcal{V} is closed under pullbacks.
3. \mathcal{M} admits a three-arrow calculus^[3] with respect to $(\mathcal{U}, \mathcal{V})$, which is functorial^[4] if \mathcal{M} satisfies axiom CM5*.

Proof. (i) and (ii). These claims follow from [proposition A.3.12](#) applied to [theorem 3.1.8](#).

(iii). Axioms CM2 and CM5 imply axiom A1 is satisfied, and axioms A2 and A3 follow from the above claims; that we get a functorial three-arrow calculus under axiom CM5* is an obvious consequence of the universal property of pushouts and pullbacks. ■

3.2 Left and right homotopy

Prerequisites. § 3.1.

Definition 3.2.1. Let X be an object in a model category \mathcal{M} .

- A **cylinder object** for X is a quadruple $(\text{Cyl}(X), i_0, i_1, p)$, where $\text{Cyl}(X)$ is an object in \mathcal{M} , $p : \text{Cyl}(X) \rightarrow X$ is a weak equivalence, and i_0 and i_1 are sections of p such that the morphism $[i_0, i_1] : X + X \rightarrow \text{Cyl}(X)$ is a cofibration.

[3] See [definition 2.5.1](#).

[4] See [definition 2.5.4](#).

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- A **path object** for X is a quadruple $(\text{Path}(X), i, p_0, p_1)$, where $\text{Path}(X)$ is an object in \mathcal{M} , $i : X \rightarrow \text{Path}(X)$ is a weak equivalence, and p_0 and p_1 are retractions of i such that the morphism $\langle p_0, p_1 \rangle : \text{Path}(X) \rightarrow X \times X$ is a fibration.

REMARK 3.2.2. Let $(\text{Cyl}(X), i_0, i_1, p)$ be a cylinder object for X . By definition, $p \circ i_0 = p \circ i_1 = \text{id}_X$, and p is a weak equivalence, so by the 2-out-of-3 property, i_0 and i_1 must also be weak equivalences $X \rightarrow \text{Cyl}(X)$.

Dually, if $(\text{Path}(X), i, p_0, p_1)$ is a path object for X , then p_0 and p_1 must be weak equivalences $\text{Path}(X) \rightarrow X$.

Proposition 3.2.3. *Let X be an object in a model category \mathcal{M} .*

- *There exists a cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ for X , where the morphism $p : \text{Cyl}(X) \rightarrow X$ is a trivial fibration.*
- *There exists a path object $(\text{Path}(X), i, p_0, p_1)$ for X , where the morphism $i : X \rightarrow \text{Path}(X)$ is a trivial cofibration.*

Proof. Use axioms CM1 and CM5. ■

Definition 3.2.4. Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category \mathcal{M} , let $(\text{Cyl}(X), i_0, i_1, p)$ be a cylinder object for X , and let $(\text{Path}(Y), i, p_0, p_1)$ be a path object for Y .

- A **left homotopy** from f_0 to f_1 with respect to $(\text{Cyl}(X), i_0, i_1, p)$ is a morphism $H : \text{Cyl}(X) \rightarrow Y$ such that $H \circ i_0 = f_0$ and $H \circ i_1 = f_1$.
- A **right homotopy** from f_0 to f_1 with respect to $(\text{Path}(Y), i, p_0, p_1)$ is a morphism $H : X \rightarrow \text{Path}(Y)$ such that $p_0 \circ H = f_0$ and $p_1 \circ H = f_1$.
- We say f_0 and f_1 are **left homotopic** if there exists a left homotopy from f_0 to f_1 with respect to some cylinder object for X .
- We say f_0 and f_1 are **right homotopic** if there exists a right homotopy from f_0 to f_1 with respect to some path object for Y .

REMARK 3.2.5. If f_0 and f_1 are either left homotopic or right homotopic, then they must represent the same morphism in $\text{Ho } \mathcal{M}$. For definiteness, let us write $\gamma : \mathcal{M} \rightarrow \text{Ho } \mathcal{M}$ for the localising functor, and suppose $H : \text{Cyl}(X) \rightarrow Y$ is a left homotopy from f_0 to f_1 . Since i_0 and i_1 are both sections of the weak

equivalence $p : \text{Cyl}(X) \rightarrow X$, we must have $\gamma i_0 = (\gamma p)^{-1} = \gamma i_1$; but $f_0 = H \circ i_0$ and $f_1 = H \circ i_1$, so indeed $\gamma f_0 = \gamma f_1$. This is one of the reasons for calling $\text{Ho } \mathcal{M}$ the *homotopy category* of \mathcal{M} .

However, it is not quite true that $\gamma f_0 = \gamma f_1$ if and only if f_0 and f_1 are either left homotopic or right homotopic; this only happens in special cases. In general, being left/right homotopic fails to even be an equivalence relation.

Definition 3.2.6. Let $f : X \rightarrow Y$ be a morphism in a model category \mathcal{M} .

- A **left homotopy left inverse** of f is a morphism $g : Y \rightarrow X$ in \mathcal{M} such that $g \circ f$ and id_X are left homotopic.
- A **right homotopy right inverse** of f is a morphism $h : Y \rightarrow X$ in \mathcal{M} such that $f \circ h$ and id_Y are right homotopic.
- A **right homotopy left inverse** of f is a morphism $g : Y \rightarrow X$ in \mathcal{M} such that $g \circ f$ and id_X are right homotopic.
- A **left homotopy right inverse** of f is a morphism $h : Y \rightarrow X$ in \mathcal{M} such that $f \circ h$ and id_Y are left homotopic.

A **homotopy equivalence** in \mathcal{M} is a pair (f, g) such that g (resp. f) is both a left homotopy left inverse and a right homotopy right inverse for f (resp. g). Two morphisms $f : X \rightarrow Y$ and $g : Y \rightarrow X$ in \mathcal{M} are **mutual homotopy inverses** when (f, g) constitute a homotopy equivalence in \mathcal{M} .

REMARK 3.2.7. Let $f : X \rightarrow Y$ and $g : Y \rightarrow X$ be morphisms in a model category.

- g is a left homotopy left inverse for f if and only if f is a left homotopy right inverse for g .
- g is a right homotopy left inverse for f if and only if f is a right homotopy left inverse for g .

However, note that the dual of ‘left homotopy left inverse’ is ‘right homotopy right inverse’, and the dual of ‘right homotopy left inverse’ is ‘left homotopy right inverse’!

Lemma 3.2.8. Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category, and suppose f_0 and f_1 are either left or right homotopic. Then, f_0 is a weak equivalence if and only if f_1 is a weak equivalence.

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Proof. Assume f_0 and f_1 are left homotopic; the other case is formally dual. So, there exist a cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ for X and a morphism $H : \text{Cyl}(X) \rightarrow Y$ such that $H \circ i_0 = f_0$ and $H \circ i_1 = f_1$. Suppose f_0 is a weak equivalence. By [remark 3.2.2](#), i_0 is a weak equivalence, so the 2-out-of-3 property implies H is also a weak equivalence; but i_1 is a weak equivalence as well, so f_1 must be a weak equivalence too. A symmetrical argument proves that f_0 is a weak equivalence if f_1 is. \blacksquare

Lemma 3.2.9. *Let $f : X \rightarrow Y$ and $g : Y \rightarrow X$ be morphisms in a model category \mathcal{M} .*

- (i) *If $g \circ f$ is either left or right homotopic to id_X , and $f \circ g$ is either left or right homotopic to id_Y , then (f, g) is an equivalence in \mathcal{M} (in the sense of [definition 2.1.12](#)).*
- (ii) *If there exist morphisms $g, h : Y \rightarrow X$ such that $g \circ f$ is either left or right homotopic to id_X and $f \circ h$ is either left or right homotopic to id_Y , then (the image of) f is an isomorphism in $\text{Ho } \mathcal{M}$.*

Proof. Obvious, given [remark 3.2.5](#). \blacklozenge

Lemma 3.2.10. *Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category \mathcal{M} .*

- (i) *Given any cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ for X , $f_0 \circ p : \text{Cyl}(X) \rightarrow Y$ is a left homotopy from f_0 to itself.*
- (ii) *If $H : \text{Cyl}(X) \rightarrow Y$ is a left homotopy from f_0 to f_1 with respect to a cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ for X , then the same H is a left homotopy from f_1 to f_0 for the cylinder object $(\text{Cyl}(X), i_1, i_0, p)$.*

Dually:

- (i') *Given any path object $(\text{Path}(Y), i, p_0, p_1)$ for Y , $i \circ f_0 : X \rightarrow \text{Path}(Y)$ is a right homotopy from f_0 to itself.*
- (ii') *If $H : X \rightarrow \text{Path}(Y)$ is a right homotopy from f_0 to f_1 with respect to a path object $(\text{Path}(Y), i, p_0, p_1)$ for Y , then the same H is a right homotopy from f_1 to f_0 for the path object $(\text{Path}(Y), i, p_1, p_0)$.*

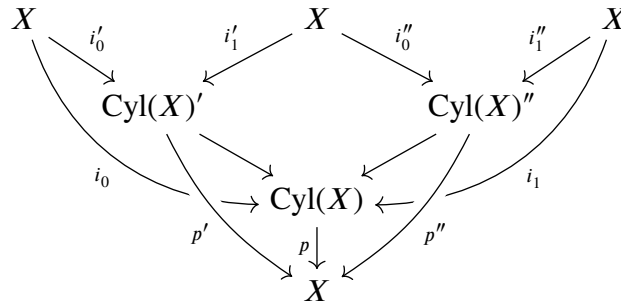
Proof. Obvious. \blacklozenge

Lemma 3.2.11. *Let \mathcal{M} be a model category.*

- *If $(\text{Cyl}(X), i_0, i_1, p)$ is a cylinder object for a cofibrant object in \mathcal{M} , then the insertions $i_0, i_1 : X \rightarrow \text{Cyl}(X)$ are trivial cofibrations, and $\text{Cyl}(X)$ is a cofibrant object in \mathcal{M} .*
- *If $(\text{Path}(Y), i, p_0, p_1)$ is a path object for a fibrant object in \mathcal{M} , then the projections $p_0, p_1 : \text{Path}(Y) \rightarrow Y$ are trivial fibrations, and $\text{Path}(Y)$ is a fibrant object in \mathcal{M} .*

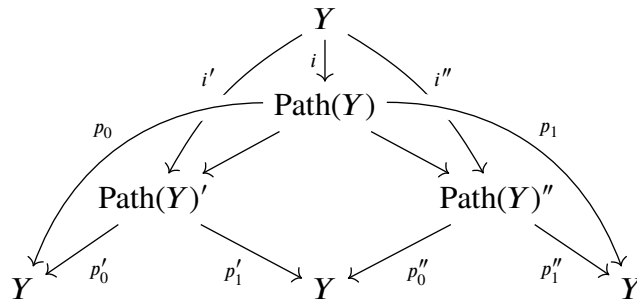
Proof. See Lemmas 1.5 and 1.7 in [GJ], or Lemma 7.3.6 in [Hirschhorn, 2003]. □

Lemma 3.2.12. *Let X be a cofibrant object in a model category \mathcal{M} . Given two cylinder objects for X , say $(\text{Cyl}(X)', i'_0, i'_1, p')$ and $(\text{Cyl}(X)'', i''_0, i''_1, p'')$, there exists a third cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ such that the diagram below commutes,*



and the diamond is a pushout diagram.

Dually, if Y is a fibrant object in \mathcal{M} , and we have two path objects for Y , say $(\text{Path}(Y)', i', p'_0, p'_1)$ and $(\text{Path}(Y)'', i'', p''_0, p''_1)$, then there exists a third path object $(\text{Path}(Y), i, p_0, p_1)$ such that the diagram below commutes,



and the diamond is a pullback diagram.

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Proof. See Lemmas 1.5 and 1.7 in [GJ, Ch. II], or Lemma 7.4.2 in [Hirschhorn, 2003]. \square

Corollary 3.2.13. *Let $f_0, f_1, f_2 : X \rightarrow Y$ be three parallel morphisms in a model category \mathcal{M} .*

- (i) *If f_0 and f_1 are left homotopic, and f_1 and f_2 are left homotopic, then f_0 and f_2 are also left homotopic.*
- (ii) *If f_0 and f_1 are right homotopic, and f_1 and f_2 are right homotopic, then f_0 and f_2 are also right homotopic. \blacksquare*

Lemma 3.2.14. *Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category \mathcal{M} .*

- (i) *If X is cofibrant, and f_0 and f_1 are left homotopic, given any path object $(\text{Path}(Y), i, p_0, p_1)$ for Y , there is a right homotopy $H : X \rightarrow \text{Path}(Y)$ from f_0 to f_1 .*
- (ii) *If Y is fibrant, and f_0 and f_1 are right homotopic, given any cylinder object $(\text{Cyl}(X), i_0, i_1, p)$ for X , there is a left homotopy $H : \text{Cyl}(X) \rightarrow Y$ from f_0 to f_1 .*

Proof. See Proposition 1.8 in [GJ, Ch. II], or Proposition 7.4.7 in [Hirschhorn, 2003]. \square

Proposition 3.2.15. *Let X and Y be objects in a model category \mathcal{M} .*

- (i) *If X is cofibrant, then being left homotopic is an equivalence relation on the hom-set $\mathcal{M}(X, Y)$.*
- (ii) *If Y is fibrant, then being right homotopic is an equivalence relation on the hom-set $\mathcal{M}(X, Y)$.*
- (iii) *If X is cofibrant and Y is fibrant, then these two equivalence relations on $\mathcal{M}(X, Y)$ coincide.*

Proof. Use the preceding lemmas. \blacksquare

Lemma 3.2.16. *Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category \mathcal{M} .*

- (i) *If f_0 and f_1 are right homotopic and $g : W \rightarrow X$ is any morphism in \mathcal{M} , then $f_0 \circ g$ and $f_1 \circ g$ are also right homotopic.*
- (ii) *If f_0 and f_1 are left homotopic and $g : Y \rightarrow Z$ is any morphism in \mathcal{M} , then $g \circ f_0$ and $g \circ f_1$ are also left homotopic.*

Proof. Obvious. ◆

Corollary 3.2.17. *Let \mathcal{M} be a model category, and let \mathcal{M}_{cf} be the full subcategory spanned by the cofibrant–fibrant objects. Then the equivalence relation induced by homotopy is a congruence on \mathcal{M}_{cf} ; in particular, there exist a locally small category \mathcal{M}' and a full functor $\mathcal{M}_{\text{cf}} \rightarrow \mathcal{M}'$ with these properties:*

- *The objects of \mathcal{M}' are those of \mathcal{M}_{cf} .*
- *The hom-set $\mathcal{M}'(X, Y)$ is $\mathcal{M}(X, Y)$ modulo homotopy.*
- *The functor $\mathcal{M}_{\text{cf}} \rightarrow \mathcal{M}'$ sends each morphism in \mathcal{M}' to its homotopy class.* ■

The next result is a version of Whitehead’s theorem; however, this is a purely formal consequence of the model category axioms and has no real content, unlike the original theorem.

Proposition 3.2.18. *Let X and Y be cofibrant–fibrant objects in a model category \mathcal{M} . If $f : X \rightarrow Y$ is a weak equivalence, then f has a homotopy inverse in \mathcal{M} .*

Proof. See Theorem 1.10 in [GJ, Ch. II], or Theorem 7.5.10 in [Hirschhorn, 2003]. □

Lemma 3.2.19. *Let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms in a model category \mathcal{M} .*

- *If $g : W \rightarrow X$ is a morphism with a right homotopy right inverse in \mathcal{M} , then $f_0 \circ g$ and $f_1 \circ g$ are right homotopic if and only if f_0 and f_1 are right homotopic.*

- If $g : Y \rightarrow Z$ is a morphism with a left homotopy left inverse in \mathcal{M} , then $g \circ f_0$ and $g \circ f_1$ are left homotopic if and only if f_0 and f_1 are left homotopic.

Proof. This follows immediately from the definitions and [lemma 3.2.16](#). ■

Corollary 3.2.20. *Let W, X, Y, Z be cofibrant–fibrant objects in a model category \mathcal{M} , and let $f_0, f_1 : X \rightarrow Y$ be a parallel pair of morphisms.*

- If $g : W \rightarrow X$ is a weak equivalence such that $f_0 \circ g$ and $f_1 \circ g$ are homotopic, then f_0 and f_1 are homotopic.
- If $g : Y \rightarrow Z$ is a weak equivalence such that $g \circ f_0$ and $g \circ f_1$ are homotopic, then f_0 and f_1 are homotopic.

Proof. Apply [proposition 3.2.18](#) in conjunction with the above lemma. ■

3.3 The homotopy category

Prerequisites. §§ 3.1, 3.2, A.4.

Definition 3.3.1. The **Quillen homotopy category** (or, more simply, **homotopy category**) of a model category \mathcal{M} is the category $\text{Ho } \mathcal{M}$ obtained by freely inverting the weak equivalences in \mathcal{M} , as in [definition A.4.9](#).

Theorem 3.3.2. *Let \mathcal{M} be a model category and let $\gamma : \mathcal{M} \rightarrow \text{Ho } \mathcal{M}$ be the localising functor.*

- (i) $\text{Ho } \mathcal{M}$ is equivalent to the locally small category \mathcal{M}' defined in [corollary 3.3.4](#), and \mathcal{M} is a saturated homotopical category.
- (ii) If X and Y are cofibrant–fibrant objects in \mathcal{M} , then the hom-ensemble map $\mathcal{M}(X, Y) \rightarrow \text{Ho } \mathcal{M}(X, Y)$ induced by γ is surjective; and moreover for any parallel pair $f_0, f_1 : X \rightarrow Y$ in \mathcal{M} , we have $\gamma f_0 = \gamma f_1$ if and only if f_0 and f_1 are homotopic.
- (iii) For any two objects X and Y in \mathcal{M} , every morphism $X \rightarrow Y$ in $\text{Ho } \mathcal{M}$ can be represented as a zigzag of the form

$$X \xleftarrow{p} \tilde{X} \longrightarrow \hat{Y} \xleftarrow{i} Y$$

where (\tilde{X}, p) is any cofibrant replacement for X and (\hat{Y}, i) is any fibrant replacement for Y .

Proof. (i). This is Theorem 1.11 in [GJ, Ch. II], or Proposition 5.8 in [DS].

(ii). Implied by claim (i).

(iii). Using claim (ii), every morphism $X \rightarrow Y$ in $\text{Ho } \mathcal{M}$ can be represented as a zigzag of the form

$$\begin{array}{ccc} X & \xleftarrow{p} & \tilde{X} & & \hat{Y} & \xleftarrow{i} & Y \\ & & \downarrow i' & & \uparrow p' & & \\ & & R\tilde{X} & \xrightarrow{f'} & Q\hat{Y} & & \end{array}$$

where $(R\tilde{X}, i')$ is a cofibrant fibrant replacement for \tilde{X} and $(Q\hat{Y}, p')$ is a fibrant cofibrant replacement for \hat{Y} ; but such a zigzag is manifestly equivalent to the zigzag

$$X \xleftarrow{p} \tilde{X} \xrightarrow{f} \hat{Y} \xleftarrow{i} Y$$

where $f = p' \circ f' \circ i'$. □

Corollary 3.3.3. *Let $f : X \rightarrow Y$ be a morphism in a model category \mathcal{M} . If f has a quasi-inverse in \mathcal{M} (in the sense of [definition 2.1.12](#)), then f is a weak equivalence in \mathcal{M} .*

Proof. If f has a quasi-inverse in \mathcal{M} , then (the image of) f is an isomorphism in $\text{Ho } \mathcal{M}$; but \mathcal{M} is a saturated homotopical category, so f must be a weak equivalence in \mathcal{M} . ■

Corollary 3.3.4. *Let \mathcal{M} be a model category and let $\gamma : \mathcal{M} \rightarrow \text{Ho } \mathcal{M}$ be the localising functor.*

- (i) *If X is a cofibrant object in \mathcal{M} and Y is a fibrant object in \mathcal{M} , then the hom-ensemble map $\mathcal{M}(X, Y) \rightarrow \text{Ho } \mathcal{M}(X, Y)$ induced by γ is surjective.*
- (ii) *Moreover, for any parallel pair $f_0, f_1 : X \rightarrow Y$ in \mathcal{M} , if X is cofibrant and Y is fibrant, we have $\gamma f_0 = \gamma f_1$ if and only if f_0 and f_1 are homotopic.*
- (iii) *The full subcategory \mathcal{M}_{cf} of cofibrant–fibrant objects in \mathcal{M} has the Whitehead property (in the sense of [definition 2.1.16](#)).*

Proof. (i). This immediately follows from statement (iii) of the above theorem.

(ii). As noted in [remark 3.2.5](#), if $f_0, f_1 : X \rightarrow Y$ are homotopic, then we must have $\gamma f_0 = \gamma f_1$. Conversely, suppose $\gamma f_0 = \gamma f_1$ with X cofibrant and Y fibrant. Let (RX, i') be a cofibrant fibrant replacement for X and (QY, p') be a fibrant cofibrant replacement for Y . Then, there exists morphisms $f'_0, f'_1 : RX \rightarrow QY$ such that $f_0 = p' \circ f'_0 \circ i'$ and $f_1 = p' \circ f'_1 \circ i'$. Since $i' : X \rightarrow RX$ and $p' : QY \rightarrow Y$ are weak equivalences, we must have $\gamma f'_0 = \gamma f'_1$ in $\text{Ho } \mathcal{M}$. The theorem then implies f'_0 and f'_1 are homotopic; thus f_0 and f_1 are also homotopic, by [lemmas 3.2.14](#) and [3.2.16](#).

(iii). Apply [theorem 2.1.17](#) in conjunction with [lemma 3.2.9](#) and the above corollary. ■

Corollary 3.3.5. *Let $f : X \rightarrow Y$ be a morphism between two cofibrant objects in a model category \mathcal{M} . The following are equivalent:*

- (i) $f : X \rightarrow Y$ is a weak equivalence in \mathcal{M} .
- (ii) $\text{Ho } \mathcal{M}(f, Z) : \text{Ho } \mathcal{M}(Y, Z) \rightarrow \text{Ho } \mathcal{M}(X, Z)$ is a bijection for all cofibrant-fibrant objects Z in \mathcal{M} .
- (iii) $\mathcal{M}'(f, Z) : \mathcal{M}'(Y, Z) \rightarrow \mathcal{M}'(X, Z)$ is a bijection for all cofibrant-fibrant objects Z in \mathcal{M} , where $\mathcal{M}'(Y, Z)$ (resp. $\mathcal{M}'(X, Z)$) denotes the set of all morphisms $Y \rightarrow Z$ (resp. $X \rightarrow Z$) in \mathcal{M} modulo homotopy.

Proof. (i) \Rightarrow (ii). Every weak equivalence in \mathcal{M} becomes an isomorphism in $\text{Ho } \mathcal{M}$, so in particular $\text{Ho } \mathcal{M}(f, Z) : \text{Ho } \mathcal{M}(Y, Z) \rightarrow \text{Ho } \mathcal{M}(X, Z)$ must be a bijection.

(ii) \Leftrightarrow (iii). The previous corollary implies that the vertical arrows in the following commutative diagram are bijections,

$$\begin{array}{ccc}
 \mathcal{M}'(Y, Z) & \xrightarrow{\mathcal{M}'(f, Z)} & \mathcal{M}'(X, Z) \\
 \downarrow & & \downarrow \\
 \text{Ho } \mathcal{M}(Y, Z) & \xrightarrow{\text{Ho } \mathcal{M}(f, Z)} & \text{Ho } \mathcal{M}(X, Z)
 \end{array}$$

and so $\mathcal{M}'(f, Z)$ is a bijection if and only if $\text{Ho } \mathcal{M}(f, Z)$ is a bijection.

(ii) \Rightarrow (i). Suppose (\hat{X}, i_X) is a cofibrant fibrant replacement for X and (\hat{Y}, i_Y) is a cofibrant fibrant replacement for Y . Then, (by axiom CM4) there exists a morphism $\hat{f} : \hat{X} \rightarrow \hat{Y}$ making the diagram below commute,

$$\begin{array}{ccc} X & \xrightarrow{f} & Y \\ i_X \downarrow & & \downarrow i_Y \\ \hat{X} & \xrightarrow{\hat{f}} & \hat{Y} \end{array}$$

and by the 2-out-of-3 property, f is a weak equivalence if and only if \hat{f} is a weak equivalence. On the other hand, the following diagram also commutes,

$$\begin{array}{ccc} \mathrm{Ho} \mathcal{M}(\hat{Y}, Z) & \xrightarrow{\mathrm{Ho} \mathcal{M}(\hat{f}, Z)} & \mathrm{Ho} \mathcal{M}(\hat{X}, Z) \\ \mathrm{Ho} \mathcal{M}(i_Y, Z) \downarrow & & \downarrow \mathrm{Ho} \mathcal{M}(i_X, Z) \\ \mathrm{Ho} \mathcal{M}(Y, Z) & \xrightarrow{\mathrm{Ho} \mathcal{M}(f, Z)} & \mathrm{Ho} \mathcal{M}(X, Z) \end{array}$$

and so $\mathrm{Ho} \mathcal{M}(f, Z)$ is a bijection if and only if $\mathrm{Ho} \mathcal{M}(\hat{f}, Z)$ is a bijection; but \hat{X} and \hat{Y} are both cofibrant–fibrant objects, so if $\mathrm{Ho} \mathcal{M}(f, Z)$ is a bijection for all cofibrant–fibrant objects Z , then \hat{f} must be a weak equivalence (because \mathcal{M} is a saturated homotopical category). \blacksquare

Proposition 3.3.6 (Joyal). *Let \mathcal{M} and \mathcal{M}' be two model categories with the same underlying category. If cofibrations in \mathcal{M} are cofibrations in \mathcal{M}' and vice versa, then the following are equivalent:*

- (i) *Every weak equivalence in \mathcal{M} is a weak equivalence in \mathcal{M}' .*
- (ii) *Every fibrant object in \mathcal{M}' is a fibrant object in \mathcal{M} .*
- (iii) *Every cofibrant–fibrant object in \mathcal{M}' is a cofibrant–fibrant object in \mathcal{M} .*
- (iv) *Every weak equivalence between cofibrant objects in \mathcal{M} is a weak equivalence between cofibrant objects in \mathcal{M}' .*

Proof. This result is due to Joyal [2010].

(i) \Rightarrow (ii). Since every trivial cofibration in \mathcal{M} is a trivial cofibration in \mathcal{M}' , [theorem 3.1.8](#) (plus the definition of weak factorisation system) implies every

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fibration in \mathcal{M}' is a fibration in \mathcal{M} ; in particular, every fibrant object in \mathcal{M}' is a fibrant object in \mathcal{M} .

(ii) \Rightarrow (iii). Obvious.

(iii) \Rightarrow (iv). Let $f : X \rightarrow Y$ be a weak equivalence between cofibrant objects in \mathcal{M} . X and Y are also cofibrant objects in \mathcal{M}' , and by [proposition 3.2.3](#), we may choose cylinder objects for X and Y in \mathcal{M} that are also cylinder objects in \mathcal{M}' , since the trivial fibrations in \mathcal{M} and \mathcal{M}' are the same. Now, if Z is a cofibrant–fibrant object in \mathcal{M}' , then it is also a cofibrant–fibrant object in \mathcal{M} , and so by [lemma 3.2.10](#), we deduce that the homotopy relation on morphisms $X \rightarrow Z$ (resp. $Y \rightarrow Z$) in \mathcal{M} agrees with the homotopy relation on morphisms $X \rightarrow Z$ (resp. $Y \rightarrow Z$) in \mathcal{M}' . Thus, applying [corollary 3.3.5](#), we conclude that $f : X \rightarrow Y$ is also a weak equivalence in \mathcal{M}' .

(iv) \Rightarrow (i). Let $f : X \rightarrow Y$ be a weak equivalence in \mathcal{M} , let (\tilde{X}, p_X) be a fibrant cofibrant replacement for X in \mathcal{M} , and let (\tilde{Y}, p_Y) be a fibrant cofibrant replacement for Y in \mathcal{M} . There exists a morphism $\tilde{f} : \tilde{X} \rightarrow \tilde{Y}$ making the following diagram commute,

$$\begin{array}{ccc} \tilde{X} & \xrightarrow{\tilde{f}} & \tilde{Y} \\ p_X \downarrow & & \downarrow p_Y \\ X & \xrightarrow{f} & Y \end{array}$$

and by the 2-out-of-3 property, $\tilde{f} : \tilde{X} \rightarrow \tilde{Y}$ is a weak equivalence between cofibrant objects in \mathcal{M} . The hypothesis says \tilde{f} is also a weak equivalence between cofibrant objects in \mathcal{M}' , and p_X and p_Y are trivial cofibrations in \mathcal{M}' , so we conclude that $f : X \rightarrow Y$ is a weak equivalence in \mathcal{M}' as well. \blacksquare

Theorem 3.3.7 (Determination principle). *A model structure is uniquely determined by any one of the following sets of data:*

- (i) *The cofibrations and the weak equivalences.*
- (ii) *The cofibrations and the trivial cofibrations.*
- (iii) *The cofibrations and the fibrant objects.*
- (iv) *The cofibrations and the cofibrant–fibrant objects.*

- (v) *The cofibrations and the weak equivalences between cofibrant objects.*
- (vi) *The cofibrations and the fibrations.*
- (vii) *The trivial cofibrations and the trivial fibrations.*
- (i') *The fibrations and the weak equivalences.*
- (ii') *The fibrations and the trivial fibrations.*
- (iii') *The fibrations and the cofibrant objects.*
- (iv') *The fibrations and the cofibrant–fibrant objects.*
- (v') *The fibrations and the weak equivalences between fibrant objects.*

Proof. (i) and (ii). By [theorem 3.1.8](#), the fibrations are precisely the morphisms with the right lifting property with respect to every trivial cofibration.

(iii), (iv), and (v). Apply Joyal’s result ([proposition 3.3.6](#)) and reduce to case (i).

(vi). The trivial cofibrations are precisely the morphisms with the left lifting property with respect to all fibrations, and the trivial fibrations are precisely the morphisms with the right lifting property with respect to all cofibrations, so this reduces to case (vii).

(vii). Axioms CM2 and CM5 imply that every weak equivalence is of the form $p \circ i$ where i is a trivial cofibration and p is a trivial fibration. Thus, the trivial cofibrations and the trivial fibrations together determine the weak equivalences. On the other hand, the trivial cofibrations determine the fibrations, and the trivial fibrations determine the cofibrations, thus the entire model structure is determined. ■

3.4 Quillen functors

Prerequisites. §§ [2.3](#), [2.4](#), [3.1](#).

Definition 3.4.1. A **left Quillen functor** is a functor $F : \mathcal{N} \rightarrow \mathcal{M}$ between model categories that has a right adjoint and preserves cofibrations and trivial cofibrations; dually, a **right Quillen functor** is a functor $G : \mathcal{M} \rightarrow \mathcal{N}$ between

model categories that has a left adjoint and preserves fibrations and trivial fibrations. A **Quillen adjunction** is an adjunction

$$F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$$

where \mathcal{M} and \mathcal{N} are model categories, such that F is a left Quillen functor and G is a right Quillen functor. A **Quillen equivalence** is a Quillen adjunction as above satisfying this additional condition:

- Given a cofibrant object A in \mathcal{N} and fibrant object X in \mathcal{M} , a morphism $FA \rightarrow Y$ is a weak equivalence in \mathcal{M} if and only if its right adjoint transpose $A \rightarrow GY$ is a weak equivalence in \mathcal{N} .

Proposition 3.4.2. *Let $F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$ be an adjunction between model categories. The following are equivalent:*

- (i) $F \dashv G$ is a Quillen adjunction.
- (ii) F is a left Quillen functor.
- (iii) G is a right Quillen functor.
- (iv) F preserves cofibrations and G preserves fibrations.
- (v) F preserves trivial cofibrations and G preserves trivial fibrations.

Proof. Use [proposition A.3.19](#). ■

REMARK 3.4.3. A functor between model categories that preserves both trivial cofibrations and trivial fibrations must also preserve weak equivalences, since axioms CM2 and CM5 together imply that a morphism is a weak equivalence if and only if it is of the form $p \circ i$ where i is a trivial cofibration and p is a trivial fibration. In particular, a functor that is both left and right Quillen must be homotopical.

Proposition 3.4.4.

- (i) *A left Quillen functor preserves cofibrant objects, and a right Quillen functor preserves fibrant objects.*
- (ii) *The composite of two Quillen adjunctions is also a Quillen adjunction.*
- (iii) *The composite of two Quillen equivalences is also a Quillen equivalence.*

Proof. Obvious. ◆

Lemma 3.4.5 (Kenneth S. Brown). *Let \mathcal{M} be a model category and let \mathcal{C} be a category with weak equivalences. If $F : \mathcal{M} \rightarrow \mathcal{C}$ sends trivial cofibrations (resp. trivial fibrations) in \mathcal{M} to weak equivalences in \mathcal{C} , then F preserves all weak equivalences between cofibrant (resp. fibrant) objects.*

Proof. See Lemma 9.9 in [DS], Lemma 7.7.1 in [Hirschhorn, 2003], or Lemma 14.5 in [DHKS]. □

Corollary 3.4.6. *Let $F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$ be a Quillen adjunction.*

- *If A and B are cofibrant objects in \mathcal{N} and $f : A \rightarrow B$ is a weak equivalence in \mathcal{N} , then Ff is a weak equivalence in \mathcal{M} .*
- *If X and Y are fibrant objects in \mathcal{M} and $g : X \rightarrow Y$ is a weak equivalence in \mathcal{M} , then Gg is a weak equivalence in \mathcal{N} .* ■

Proposition 3.4.7 (Dugger). *Let $F \dashv G$ be an adjunction between DHK model categories. The following are equivalent:*

- (i) *$F \dashv G$ is a Quillen adjunction.*
- (ii) *F preserves cofibrations between cofibrant objects and all trivial cofibrations.*
- (iii) *G preserves fibrations between fibrant objects and all trivial fibrations.*

Proof. See Proposition 8.5.4 in [Hirschhorn, 2003], or Corollary A.2 in [Dugger, 2001b]. □

Proposition 3.4.8. *Let \mathcal{M} and \mathcal{N} be model categories, let \mathcal{M}_f be the full subcategory of fibrant objects in \mathcal{M} , and let \mathcal{N}_c be the full subcategory of cofibrant objects in \mathcal{N} .*

- *If $F : \mathcal{N} \rightarrow \mathcal{M}$ is a left Quillen functor and (Q, p) is a cofibrant replacement functor for \mathcal{N} , then (\mathcal{N}_c, Q, p) is a left deformation retract for F .*
- *If $G : \mathcal{M} \rightarrow \mathcal{N}$ is a right Quillen functor and (R, i) is a fibrant replacement functor for \mathcal{M} , then (\mathcal{M}_f, R, i) is a right deformation retract for G .*

Proof. Apply Ken Brown's lemma (3.4.5). ■

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Theorem 3.4.9. *Let \mathcal{M} and \mathcal{N} be model categories, and suppose both have fibrant and cofibrant replacement functors.*

TODO: State the version for model categories without fibrant/cofibrant replacement functors.

- (i) *If $F : \mathcal{N} \rightarrow \mathcal{M}$ is a left Quillen functor, then it has a total left derived functor $\mathbf{L}F : \text{Ho } \mathcal{N} \rightarrow \text{Ho } \mathcal{M}$ as well as a homotopical left approximation $\mathbb{L}F : \mathcal{N} \rightarrow \mathcal{M}$.*
- (ii) *If $F : \mathcal{N} \rightarrow \mathcal{M}$ and $G : \mathcal{M} \rightarrow \mathcal{L}$ are left Quillen functors, then the composite $(\mathbf{L}G)(\mathbf{L}F)$ is a total left derived functor for GF , and the composite $(\mathbb{L}G)(\mathbb{L}F)$ is a homotopical left approximation for GF .*

Dually:

- (i') *If $G : \mathcal{M} \rightarrow \mathcal{N}$ is a right Quillen functor, then it has a total right derived functor $\mathbf{R}G : \text{Ho } \mathcal{M} \rightarrow \text{Ho } \mathcal{N}$ as well as a homotopical right approximation $\mathbb{R}G : \mathcal{M} \rightarrow \mathcal{N}$.*
- (ii') *If $F : \mathcal{N} \rightarrow \mathcal{P}$ and $G : \mathcal{M} \rightarrow \mathcal{N}$ are right Quillen functors, then the composite $(\mathbf{R}F)(\mathbf{R}G)$ is a total right derived functor for FG , and the composite $(\mathbb{R}F)(\mathbb{R}G)$ is a homotopical right approximation for FG .*

Furthermore:

- (iii) *If $F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$ is a Quillen adjunction, then there is a derived adjunction $\mathbf{L}F \dashv \mathbf{L}G : \text{Ho } \mathcal{M} \rightarrow \text{Ho } \mathcal{N}$.*

Proof. Apply theorems 2.3.6 and 2.4.5 together with proposition 3.4.4. ■

Definition 3.4.10. Let \mathbb{A} be a small category and let \mathcal{M} be a model category.

- The **injective model structure** on the functor category $[\mathbb{A}, \mathcal{M}]$ is a model structure such that a morphism in $[\mathbb{A}, \mathcal{M}]$ is a cofibration (resp. weak equivalence) if and only if all its components are cofibrations (resp. weak equivalences) in \mathcal{M} .
- The **projective model structure** on the functor category $[\mathbb{A}, \mathcal{M}]$ is a model structure such that a morphism in $[\mathbb{A}, \mathcal{M}]$ is a fibration (resp. weak equivalence) if and only if all its components are fibrations (resp. weak equivalences) in \mathcal{M} .

REMARK 3.4.11. The injective (resp. projective) model structure on $[\mathbb{A}, \mathcal{M}]$ is unique if it exists, by the determination principle (theorem 3.3.7).

Proposition 3.4.12. *Let \mathcal{M} be a model category, let \mathbb{A} be a small category, and let $\Delta : \mathcal{M} \rightarrow [\mathbb{A}, \mathcal{M}]$ be the functor that sends an object X in \mathcal{M} to the constant functor $\Delta X : \mathbb{A} \rightarrow \mathcal{M}$ with value X .*

- *If \mathcal{M} has colimits for diagrams of shape \mathbb{A} , then $\Delta : \mathcal{M} \rightarrow [\mathbb{A}, \mathcal{M}]$ is a right Quillen functor with respect to the projective model structure on $[\mathbb{A}, \mathcal{M}]$ if it exists.*
- *If \mathcal{M} has limits for diagrams of shape \mathbb{A} , then $\Delta : \mathcal{M} \rightarrow [\mathbb{A}, \mathcal{M}]$ is a left Quillen functor with respect to the injective model structure on $[\mathbb{A}, \mathcal{M}]$ if it exists.*

Proof. Δ certainly preserves fibrations (resp. cofibrations) and weak equivalences with respect to the projective (resp. injective) model structure, so by [proposition 3.4.2](#), $\lim_{\rightarrow \mathbb{D}} \dashv \Delta$ (resp. $\Delta \dashv \lim_{\leftarrow \mathbb{D}}$) is a Quillen adjunction.^[5] ■

Proposition 3.4.13. *Let \mathcal{M} be a model category and let I be a set.*

- (i) *The functor category $[I, \mathcal{M}]$ admits a model structure that is simultaneously an injective model structure and a projective model structure.*
- (ii) *If \mathcal{M} has products and coproducts for families of objects indexed by I , then $\Delta : \mathcal{M} \rightarrow [I, \mathcal{M}]$ is both a left Quillen functor and a right Quillen functor.*

Proof. (i). If we declare the cofibrations (resp. weak equivalences, fibrations) in $[I, \mathcal{M}]$ to be precisely the morphisms that are cofibrations (resp. weak equivalences, fibrations) componentwise, then the axioms CM1–5 may be verified componentwise as well.

(ii). Apply [proposition 3.4.12](#). ■

3.5 Reedy diagrams

Prerequisites. §§ 3.1, 3.4

Definition 3.5.1. A **direct category** is a category C equipped with a function $\deg : \text{ob } C \rightarrow \mathbb{N}$ such that, if $f : A \rightarrow B$ is a morphism in C , then $\deg A \leq \deg B$,

[5] Recall [proposition 0.1.12](#).

with equality if and only if $f = \text{id}_A = \text{id}_B$. Dually, an **inverse category** is a category C equipped with a function $\text{deg} : \text{ob } C \rightarrow \mathbb{N}$ such that, if $f : A \rightarrow B$ is a morphism in C , then $\text{deg } A \geq \text{deg } B$, with equality if and only if $f = \text{id}_A = \text{id}_B$.

REMARK 3.5.2. The degree function for a direct or inverse category is not determined by the underlying category: for example, if deg is a degree function for C , then so is $A \mapsto 1 + \text{deg } A$. However, the partial order induced by deg is determined by the underlying category of a direct (resp. inverse) category: $\text{deg } A \leq \text{deg } B$ if and only if there exists a morphism $A \rightarrow B$ (resp. $B \rightarrow A$) in C ; note that this relation is indeed antisymmetric because the only morphisms that do not change the degree are identity morphisms.

Definition 3.5.3. A **Reedy category** is a category C equipped with two subcategories, the **direct subcategory** C_{\rightarrow} and the **inverse subcategory** C_{\leftarrow} , such that the following conditions are satisfied:

- $\text{ob } C = \text{ob } C_{\rightarrow} = \text{ob } C_{\leftarrow}$.
- There exists a function $\text{deg} : \text{ob } C \rightarrow \mathbb{N}$ such that $(C_{\rightarrow}, \text{deg})$ is a direct category and $(C_{\leftarrow}, \text{deg})$ is an inverse category.
- Every morphism in C admits a unique factorisation of the form $s \circ d$, where d is in C_{\leftarrow} and s is in C_{\rightarrow} .

A **Reedy diagram** in a category \mathcal{M} is a functor $C \rightarrow \mathcal{M}$, where C is a Reedy category.

REMARK 3.5.4. Any direct (resp. inverse) category is a Reedy category in a trivial way: take the whole category as the direct (resp. inverse) subcategory, and take $\text{disc } \text{ob } C$ as the inverse (resp. direct) subcategory.

Example 3.5.5. The simplex category $\mathbf{\Delta}$ is a Reedy category, where the direct subcategory consists of all degeneracy operators and their composites, and the inverse subcategory consists of all face operators and their composites; note that the unique factorisation condition is implied by [theorem 1.1.4](#).

REMARK 3.5.6. The opposite of any Reedy category is automatically a Reedy category, after exchanging the direct and inverse subcategories.

Definition 3.5.7. Let A be an object in a Reedy category \mathbb{C} .

- The **latching category** of \mathbb{C} at A , denoted by $\partial(\mathbb{C}_{\rightarrow} \downarrow A)$, is the largest full subcategory of the slice category $(\mathbb{C}_{\rightarrow} \downarrow A)$ that does *not* contain the object $\text{id}_A : A \rightarrow A$.
- The **matching category** of \mathbb{C} at A , denoted by $\partial(A \downarrow \mathbb{C}_{\leftarrow})$, is the largest full subcategory of the slice category $(A \downarrow \mathbb{C}_{\leftarrow})$ that does *not* contain the object $\text{id}_A : A \rightarrow A$.

REMARK 3.5.8. If \mathbb{C} is a Reedy category whose direct (resp. inverse) subcategory is discrete, then all its latching (resp. matching) categories are empty.

Definition 3.5.9. Let \mathcal{M} be a category with limits and colimits for all finite (resp. small) diagrams, and let $X : \mathbb{C} \rightarrow \mathcal{M}$ be a finite (resp. small) Reedy diagram.

- The **latching object** of X at A , denoted by $L_A(X)$, is the colimit of the diagram $\partial(\mathbb{C}_{\rightarrow} \downarrow A) \rightarrow \mathcal{M}$ obtained by composing $X : \mathbb{C} \rightarrow \mathcal{M}$ and the projection $\partial(\mathbb{C}_{\rightarrow} \downarrow A) \rightarrow \mathbb{C}$.
- The **matching object** of X at A , denoted by $M_A(X)$, is the limit of the diagram $\partial(A \downarrow \mathbb{C}_{\leftarrow}) \rightarrow \mathcal{M}$ obtained by composing $X : \mathbb{C} \rightarrow \mathcal{M}$ and the projection $\partial(A \downarrow \mathbb{C}_{\leftarrow}) \rightarrow \mathbb{C}$.
- The **latching morphism** of X at A is the morphism $L_A(X) \rightarrow XA$ induced by the inclusion $\partial(\mathbb{C}_{\rightarrow} \downarrow A) \hookrightarrow (\mathbb{C}_{\rightarrow} \downarrow A)$.
- The **matching morphism** of X at A is the morphism $XA \rightarrow M_A(X)$ induced by the inclusion $\partial(A \downarrow \mathbb{C}_{\leftarrow}) \hookrightarrow (A \downarrow \mathbb{C}_{\leftarrow})$.

REMARK 3.5.10. The latching object $L_A(X)$ is functorial in A (as A varies in the direct subcategory), and the matching object $M_A(X)$ is functorial in A (as A varies in the inverse subcategory). Of course, it goes without saying that $L_A(X)$ and $M_A(X)$ are both functorial in X (as X varies in $[\mathbb{C}, \mathcal{M}]$).

Definition 3.5.11. Let \mathcal{M} be a category with limits and colimits for all finite (resp. small) diagrams, and let $\varphi : X \Rightarrow Y$ be a natural transformation between two finite (resp. small) Reedy diagrams $X, Y : \mathbb{C} \rightarrow \mathcal{M}$.

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- The **relative latching morphism** $XA \cup^{L_A(X)} L_A(Y) \rightarrow YA$ is the unique morphism in \mathcal{M} making the diagram below commute,

$$\begin{array}{ccc}
 L_A(X) & \xrightarrow{L_A(\varphi)} & L_A(Y) \\
 \downarrow & & \downarrow \\
 XA & \longrightarrow & XA \cup^{L_A(X)} L_A(Y) \\
 & \searrow \varphi_A & \dashrightarrow \\
 & & YA
 \end{array}$$

where the arrows $L_A(X) \rightarrow XA$ and $L_A(Y) \rightarrow YA$ are the latching morphisms and the square is a pushout square.

- The **relative matching morphism** $XA \rightarrow M_A(X) \times_{M_A(Y)} YA$ is the unique morphism in \mathcal{M} making the diagram below commute,

$$\begin{array}{ccc}
 XA & \xrightarrow{\varphi_A} & YA \\
 \dashrightarrow & & \downarrow \\
 M_A(X) \times_{M_A(Y)} YA & \longrightarrow & YA \\
 \downarrow & & \downarrow \\
 M_A(X) & \xrightarrow{M_A(\varphi)} & M_A(Y)
 \end{array}$$

where the arrows $XA \rightarrow M_A(X)$ and $YA \rightarrow M_A(Y)$ are the latching morphisms and the square is a pullback square.

REMARK 3.5.12. If the direct subcategory of \mathbb{C} is discrete, then $L_A(X)$ is an initial object in \mathcal{M} for all A and X , so the relative latching morphism of a natural transformation $\varphi : X \Rightarrow Y$ at any object A in \mathbb{C} is (isomorphic to) $\varphi_A : XA \rightarrow YA$ itself.

Dually, if the inverse subcategory of \mathbb{C} is discrete, then $M_A(X)$ is a terminal object in \mathcal{M} for all A and X , so the relative matching morphism of a natural transformation $\varphi : X \Rightarrow Y$ at any object A in \mathbb{C} is (isomorphic to) $\varphi_A : XA \rightarrow YA$ itself.

Definition 3.5.13. Let \mathcal{M} be a model category, let \mathbb{C} be a finite (resp. small) Reedy category, and assume \mathcal{M} has limits and colimits for all finite (resp. small) diagrams.

- A **Reedy weak equivalence** in $[\mathbb{C}, \mathcal{M}]$ is a natural transformation such that all its components are weak equivalences in \mathcal{M} .
- A **Reedy cofibration** in $[\mathbb{C}, \mathcal{M}]$ is a natural transformation such that all its relative latching morphisms are cofibrations in \mathcal{M} .
- A **Reedy fibration** in $[\mathbb{C}, \mathcal{M}]$ is a natural transformation such that all its relative matching morphisms are fibrations in \mathcal{M} .

Proposition 3.5.14. *With notation as in the definition:*

- *A Reedy cofibration in $[\mathbb{C}, \mathcal{M}]$ is a Reedy weak equivalence if and only if all its relative latching morphisms are trivial cofibrations in \mathcal{M} .*
- *A Reedy fibration in $[\mathbb{C}, \mathcal{M}]$ is a Reedy weak equivalence if and only if all its relative matching morphisms are trivial fibrations in \mathcal{M} .*

TODO: Check if this requires functorial factorisation. Surely not!

Proof. This is Theorem 15.3.15 in [Hirschhorn, 2003]. □

Lemma 3.5.15. *Let \mathcal{M} be a model category, let \mathbb{C} be a finite (resp. small) Reedy category, and assume \mathcal{M} has limits and colimits for all finite (resp. small) diagrams.*

- *A diagram $X : \mathbb{C} \rightarrow \mathcal{M}$ is Reedy cofibrant if and only if every latching morphism of X is a cofibration in \mathcal{M} .*
- *A diagram $X : \mathbb{C} \rightarrow \mathcal{M}$ is Reedy fibrant if and only if every matching morphism of X is a fibration in \mathcal{M} .*

Proof. Let 0 be an initial object in \mathcal{M} and let 1 be a terminal object in \mathcal{M} . It is a standard fact that $\Delta 0$ is an initial object in $[\mathbb{C}, \mathcal{M}]$ and $\Delta 1$ is a terminal object in $[\mathbb{C}, \mathcal{M}]$, so the claims follow from the observation that the latching morphism $L_A(\Delta 0) \rightarrow 0$ and the matching morphism $1 \rightarrow M_A(\Delta 1)$ are isomorphisms for all objects A in \mathbb{C} . ■

Lemma 3.5.16. *With notation as in the previous lemma:*

- *If $X : \mathbb{C} \rightarrow \mathcal{M}$ is a Reedy cofibrant diagram, then, for every object A in \mathbb{C} , the object $X A$ and the latching object $L_A(X)$ are cofibrant objects in \mathcal{M} .*

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Proof. See Corollary 15.3.12 in [Hirschhorn, 2003]. ■

Theorem 3.5.17. *With notation as in the definition, the announced weak equivalences, cofibrations, and fibrations constitute a model structure on $[\mathbb{C}, \mathcal{M}]$, called the **Reedy model structure**; moreover, if \mathcal{M} is a DHK model category, then so is $[\mathbb{C}, \mathcal{M}]$ when equipped with the Reedy model structure.*

Proof. See Theorem 5.2.5 in [Hovey, 1999], or Theorem 15.3.4 in [Hirschhorn, 2003]. □

TODO: Check if this requires functorial factorisation. Surely not!

Corollary 3.5.18. *Let \mathcal{M} be a model category, let \mathbb{C} be a finite (resp. small) Reedy category, and assume \mathcal{M} has limits and colimits for all finite (resp. small) diagrams.*

- If the direct subcategory of \mathbb{C} is discrete, then the Reedy model structure on $[\mathbb{C}, \mathcal{M}]$ is the injective model structure.
- If the inverse subcategory of \mathbb{C} is discrete, then the Reedy model structure on $[\mathbb{C}, \mathcal{M}]$ is the projective model structure.

Proof. This follows from the theorem and [remark 3.5.12](#). ■

Definition 3.5.19. Let C be a Reedy category.

- C has **cofibrant constants** if, for every object A in \mathbb{C} , the latching category $\partial(C_{\rightarrow} \downarrow A)$ has at most one connected component.
- C has **fibrant constants** if, for every object A in \mathbb{C} , the matching category $\partial(A \downarrow C_{\leftarrow})$ has at most one connected component.

Example 3.5.20. Let C be a Reedy category.

- If the direct subcategory of C is discrete, then C has cofibrant constants. (In fact, every latching category is empty.)
- If the inverse subcategory of C is discrete, then C has fibrant constants. (In fact, every matching category is empty.)

Proposition 3.5.21. *Let \mathcal{M} be a model category, let \mathbb{C} be a finite (resp. small) Reedy category, and assume \mathcal{M} has limits and colimits for all finite (resp. small) diagrams.*

- *If \mathbb{C} has cofibrant constants, then the functor $\Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ is a left Quillen functor.*
- *If \mathbb{C} has fibrant constants, then the functor $\Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ is a right Quillen functor.*

Proof. The two claims are formally dual; we need only prove one version.

If the matching category $\partial(A \downarrow \mathbb{C}_\leftarrow)$ is empty, then the matching object of ΔX at A is a terminal object in \mathcal{M} , so the relative matching morphism of Δf at A is isomorphic to $f : X \rightarrow Y$ in this case.

On the other hand, if the matching category $\partial(A \downarrow \mathbb{C}_\leftarrow)$ of \mathbb{C} has only one connected component, then the matching morphism $X \rightarrow M_A(\Delta X)$ must be an isomorphism, so the relative matching morphism of Δf at A is an isomorphism, hence a (trivial) fibration in particular.

We now conclude that, for any fibration $f : X \rightarrow Y$ in \mathcal{M} , every relative matching morphism of $\Delta f : \Delta X \rightarrow \Delta Y$ is a fibration. Clearly, the functor $\Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ preserves weak equivalences, so this completes the proof that Δ is a right Quillen functor. ■

Theorem 3.5.22. *Let \mathbb{C} be a small Reedy category.*

- (i) *\mathbb{C} has cofibrant constants.*
- (ii) *$\Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ is a left Quillen functor for all DHK model categories \mathcal{M} .*
- (iii) *For every cofibrant object X in any DHK model category \mathcal{M} , the constant diagram $\Delta X : \mathbb{C} \rightarrow \mathcal{M}$ is Reedy cofibrant.*

Dually, the following are equivalent:

- (i') *\mathbb{C} has fibrant constants.*
- (ii') *$\Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ is a right Quillen functor for all DHK model categories \mathcal{M} .*
- (iii') *For every fibrant object X in any DHK model category \mathcal{M} , the constant diagram $\Delta X : \mathbb{C} \rightarrow \mathcal{M}$ is Reedy fibrant.*

Proof. (i) \Rightarrow (ii). This is the content of the earlier proposition.

(ii) \Rightarrow (iii). Left Quillen functors preserve cofibrant objects.

(iii) \Rightarrow (i). Take \mathcal{M} to be **Set** equipped with the mono–epi model structure,^[6] and consider the constant diagram $\Delta 1$. Since 1 is a cofibrant object in \mathcal{M} , $\Delta 1$ must be a Reedy cofibrant object in $[\mathbb{C}, \mathcal{M}]$. It is not hard to see that the latching object $L_A(\Delta 1)$ is the set of connected components of the latching category $\partial(\mathbb{C}_{\rightarrow} \downarrow A)$, so by lemma 3.5.15 $\partial(\mathbb{C}_{\rightarrow} \downarrow A)$ has at most one connected component. ■

Corollary 3.5.23. *Let \mathcal{M} be a DHK model category and let \mathbb{C} be a small Reedy category.*

- *If \mathbb{C} has fibrant constants, then the adjunction $\lim_{\rightarrow \mathbb{C}} \dashv \Delta : \mathcal{M} \rightarrow [\mathbb{C}, \mathcal{M}]$ is deformable.*
- *If \mathbb{C} has cofibrant constants, then the adjunction $\Delta \dashv \lim_{\leftarrow \mathbb{C}} : [\mathbb{C}, \mathcal{M}] \rightarrow \mathcal{M}$ is deformable.*

Proof. Apply theorem 3.4.9 to the above result. ■

Proposition 3.5.24. *Let X be a simplicial set, and let $\Delta^\bullet : \Delta \rightarrow \mathbf{sSet}$ be the inclusion of the standard simplices.*

- (i) *The comma category $(\Delta^\bullet \downarrow X)$ is a Reedy category, where the direct subcategory consists of all face operators and the inverse subcategory consists of all degeneracy operators.*
- (ii) *Moreover, $(\Delta^\bullet \downarrow X)$ has fibrant constants.*

Proof. (i). There is an evident projection $(\Delta^\bullet \downarrow X) \rightarrow \Delta$ that is faithful and reflects identity morphisms, so the Reedy category structure on Δ induces one on $(\Delta^\bullet \downarrow X)$.

(ii). See Proposition 15.10.4 in [Hirschhorn, 2003]. □

[6] See example 3.1.5.

3.6 Combinatorial model categories

Prerequisites. §§ 0.2, 0.4, 3.1, A.3.

Definition 3.6.1. A **cofibrantly-generated model category** is a complete and cocomplete model category \mathcal{M} such that there exist a set \mathcal{I} of cofibrations and a set \mathcal{I}' of trivial cofibrations satisfying these conditions:

- $(\mathcal{I}, \mathcal{M})$ admits the small object argument, and $\text{cof}_{\mathcal{M}} \mathcal{I}$ is the class of all cofibrations in \mathcal{M} .
- $(\mathcal{I}', \mathcal{M})$ admits the small object argument, and $\text{cof}_{\mathcal{M}} \mathcal{I}'$ is the class of all trivial cofibrations in \mathcal{M} .

REMARK 3.6.2. By Quillen's small object argument (0.4.11), any cofibrantly-generated model category satisfies axiom CM5* and thus is a DHK model category.

Theorem 3.6.3 (Kan's recognition principle). *Let \mathcal{M} be a complete and cocomplete locally small category, let \mathcal{W} be a subcategory of \mathcal{M} containing all the objects, and let \mathcal{I} and \mathcal{I}' be subsets of $\text{mor } \mathcal{M}$. Assume the following hypotheses:*

- \mathcal{W} is closed under retracts and has the 2-out-of-3 property in \mathcal{M} .
- $(\mathcal{I}, \mathcal{M})$ and $(\mathcal{I}', \mathcal{M})$ both admit the small object argument.
- $\text{inj}^{\mathcal{M}} \mathcal{I} \subseteq \mathcal{W} \cap \text{inj}^{\mathcal{M}} \mathcal{I}'$.
- $\text{cof}_{\mathcal{M}} \mathcal{I}' \subseteq \mathcal{W} \cap \text{cof}_{\mathcal{M}} \mathcal{I}$.

If, in addition, either

- $\text{inj}^{\mathcal{M}} \mathcal{I} = \mathcal{W} \cap \text{inj}^{\mathcal{M}} \mathcal{I}'$, or
- $\text{cof}_{\mathcal{M}} \mathcal{I}' = \mathcal{W} \cap \text{cof}_{\mathcal{M}} \mathcal{I}$.

then there exists a unique model structure on \mathcal{M} such that $\text{cof}_{\mathcal{M}} \mathcal{I}$ is the class of cofibrations, $\text{cof}_{\mathcal{M}} \mathcal{I}'$ is the class of trivial cofibrations, and \mathcal{W} is the class of weak equivalences.

Proof. See Theorem 11.3.1 in [Hirschhorn, 2003]. □

Theorem 3.6.4 (Kan's lifting theorem). *Let \mathcal{M} be a complete and cocomplete locally small category, let \mathcal{N} be a cofibrantly generated model category. Assume the following hypotheses:*

- $F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$ is an adjunction of categories.
- \mathcal{J} is a generating set of cofibrations in \mathcal{N} .
- \mathcal{J}' is a generating set of trivial cofibrations in \mathcal{N} .
- $(\mathcal{I}, \mathcal{M})$ and $(\mathcal{I}', \mathcal{M})$ admit the small object argument, where \mathcal{I} and \mathcal{I}' are the following sets:

$$\begin{aligned}\mathcal{I} &= \{Ff \mid f \in \mathcal{J}\} \\ \mathcal{I}' &= \{Ff \mid f \in \mathcal{J}'\}\end{aligned}$$

- G sends relative \mathcal{I}' -cell complexes in \mathcal{M} to weak equivalences in \mathcal{N} .

Then:

- (i) *There is a unique model structure on \mathcal{M} with $\text{cof}_{\mathcal{M}} \mathcal{I}$ as the class of cofibrations and $\text{cof}_{\mathcal{M}} \mathcal{I}'$ as the class of trivial cofibrations.*
- (ii) *A morphism $g : A \rightarrow B$ in \mathcal{M} is a weak equivalence in this model structure if and only if $Gg : GA \rightarrow GB$ is a weak equivalence in \mathcal{N} .*
- (iii) *$F \dashv G : \mathcal{M} \rightarrow \mathcal{N}$ is a Quillen adjunction with respect to this model structure.*

Proof. See Theorem II.3.2 in [Hirschhorn, 2003]. □

Theorem 3.6.5 (Existence of cofibrantly-generated projective model structures). *Let \mathcal{M} be a cofibrantly-generated model category. If \mathbb{A} is a small category, then the projective model structure on $[\mathbb{A}, \mathcal{M}]$ exists and is cofibrantly generated.*

Proof. See Theorem II.6.1 in [Hirschhorn, 2003]. □

Definition 3.6.6. A **combinatorial model category** is a cofibrantly-generated model category that is also a locally presentable category.

REMARK 3.6.7. Since locally presentable categories are automatically complete and cocomplete,^[7] in light of [remark 0.4.9](#), to show that a locally presentable model category \mathcal{M} is a combinatorial model category, it is enough to verify that there exist sets \mathcal{I} and \mathcal{I}' such that $\text{cof}_{\mathcal{M}} \mathcal{I}$ is the class of all cofibrations in \mathcal{M} and $\text{cof}_{\mathcal{M}} \mathcal{I}'$ is the class of all trivial cofibrations in \mathcal{M} .

Theorem 3.6.8 (Existence of combinatorial injective model structures). *Let \mathcal{M} be a combinatorial model category. If \mathbb{A} is a small category, then the injective model structure on $[\mathbb{A}, \mathcal{M}]$ exists and is combinatorial.*

Proof. This theorem is due to Lurie; see [\[HTT, Proposition A.2.8.2\]](#). □

3.7 Monoidal model categories

Prerequisites. §§ [3.1](#), [3.4](#), [B.1](#), [B.2](#).

Proposition 3.7.1. *Let \mathcal{C} and \mathcal{D} be categories with pullbacks, let \mathcal{E} be a category with pushouts, and let $\mathcal{I} \subseteq \text{mor } \mathcal{C}$, $\mathcal{J} \subseteq \text{mor } \mathcal{D}$ and $\mathcal{K} \subseteq \text{mor } \mathcal{E}$ be subensembles. Suppose we have the following functors*

$$\begin{aligned} \otimes &: \mathcal{C} \times \mathcal{D} \rightarrow \mathcal{E} \\ \pitchfork &: \mathcal{D}^{\text{op}} \times \mathcal{E} \rightarrow \mathcal{C} \\ \circ - &: \mathcal{E} \times \mathcal{C}^{\text{op}} \rightarrow \mathcal{D} \end{aligned}$$

and natural bijections:

$$\begin{aligned} \mathcal{E}(\mathcal{C} \otimes \mathcal{D}, \mathcal{E}) &\cong \mathcal{C}(\mathcal{C}, \mathcal{D} \pitchfork \mathcal{E}) \\ \mathcal{E}(\mathcal{C} \otimes \mathcal{D}, \mathcal{E}) &\cong \mathcal{D}(\mathcal{D}, \mathcal{E} \circ - \mathcal{C}) \\ \mathcal{C}(\mathcal{C}, \mathcal{D} \pitchfork \mathcal{E}) &\cong \mathcal{D}(\mathcal{D}, \mathcal{E} \circ - \mathcal{C}) \end{aligned}$$

Then the following are equivalent:

[7] See [theorem 0.2.26](#).

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- (i) If $f : C \rightarrow C'$ is in \mathcal{I} , $g : D \rightarrow D'$ is in \mathcal{J} , and the square in the diagram below is a pushout square in \mathcal{E} ,

$$\begin{array}{ccc}
 C \otimes D & \xrightarrow{\text{id}_C \otimes g} & C \otimes D' \\
 f \otimes \text{id}_D \downarrow & & \downarrow \\
 C' \otimes D & \xrightarrow{\quad} & (C \otimes D') \cup^{C \otimes D} (C' \otimes D) \\
 & \searrow \text{id}_{C'} \otimes g & \downarrow f \otimes \text{id}_{D'} \\
 & & C' \otimes D'
 \end{array}$$

$\xrightarrow{\quad} (C \otimes D') \cup^{C \otimes D} (C' \otimes D) \xrightarrow{\quad} C' \otimes D'$

$\xrightarrow{\quad} (C \otimes D') \cup^{C \otimes D} (C' \otimes D) \xrightarrow{\quad} C' \otimes D'$

$\xrightarrow{\quad} (C \otimes D') \cup^{C \otimes D} (C' \otimes D) \xrightarrow{\quad} C' \otimes D'$

then the unique morphism $f \boxplus g$ making the diagram commute is in \mathcal{K} .

- (ii) If $g : D \rightarrow D'$ is in \mathcal{J} , $h : E \rightarrow E'$ is in \mathcal{K} , and the square in the diagram below is a pullback square in \mathcal{C} ,

$$\begin{array}{ccc}
 D' \pitchfork E & \xrightarrow{g \pitchfork h} & D \pitchfork E \\
 \downarrow \text{id}_{D'} \pitchfork h & & \downarrow \text{id}_D \pitchfork h \\
 D' \pitchfork E' & \xrightarrow{g \pitchfork h} & D \pitchfork E' \\
 \downarrow \text{id}_{D'} \pitchfork h & & \downarrow \text{id}_D \pitchfork h \\
 D' \pitchfork E' & \xrightarrow{g \pitchfork h} & D \pitchfork E'
 \end{array}$$

$\xrightarrow{\quad} (D' \pitchfork E') \times_{D \pitchfork E'} (D \pitchfork E) \xrightarrow{\quad} D \pitchfork E$

$\xrightarrow{\quad} (D' \pitchfork E') \times_{D \pitchfork E'} (D \pitchfork E) \xrightarrow{\quad} D \pitchfork E$

$\xrightarrow{\quad} (D' \pitchfork E') \times_{D \pitchfork E'} (D \pitchfork E) \xrightarrow{\quad} D \pitchfork E$

then the unique morphism $g \boxplus h$ making the diagram commute is in \mathcal{I} .

- (iii) If $h : E \rightarrow E'$ is in \mathcal{K} , $f : C \rightarrow C'$ is in \mathcal{I} and the square in the diagram below is a pullback square in \mathcal{D} ,

$$\begin{array}{ccc}
 E \circ C' & \xrightarrow{h \circ f} & E \circ C \\
 \downarrow h \circ \text{id}_C & & \downarrow h \circ \text{id}_C \\
 E' \circ C' & \xrightarrow{h \circ f} & E' \circ C \\
 \downarrow h \circ \text{id}_C & & \downarrow h \circ \text{id}_C \\
 E' \circ C' & \xrightarrow{h \circ f} & E' \circ C
 \end{array}$$

$\xrightarrow{\quad} (E' \circ C') \times_{E' \circ C} (E \circ C) \xrightarrow{\quad} E \circ C$

$\xrightarrow{\quad} (E' \circ C') \times_{E' \circ C} (E \circ C) \xrightarrow{\quad} E \circ C$

$\xrightarrow{\quad} (E' \circ C') \times_{E' \circ C} (E \circ C) \xrightarrow{\quad} E \circ C$

then the unique morphism $h \boxplus f$ making the diagram commute is in \mathcal{J} .

Proof. (i) \Rightarrow (ii). Let $f : C \rightarrow C'$ be in \mathcal{I} , let $g : D \rightarrow D'$ be in \mathcal{J} , let $h : E \rightarrow E'$ be in \mathcal{K} , and suppose we have a commutative diagram of the following form:

$$\begin{array}{ccc} C & \longrightarrow & D' \pitchfork E \\ f \downarrow & \nearrow \text{dashed} & \downarrow g \boxplus h \\ C' & \longrightarrow & (D' \pitchfork E') \times_{D \pitchfork E'} (D \pitchfork E) \end{array}$$

By the universal property of pullbacks, this corresponds to a commutative diagram in \mathcal{C} of the form below,

$$\begin{array}{ccccc} C & \longrightarrow & D' \pitchfork E & & \\ f \downarrow & \nearrow \text{dashed} & \downarrow \text{id}_{D'} \pitchfork h & & \\ C' & \longrightarrow & D' \pitchfork E' & & \\ & \searrow & \downarrow g \pitchfork \text{id}_{E'} & & \\ & & D \pitchfork E & \xrightarrow{\text{id}_D \pitchfork h} & D \pitchfork E' \end{array}$$

and, by adjoint transposition, to a commutative diagram in \mathcal{E} of the form

$$\begin{array}{ccccc} C \otimes D & \xrightarrow{\text{id}_C \otimes g} & C \otimes D' & & \\ f \otimes \text{id}_D \downarrow & & \downarrow f \otimes \text{id}_{D'} & & \\ C' \otimes D & \xrightarrow{\text{id}_{C'} \otimes g} & C' \otimes D' & \longrightarrow & E \\ & \searrow & \downarrow \text{dashed} & & \downarrow h \\ & & C' \otimes D' & \longrightarrow & E' \end{array}$$

whence, by the universal property of pushouts, commutative diagram in \mathcal{E} of the following form:

$$\begin{array}{ccc} (C \otimes D') \cup^{C \otimes D} (C' \otimes D) & \longrightarrow & E \\ f \boxplus g \downarrow & \nearrow \text{dashed} & \downarrow h \\ C' \otimes D' & \longrightarrow & E' \end{array}$$

But $(f \boxplus g) \boxtimes h$, so we conclude that $f \boxtimes (g \boxplus h)$.

(ii) \Rightarrow (iii), (i) \Rightarrow (ii). A similar argument works. ■

Definition 3.7.2. Let \mathcal{C} , \mathcal{D} , and \mathcal{E} be three model categories. A **Quillen adjunction of two variables** consists of three functors $\otimes, \pitchfork, \circ$ with natural bijections as in the proposition satisfying the following (equivalent) axioms:

- (a) If $h : E \rightarrow E'$ is a fibration in \mathcal{E} and $f : C \rightarrow C'$ is a cofibration in \mathcal{C} , then the morphism $h \boxtimes f : E \circ C' \rightarrow (E' \circ C') \times_{E' \circ C} (E \circ C)$ is a fibration in \mathcal{D} , which is a weak equivalence if either h or f is.
- (b) If $f : C \rightarrow C'$ is a cofibration in \mathcal{C} and $g : D \rightarrow D'$ is a cofibration in \mathcal{D} , then the morphism $f \boxtimes g : C \otimes D \rightarrow (C \otimes D') \cup^{C \otimes D} (C' \otimes D)$ is a cofibration in \mathcal{E} , which is a weak equivalence if either f or g is.
- (c) If $g : D \rightarrow D'$ is a cofibration in \mathcal{C} and $h : E \rightarrow E'$ is a fibration in \mathcal{D} , then the morphism $g \boxtimes h : D' \pitchfork E \rightarrow (D' \pitchfork E') \times_{D \pitchfork E'} (D \pitchfork E)$ is a fibration in \mathcal{C} , which is a weak equivalence if either g or h is.

Proposition 3.7.3. Let $(\otimes, \pitchfork, \circ)$ be a Quillen adjunction of two variables as above.

- (i) For each cofibrant object C in \mathcal{C} , the adjunction

$$C \otimes (-) \dashv (-) \circ C : \mathcal{E} \rightarrow \mathcal{D}$$

is a Quillen adjunction.

- (ii) For each cofibrant object D in \mathcal{D} , the adjunction

$$(-) \otimes D \dashv D \pitchfork (-) : \mathcal{E} \rightarrow \mathcal{C}$$

is a Quillen adjunction.

- (iii) For each fibrant object E in \mathcal{E} , the adjunction

$$E \circ (-) \dashv (-) \pitchfork E : \mathcal{D}^{\text{op}} \rightarrow \mathcal{C}$$

is a Quillen adjunction.

Proof. Immediate from the definitions. ◆

Corollary 3.7.4.

- (i) For each object C in \mathcal{C} , $C \otimes (-)$ preserves weak equivalences between cofibrant objects, and $(-) \otimes C$ preserves weak equivalences between fibrant objects.

- (ii) For each object D in \mathcal{D} , $(-) \otimes D$ preserves weak equivalences between cofibrant objects, and $D \pitchfork (-)$ preserves weak equivalences between fibrant objects.
- (iii) For each object E in \mathcal{E} , $E \circ (-)$ sends weak equivalences between cofibrant objects in \mathcal{C} to weak equivalences between fibrant objects in \mathcal{D} , and $(-) \pitchfork E$ sends weak equivalences between cofibrant objects in \mathcal{D} to weak equivalences between fibrant objects in \mathcal{D} .

Proof. Apply Ken Brown's lemma (3.4.5). ■

Lemma 3.7.5. *Let \mathcal{V} be a monoidal category, let \mathcal{M} be a model category with fibrant and cofibrant replacement functors, and let $p : \tilde{I} \rightarrow I$ be a morphism in \mathcal{V} , where I is the monoidal unit of \mathcal{V} .*

If \mathcal{M} has a left \mathcal{V} -action \otimes and right adjoint right \mathcal{V}^{op} -action $\circ-$ such that the adjunction

$$\tilde{I} \otimes (-) \dashv (-) \circ \tilde{I} : \mathcal{M} \rightarrow \mathcal{M}$$

is a Quillen adjunction, then the following are equivalent:

- (i) *For all cofibrant objects X in \mathcal{M} , $p \otimes \text{id}_X : \tilde{I} \otimes X \rightarrow I \otimes X$ is a weak equivalence.*
- (ii) *For all fibrant objects Y in \mathcal{M} , $\text{id}_Y \circ p : Y \circ I \rightarrow Y \circ \tilde{I}$ is a weak equivalence.*

If \mathcal{M} has a right \mathcal{V} -action \otimes and a right adjoint left \mathcal{V}^{op} -action $\rightarrow-$ such that the adjunction

$$(-) \otimes \tilde{I} \dashv \tilde{I} \rightarrow (-) : \mathcal{M} \rightarrow \mathcal{M}$$

is a Quillen adjunction, then the following are equivalent:

- (i') *For all cofibrant objects X in \mathcal{M} , $\text{id}_X \otimes p : X \otimes \tilde{I} \rightarrow X \otimes I$ is a weak equivalence.*
- (ii') *For all fibrant objects Y in \mathcal{M} , $p \rightarrow \text{id}_Y : I \rightarrow Y \rightarrow \tilde{I} \rightarrow Y$ is a weak equivalence.*

Proof. Since $\eta_X : X \rightarrow I \otimes X$ is a natural isomorphism, the adjunction

$$I \otimes (-) \dashv (-) \circ I : \mathcal{M} \rightarrow \mathcal{M}$$

is an adjoint equivalence of categories, and *a fortiori* a Quillen equivalence, and the natural transformations $p \circlearrowleft (-)$ and $(-) \circlearrowright p$ constitute a conjugate pair. [Theorem 2.3.11](#) says that the derived natural transformations for $p \circlearrowleft (-)$ and $(-) \circlearrowright p$ constitute a conjugate pair of natural transformations between the derived adjunctions. Applying [proposition 2.3.12](#) to [theorem 3.4.9](#), we deduce that the following are equivalent:

- For all cofibrant objects X , $p \circlearrowleft \text{id}_X$ is a weak equivalence.
- The left derived natural transformation for $p \circlearrowleft (-)$ is a natural isomorphism.
- The right derived natural transformation for $(-) \circlearrowright p$ is a natural isomorphism.
- For all fibrant objects Y , $\text{id}_Y \circlearrowright p$ is a weak equivalence. ■

Definition 3.7.6. A **monoidal model category** is a biclosed monoidal category \mathcal{M} equipped with a model structure satisfying the following additional axioms:

- **Pushout–product axiom.** The right \mathcal{M} -hom system $(\otimes, \dashv, \circlearrowright)$, where \dashv (resp. \circlearrowright) is the right (resp. left) internal hom functor of \mathcal{M} , is a Quillen adjunction of two variables.
- **Unit axiom.** For each cofibrant replacement (\tilde{I}, p) of the monoidal unit I and each cofibrant object X in \mathcal{M} , the morphisms $p \otimes \text{id}_X : \tilde{I} \otimes X \rightarrow I \otimes X$ and $\text{id}_X \otimes p : X \otimes \tilde{I} \rightarrow X \otimes I$ are weak equivalences in \mathcal{M} .

Lemma 3.7.7. *Let \mathcal{M} be a biclosed monoidal category equipped with a model structure satisfying the pushout–product axiom, and let X be any object in \mathcal{M} . The following are equivalent:*

- (i) *There exists a cofibrant replacement (\tilde{I}, p) of the monoidal unit I such that $p \otimes \text{id}_X$ and $\text{id}_X \otimes p$ are weak equivalences in \mathcal{M} .*
- (ii) *There exists a fibrant cofibrant replacement (QI, q) of the monoidal unit I such that $q \otimes \text{id}_X$ and $\text{id}_X \otimes q$ are weak equivalences in \mathcal{M} .*
- (iii) *For any cofibrant replacement (\tilde{I}, p) of the monoidal unit I , both $p \otimes \text{id}_X$ and $\text{id}_X \otimes p$ are weak equivalences in \mathcal{M} .*

Proof. (i) \Rightarrow (ii). Let (QI, q) be a fibrant cofibrant replacement of I ; such exists by [proposition 3.1.14](#). Since \tilde{I} is cofibrant, axiom CM5 implies there is a morphism $w : \tilde{I} \rightarrow QI$ such that $q \circ w = p$, and the 2-out-of-3 property implies w is a weak equivalence. [Corollary 3.7.4](#) says $w \otimes \text{id}_X$ and $\text{id}_X \otimes w$ are weak equivalences, thus by the 2-out-of-3 property again $q \otimes \text{id}_X$ and $\text{id}_X \otimes q$ must be weak equivalences.

(ii) \Rightarrow (iii). A similar argument works.

(iii) \Rightarrow (i). Obvious, given the existence of cofibrant replacements. ■

Corollary 3.7.8. *Let \mathcal{M} be a biclosed monoidal category equipped with a model structure. If the monoidal unit I is a cofibrant object in \mathcal{M} , then the following are equivalent:*

- (i) \mathcal{M} is a monoidal model category.
- (ii) \mathcal{M} satisfies the pushout–product axiom. ■

Definition 3.7.9. A **cartesian model category** is a cartesian closed category \mathcal{M} equipped with a model structure satisfying the following additional axioms:

- **Pushout–product axiom.** The left \mathcal{M} -hom system $(\times, [-, -], [-, -])$ is a Quillen adjunction of two variables.
- **Cofibrant unit axiom.** Every terminal object in \mathcal{M} is cofibrant.

Example 3.7.10. The Kan–Quillen model structure on **sSet** makes it a cartesian model category: **sSet** is a cartesian closed combinatorial model category (*a fortiori* a DHK model category), all simplicial sets are cofibrant, and the pushout–product axiom is just [proposition 1.3.9](#).

Definition 3.7.11. An **isofibration** is a functor that is injective on objects. An **isofibration** is a functor $F : C \rightarrow D$ such that, for every object C in C and every isomorphism $f : FC \rightarrow D$ in D , there exists an isomorphism $\tilde{f} : C \rightarrow \tilde{D}$ in C such that $F\tilde{f} = f$.

Proposition 3.7.12. *Let **Cat** be the category of small categories. The following data constitute a model structure on **Cat**:*

- *The weak equivalences are the functors that are fully faithful and essentially surjective on objects.*

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- *The cofibrations are the isofibrations.*
- *The fibrations are the isofibrations.*

Moreover, the factorisations for axiom CM5 may be chosen functorially, so that **Cat** becomes a DHK model category. This model structure is called the **canonical model structure on Cat**.

Proof. It is not hard to show that **Cat** has limits and colimits for all small diagrams, so axiom CM1* is satisfied. It is also clear that the announced class of weak equivalences has the 2-out-of-3 property, so by [theorem 3.1.8](#), it is enough to show that we have a pair of compatible weak factorisation systems.

Let $I : \mathbb{A} \rightarrow \mathbb{B}$ be an isofibration and $P : \mathbb{C} \rightarrow \mathbb{D}$ be an isofibration, and suppose we have a commutative diagram of the following form:

$$\begin{array}{ccc} \mathbb{A} & \xrightarrow{F} & \mathbb{C} \\ I \downarrow & & \downarrow P \\ \mathbb{B} & \xrightarrow{G} & \mathbb{D} \end{array}$$

First, suppose P is a weak equivalence. Then, P must be surjective on objects, so we may define a map $H : \text{ob } \mathbb{B} \rightarrow \text{ob } \mathbb{C}$ by taking $HB = FA$ if $B = IA$ for some A , and if B is not in the image of A , define HB to be any object in \mathbb{C} such that $PHB = GB$; there is then a unique way of extending H to a functor $\mathbb{B} \rightarrow \mathbb{C}$ making the evident diagram commute.

Next, instead suppose I is a weak equivalence. Then, I may be regarded as the inclusion of a full subcategory that is essentially surjective on objects. For each object B in \mathbb{B} that is not in the image of I , fix an object A in \mathbb{A} and an isomorphism $IA \xrightarrow{\cong} B$. Since P is an isofibration, for each such B we may also choose an object C in \mathbb{C} and an isomorphism $FA \xrightarrow{\cong} C$ whose image under P is $GIA \xrightarrow{\cong} GB$. There is then a unique functor $H : \mathbb{B} \rightarrow \mathbb{C}$ that makes the evident diagram commute and sends B to the chosen C and $IA \xrightarrow{\cong} B$ to $FA \xrightarrow{\cong} C$.

It remains to be shown that every functor can be factorised in the required manner. Let $F : \mathbb{C} \rightarrow \mathbb{D}$ be any functor. Consider the iso-comma category $(F \downarrow \mathbb{D})_{\text{iso}}$:

- The objects are triples (C, D, α) , where C is an object in \mathbb{C} , D is an object in \mathbb{D} , and $\alpha : FC \rightarrow D$ is an *isomorphism* in \mathbb{D} .

- The morphisms $(C, D, \alpha) \rightarrow (C', D', \alpha')$ is a morphism $f : C \rightarrow C'$ in \mathbb{C} together with a morphism $g : D \rightarrow D'$ in \mathbb{D} such that $g \circ \alpha = \alpha' \circ Ff$.^[8]
- Composition and identities are inherited from \mathbb{C} and \mathbb{D} .

There is an evident isofibration $I : \mathbb{C} \rightarrow (F \downarrow \mathbb{D})_{\text{iso}}$ sending an object C in \mathbb{C} to the object (C, FC, id_{FC}) , and it is easy to see that I is a weak equivalence. On the other hand, the projection $P : (F \downarrow \mathbb{D})_{\text{iso}} \rightarrow \mathbb{D}$ is an isofibration by construction, and obviously $F = PI$. Thus, we have factored F as a trivial isofibration followed by an isofibration, and it is clear that this construction is functorial in F .

Now, consider instead the category $\mathbf{M}(F)$ defined below:

- $\text{ob } \mathbf{M}(F) = \text{ob } \mathbb{C} \amalg \text{ob } \mathbb{D}$.
- If C and C' are objects in \mathbb{C} , while D and D' are objects in \mathbb{D} , then:

$$\text{Hom}(C, C') = \mathbb{D}(FC, FC')$$

$$\text{Hom}(C, D') = \mathbb{D}(FC, D')$$

$$\text{Hom}(D, C') = \mathbb{D}(D, FC')$$

$$\text{Hom}(D, D') = \mathbb{D}(D, D')$$

- Composition and identities are inherited from \mathbb{D} .

There is an evident isofibration $I : \mathbb{C} \rightarrow \mathbf{M}(F)$ that sends an object C in \mathbb{C} to the corresponding object in $\mathbf{M}(F)$ and sends a morphism $f : C \rightarrow C'$ in \mathbb{C} to the morphism in $\mathbf{M}(F)$ corresponding to $Ff : FC \rightarrow FC'$ in \mathbb{D} . On the other hand, there is an evident projection $P : \mathbf{M}(F) \rightarrow \mathbb{D}$ that is fully faithful and surjective on objects, i.e. P is a trivial isofibration. Of course, $F = PI$, so this is a factorisation of F as an isofibration followed by a trivial isofibration, and it is clear that this construction is functorial in F . ■

Theorem 3.7.13. *Let \mathbf{Cat} be considered as a model category via the canonical model structure.*

- (i) *Every object in \mathbf{Cat} is both cofibrant and fibrant.*
- (ii) *\mathbf{Cat} is a combinatorial model category.*

[8] However, because α and α' are isomorphisms, f freely and uniquely determines g .

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(iii) **Cat** is a cartesian model category.

Proof. (i). The unique functor $\emptyset \rightarrow \mathbb{C}$ is vacuously an isofibration, and the unique functor $\mathbb{C} \rightarrow \mathbb{1}$ is certainly an isofibration.

(ii). **Cat** is a locally finitely presentable category,^[9] and it remains to be shown that the canonical model structure is a cofibrantly-generated model structure.

By the very definition of isofibration, the set $\{\mathbb{1} \rightarrow \mathbf{I2}\}$ is a generating set of trivial isofibrations, where $\mathbf{I2}$ is the groupoid containing only a pair of non-trivial isomorphisms. It is also straightforward to see that a functor is ...

... surjective on objects if and only if it has the right lifting property with respect to the unique functor $\emptyset \rightarrow \mathbb{1}$;

... full if and only if it has the right lifting property with respect to the inclusion $\text{disc } 2 \rightarrow \mathbb{2}$; and

... faithful if and only if it has the right lifting property with respect to the surjective functor $\mathbb{E} \rightarrow \mathbb{2}$, where \mathbb{E} is the category with a parallel pair of non-trivial morphisms.

However, a functor is a trivial isofibration if and only if it is fully faithful and surjective on objects, so $\{\emptyset \rightarrow \mathbb{1}, \text{disc } 2 \rightarrow \mathbb{2}, \mathbb{E} \rightarrow \mathbb{2}\}$ is a set of generating isofibrations.

(iii). Let $F : \mathbb{C} \rightarrow \mathbb{C}'$ and $G : \mathbb{D} \rightarrow \mathbb{D}'$ be isofibrations, and consider the functor $F \boxtimes G$ defined by the diagram below:

$$\begin{array}{ccc}
 \mathbb{C} \times \mathbb{D} & \xrightarrow{\text{id}_{\mathbb{C}} \times G} & \mathbb{C} \times \mathbb{D}' \\
 \downarrow F \times \text{id}_{\mathbb{D}} & & \downarrow F \times \text{id}_{\mathbb{D}'} \\
 \mathbb{C}' \otimes \mathbb{D} & \xrightarrow{\quad} & (\mathbb{C} \times \mathbb{D}') \cup^{\mathbb{C} \times \mathbb{D}} (\mathbb{C}' \times \mathbb{D}) \\
 & \searrow \text{id}_{\mathbb{C}'} \times G & \downarrow F \boxtimes G \\
 & & \mathbb{C}' \times \mathbb{D}'
 \end{array}$$

The functor $\text{ob} : \mathbf{Cat} \rightarrow \mathbf{Set}$ has both left and right adjoints, so it is easy to see that $F \boxtimes G$ is an isofibration. Moreover, if $F : \mathbb{C} \rightarrow \mathbb{C}'$ is a trivial

[9] — because e.g. **Cat** is the category of models for a finite limit sketch; see Proposition 1.51 in [LPAC] or Proposition 5.6.4 in [Borceux, 1994b].

isofibration, one may directly verify that $F \times \text{id}_{\mathbb{D}} : \mathbb{C} \times \mathbb{D} \rightarrow \mathbb{C}' \times \mathbb{D}$ and $F \times \text{id}_{\mathbb{D}'} : \mathbb{C} \times \mathbb{D}' \rightarrow \mathbb{C}' \times \mathbb{D}'$ are trivial isofibrations; but trivial isofibrations are closed under pushout, so applying the 2-out-of-3 property of weak equivalences, we conclude that $F \sqsupseteq G$ is a trivial isofibration if F is. The symmetrical argument shows that $F \sqsupseteq G$ is a trivial isofibration if G is.

Having shown that **Cat** satisfies the pushout–product axiom, we must now verify that **Cat** is cartesian closed and has a cofibrant unit; but the former is a very well-known fact, and the latter follows from claim (i). ■

Theorem 3.7.14. *Let **Grpd** be the category of small groupoids.*

(i) *The following data constitute a model structure on **Grpd**:*

- *The weak equivalences are the functors that are fully faithful and essentially surjective on objects.*
- *The cofibrations are the isofibrations.*
- *The fibrations are the isofibrations.*

*This model structure is called the **canonical model structure on Grpd**.*

(ii) *Every object in **Grpd** is both cofibrant and fibrant.*

(iii) ***Grpd** is a combinatorial model category.*

(iv) ***Grpd** is a cartesian model category.*

(v) *The inclusion $\text{und} : \mathbf{Grpd} \rightarrow \mathbf{Cat}$ preserves and reflects weak equivalences, isofibrations, and isofibrations; moreover, it is both a left Quillen functor and a right Quillen functor.*

Proof. (i). The proof of [proposition 3.7.12](#) goes through for **Grpd** without modifications.

(ii) – (iv). These can be proven in essentially the same way as [proposition 3.7.12](#), though one should note that the generating isofibrations and generating trivial isofibrations for **Grpd** are different.

(v). It is clear that $\text{und} : \mathbf{Grpd} \rightarrow \mathbf{Cat}$ has the announced preservation and reflection properties. One may check that und has a left adjoint $\mathbf{I} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$ and a right adjoint $\text{iso} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$, so und is both a left Quillen functor and a right Quillen functor. ■

DERIVATORS

4.1 Basics

Prerequisites. §§ 2.1, 2.5, A.1, A.5.

The notion of derivator has a somewhat complicated history; the name and the original idea are due to Grothendieck [1983, 1991], but Heller [1988] studied essentially the same thing independently. The distinguishing characteristic of the theory of derivators is its agnosticism: a derivator is a way of studying homotopy-coherent diagrams and their limits/colimits without using any particular model for homotopical algebra.

In this section, we use the explicit universe convention, and all 2-categories and 2-functors will be strict unless otherwise stated, and for simplicity, we say ‘coproduct’, ‘product’, ‘pullback’, etc. instead of ‘2-coproduct’, ‘2-product’, ‘2-pullback’ etc., and we tacitly assume that these have the relevant 2-dimensional universal property in addition to the usual 1-dimensional universal property.

Definition 4.1.1. A **derivator domain** is 2-category \mathfrak{K} satisfying these axioms:

D0. \mathfrak{K} has an initial object 0 , a terminal object 1 , and tensors with the category $\mathbb{2} = \{0 \rightarrow 1\}$.

D1. \mathfrak{K} has finite coproducts and pullbacks.

D2. \mathfrak{K} has comma objects of the form $(u \downarrow b)$ and $(b \downarrow u)$ for all morphisms $u : A \rightarrow B$ and $b : 1 \rightarrow B$.

A **subdomain** of a derivator domain is a 2-full 2-subcategory that is closed under constructions specified in the above axioms.

Definition 4.1.2. Let \mathbf{U} be a universe. A **U-small prederivator** on \mathfrak{K} is a 2-functor $\mathcal{D} : \mathfrak{K}^{\text{op}} \rightarrow \mathbf{Cat}$, where \mathfrak{K} is a derivator domain and \mathbf{Cat} is the 2-category of \mathbf{U} -small categories. A **prederivator** is a 2-functor that is a \mathbf{U} -small prederivator for some universe \mathbf{U} .

We write \mathcal{D}^A for the value of \mathcal{D} at an object A in \mathfrak{K} , and we write either \mathcal{D}^u or u^* for the functor $\mathcal{D}^B \rightarrow \mathcal{D}^A$ induced by a morphism $u : A \rightarrow B$ in \mathfrak{K} . If $f : x \rightarrow y$ is a morphism in \mathcal{D}^B , then we may sometimes write $f \upharpoonright u : x \upharpoonright u \rightarrow y \upharpoonright u$ instead of $u^*(f) : u^*(x) \rightarrow u^*(y)$. The **underlying category** of a prederivator \mathcal{D} is the category \mathcal{D}^1 , where 1 is any terminal object of \mathfrak{K} .

REMARK 4.1.3. While it is true that \mathfrak{K} is a derivator domain if and only if \mathfrak{K}^{co} is a derivator domain, the duality principle for general prederivators is somewhat subtle: because $(-)^{\text{op}}$ is a 2-functor $\mathbf{Cat}^{\text{co}} \rightarrow \mathbf{Cat}$, the opposite of a prederivator on \mathfrak{K} is a prederivator on \mathfrak{K}^{co} , which is in general not isomorphic or even equivalent to \mathfrak{K} .

One should be aware that some authors (e.g. Cisinski [2003]) define prederivators to be 2-functors $\mathfrak{K}^{\text{coop}} \rightarrow \mathbf{Cat}$; readers should take care to dualise results appropriately when translating between the two conventions.

Definition 4.1.4. A **semiderivator** on \mathfrak{K} is prederivator $\mathcal{D} : \mathfrak{K}^{\text{op}} \rightarrow \mathbf{Cat}$ satisfying the following axioms:

Der1. \mathcal{D} sends coproducts of finite families of objects in \mathfrak{K} to products in \mathbf{Cat} .

Der2. Let A be an object in \mathfrak{K} and let $f : x \rightarrow y$ be a morphism in \mathcal{D}^A . Then, f is an isomorphism in \mathcal{D}^A if and only if, for all morphisms $a : 1 \rightarrow A$ in \mathfrak{K} , the morphism $f \upharpoonright a : x \upharpoonright a \rightarrow y \upharpoonright a$ is an isomorphism in \mathcal{D}^1 .

Example 4.1.5. If B is an object in \mathfrak{K} and \mathfrak{K} is a locally \mathbf{U} -small 2-category, then the 2-functor $\mathfrak{K}(-, B) : \mathfrak{K}^{\text{op}} \rightarrow \mathbf{Cat}$ is a prederivator. We say $\mathfrak{K}(-, B)$ is the **prederivator represented by B** .

Definition 4.1.6. Let C be a \mathbf{U} -small relative category. The **prederivator** of C , denoted by $\mathcal{D}(C)$, is the \mathbf{U} -small prederivator on \mathbf{RelCat} (or any subdomain thereof) defined by $\mathcal{D}(C)^A = \text{Ho}[\mathcal{A}, C]_h$.

Proposition 4.1.7. Let \mathcal{D} be a prederivator on \mathfrak{K} . If A is an object in \mathfrak{K} and C is a finite category, then there is a canonical comparison functor $\mathcal{D}^{\text{co}A} \rightarrow [C, \mathcal{D}^A]$ and it is 2-natural in A and in C .

TODO: Check 2-naturality properly!

Proof. By definition, the object $\mathbb{C} \odot A$ in \mathfrak{K} induces isomorphisms

$$\mathfrak{K}(\mathbb{C} \odot A, B) \cong [\mathbb{C}, \mathfrak{K}(A, B)]$$

that are 2-natural in A , B , and \mathbb{C} . Since \mathcal{D} is a prederivator on \mathfrak{K} , it induces a functor $\mathfrak{K}(A, B) \rightarrow [\mathcal{D}^B, \mathcal{D}^A]$ that is 2-natural in A and in B , so we obtain a 2-natural functor $\mathfrak{K}(\mathbb{C} \odot A, B) \rightarrow [\mathbb{C}, [\mathcal{D}^B, \mathcal{D}^A]]$ by composition; but we have 2-natural isomorphisms

$$[\mathbb{C}, [\mathcal{D}^B, \mathcal{D}^A]] \cong [\mathbb{C} \times \mathcal{D}^B, \mathcal{D}^A] \cong [\mathcal{D}^B, [\mathbb{C}, \mathcal{D}^A]]$$

so, taking $B = \mathbb{C} \odot A$, we obtain the required functor $\mathcal{D}^{\mathbb{C} \odot A} \rightarrow [\mathbb{C}, \mathcal{D}^A]$. ■

Definition 4.1.8. A **strong semiderivator** on \mathfrak{K} is a semiderivator that satisfies the additional axiom below:

Der5. For any object A in \mathfrak{K} , the canonical functor $\mathcal{D}^{2 \odot A} \rightarrow [2, \mathcal{D}^A]$ is full and essentially surjective on objects (but not necessarily faithful).

REMARK 4.1.9. If \mathcal{D} is the prederivator represented by an object in \mathfrak{K} , then \mathcal{D} automatically satisfies axioms Der1 and Der5; and if \mathfrak{K} is a 2-full 2-subcategory of \mathfrak{Cat} with the same terminal object, then \mathcal{D} will also satisfy axiom Der2.

Lemma 4.1.10. *If \mathcal{C} is a uni-fractionable category, then the canonical comparison functor $\mathrm{Ho} [\min 2, \mathcal{C}]_{\mathrm{h}} \rightarrow [2, \mathrm{Ho} \mathcal{C}]$ is full and essentially surjective on objects.*

Proof. Let \mathcal{U} and \mathcal{V} be subcategories of $\mathrm{weq} \mathcal{C}$ such that \mathcal{C} admits a three-arrow calculus with respect to $(\mathcal{U}, \mathcal{V})$, and let $\bar{f} : X \rightarrow Y$ be any morphism in $\mathrm{Ho} \mathcal{C}$. By the fundamental theorem of three-arrow calculi (2.5.9), there exist $u : Y \rightarrow \hat{Y}$ in \mathcal{U} , $v : \tilde{X} \rightarrow X$ in \mathcal{V} , and $f : \tilde{X} \rightarrow \hat{Y}$ such that $\bar{f} = u^{-1} \circ f \circ v^{-1}$ in $\mathrm{Ho} \mathcal{C}$, i.e. such that the following diagram in $\mathrm{Ho} \mathcal{C}$ commutes:

$$\begin{array}{ccc} X & \xrightarrow{v^{-1}} & \tilde{X} \\ \bar{f} \downarrow & & \downarrow f \\ Y & \xrightarrow{u} & \hat{Y} \end{array}$$

It immediately follows that $\mathrm{Ho} [\min 2, \mathcal{C}]_{\mathrm{h}} \rightarrow [2, \mathcal{C}]$ is essentially surjective on objects.

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It remains to be shown that $\text{Ho}[\min \mathbb{2}, C]_{\text{h}} \rightarrow [2, C]$ is a full functor. Let $x : X_1 \rightarrow X_2$ and $y : Y_1 \rightarrow Y_2$ be morphisms in C , let $\bar{f}_1 : X_1 \rightarrow Y_1$ and $\bar{f}_2 : X_2 \rightarrow Y_2$ be morphisms in $\text{Ho} C$, and suppose we have $\bar{f}_2 \circ x = y \circ \bar{f}_1$; note this constitutes a morphism in $[2, C]$ between objects in the image of the functor $\text{Ho}[\min \mathbb{2}, C]_{\text{h}} \rightarrow [2, C]$. As before, we may choose $u_1 : Y_1 \rightarrow \hat{Y}_1$ and $u_2 : Y_2 \rightarrow \hat{Y}_2$ in \mathcal{U} , $v_1 : \tilde{X}_1 \rightarrow X_1$ and $v_2 : \tilde{X}_2 \rightarrow X_2$ in \mathcal{V} , and $f_1 : \tilde{X}_1 \rightarrow \hat{Y}_1$ and $f_2 : \tilde{X}_2 \rightarrow \hat{Y}_2$ in C such that the equations below hold in $\text{Ho} C$:

$$\bar{f}_1 = u_1^{-1} \circ f_1 \circ v_1^{-1} \qquad \bar{f}_2 = u_2^{-1} \circ f_2 \circ v_2^{-1}$$

Using axioms A2 and A3, there exist $u'_2 : Y_2 \rightarrow Z$ in \mathcal{U} , $v'_1 : W \rightarrow X_1$ in \mathcal{V} , and $z : \hat{Y}_1 \rightarrow Z$ and $w : W \rightarrow \tilde{X}_2$ making the following diagrams in C commute,

$$\begin{array}{ccccc} X_1 & \xleftarrow{v_1} & \tilde{X}_1 & \xrightarrow{f_1} & \hat{Y}_1 & \xleftarrow{u_1} & Y_1 \\ \parallel & & \parallel & & \downarrow z & & \downarrow y \\ X_1 & \xleftarrow{v_1} & \tilde{X}_1 & \xrightarrow{z \circ f_1} & Z & \xleftarrow{u'_2} & Y_2 \end{array}$$

$$\begin{array}{ccccc} X_1 & \xleftarrow{v'_1} & W & \xrightarrow{f_2 \circ w} & \hat{Y}_2 & \xleftarrow{u_2} & Y_2 \\ \downarrow x & & \downarrow w & & \parallel & & \parallel \\ X_2 & \xleftarrow{v_2} & \tilde{X}_2 & \xrightarrow{f_2} & \hat{Y}_2 & \xleftarrow{u_2} & Y_2 \end{array}$$

and since $\bar{f}_2 \circ x = y \circ \bar{f}_1$, the fundamental theorem says there exist a commutative diagram in C of the form below,

$$\begin{array}{ccccccc} X_1 & \xleftarrow{v'_1} & W & \xrightarrow{f_2 \circ w} & \hat{Y}_2 & \xleftarrow{u_2} & Y_2 \\ \parallel & & \uparrow v_3 & & \uparrow v_4 & & \parallel \\ X_1 & \xleftarrow{v_5} & \bullet & \xrightarrow{f_5} & \bullet & \xleftarrow{u_6} & Y_2 \\ \parallel & & \downarrow w_3 & & \downarrow w_4 & & \parallel \\ X_1 & \xleftarrow{v_6} & \bullet & \xrightarrow{f_6} & \bullet & \xleftarrow{u_6} & Y_2 \\ \parallel & & \uparrow u_3 & & \uparrow u_4 & & \parallel \\ X_1 & \xleftarrow{v_1} & \tilde{X}_1 & \xrightarrow{z \circ f_1} & Z & \xleftarrow{u'_2} & Y_2 \end{array}$$

where u_3, u_4, u_5, u_6 are in \mathcal{U} , v_3, v_4, v_5, v_6 are in \mathcal{V} , and w_3, w_4 are weak equivalences in C .

It is easy to verify that the following diagram in \mathcal{C} commutes,

$$\begin{array}{cccccccccccc}
 X_1 & \xrightarrow{v'_1} & W & \xlongequal{\quad} & W & \xleftarrow{v_3} & \bullet & \xrightarrow{w_3} & \bullet & \xleftarrow{u_3} & \tilde{X}_1 & \xrightarrow{f_1} & \hat{Y}_1 & \xleftarrow{u_1} & Y_1 \\
 x \downarrow & & w \downarrow & & f_2 \circ w \downarrow & & f_5 \downarrow & & f_6 \downarrow & & z \circ f_1 \downarrow & & z \downarrow & & y \downarrow \\
 X_2 & \xleftarrow{v_2} & \tilde{X}_2 & \xrightarrow{f_2} & \hat{Y}_2 & \xleftarrow{v_4} & \bullet & \xrightarrow{w_4} & \bullet & \xleftarrow{u_4} & Z & \xlongequal{\quad} & Z & \xleftarrow{u'_2} & Y_2
 \end{array}$$

and this is the required lift of (\bar{f}_1, \bar{f}_2) to $\text{Ho}[\text{min } 2, \mathcal{C}]_{\text{h}}$, because the diagram in \mathcal{C} shown below commutes:

$$\begin{array}{cccccccccccccccc}
 X_1 & \xlongequal{\quad} & X_1 & \xlongequal{\quad} & X_1 & \xlongequal{\quad} & X_1 & \xlongequal{\quad} & X_1 & \xleftarrow{v_1} & \tilde{X}_1 & \xrightarrow{f_1} & \hat{Y}_1 & \xleftarrow{u_1} & Y_1 \\
 \parallel & & v'_1 \uparrow & & v'_1 \uparrow & & v_5 \uparrow & & v_6 \uparrow & & \parallel & & \parallel & & \parallel \\
 X_1 & \xleftarrow{v'_1} & W & \xlongequal{\quad} & W & \xleftarrow{v_3} & \bullet & \xrightarrow{w_3} & \bullet & \xleftarrow{u_3} & \tilde{X}_1 & \xrightarrow{f_1} & \hat{Y}_1 & \xleftarrow{u_1} & Y_1 \\
 x \downarrow & & w \downarrow & & f_2 \circ w \downarrow & & f_5 \downarrow & & f_6 \downarrow & & z \circ f_1 \downarrow & & z \downarrow & & y \downarrow \\
 X_2 & \xleftarrow{v_2} & \tilde{X}_2 & \xrightarrow{f_2} & \hat{Y}_2 & \xleftarrow{v_4} & \bullet & \xrightarrow{w_4} & \bullet & \xleftarrow{u_4} & Z & \xlongequal{\quad} & Z & \xleftarrow{u'_2} & Y_2 \\
 \parallel & & \parallel & & \parallel & & u_5 \uparrow & & u_6 \uparrow & & u'_2 \uparrow & & u'_2 \uparrow & & \parallel \\
 X_2 & \xleftarrow{v_2} & \tilde{X}_2 & \xrightarrow{f_2} & \hat{Y}_2 & \xleftarrow{u_2} & Y_2 & \xlongequal{\quad} & Y_2 & \xlongequal{\quad} & Y_2 & \xlongequal{\quad} & Y_2 & \xlongequal{\quad} & Y_2
 \end{array}$$

We may therefore conclude that $\text{Ho}[\text{min } 2, \mathcal{C}]_{\text{h}} \rightarrow [2, \mathcal{C}]$ is indeed full. \blacksquare

Proposition 4.1.11. *Let $\mathcal{D}(\mathcal{M})$ be the prederivator of a \mathbf{U} -small relative category \mathcal{M} .*

- (i) \mathcal{D} satisfies axiom *Der1*.
- (ii) Moreover, if \mathcal{M} is a (necessarily saturated) homotopical category and each homotopical functor category $[\mathcal{A}, \mathcal{M}]_{\text{h}}$ admits a three-arrow calculus, then \mathcal{D} is a strong semiderivator.

Proof. (i). [Proposition A.4.16](#) implies \mathcal{D} sends finite coproducts in \mathbf{RelCat} to products in \mathbf{Cat}^+ , so axiom *Der1* is satisfied.

(ii). Suppose $f : X \rightarrow Y$ is a morphism in $\text{Ho}[\mathcal{A}, \mathcal{M}]_{\text{h}}$ such that all its components are isomorphisms in $\text{Ho } \mathcal{M}$. The fundamental theorem of three-arrow

calculus (2.5.9) says $f : X \rightarrow Y$ may be represented by a zigzag in $[\mathcal{A}, \mathcal{M}]_h$ of the form below,

$$X \xleftarrow{\psi} \tilde{X} \xrightarrow{\theta} \hat{Y} \xleftarrow{\varphi} Y$$

where ψ and φ are natural weak equivalences. Thus, if A is an object in \mathcal{A} , then the following zigzag represents an isomorphism in $\text{Ho } \mathcal{M}$:

$$XA \xleftarrow{\psi_A} \tilde{X}A \xrightarrow{\theta_A} \hat{Y}A \xleftarrow{\varphi_A} YA$$

However, [proposition 2.5.10](#) says \mathcal{M} is a saturated homotopical category, so θ_A must be a weak equivalence in \mathcal{M} as well; hence, $f : X \rightarrow Y$ is an isomorphism in $\text{Ho } [\mathcal{A}, \mathcal{M}]_h$. This shows that \mathcal{D} satisfies axiom Der2.

Finally, observe that $[\min 2 \times \mathcal{A}, \mathcal{M}]_h \cong [\min 2, [\mathcal{A}, \mathcal{M}]_h]_h$, and the hypothesis says $[\mathcal{A}, \mathcal{M}]_h$ admits a three-arrow calculus, so we apply [lemma 4.1.10](#) to deduce that axiom Der5 is satisfied. ■

Definition 4.1.12. Let \mathcal{D} be a prederivator on \mathfrak{K} , let $u : A \rightarrow B$ be a morphism in \mathfrak{K} , and let X be an object in \mathcal{D}^A .

- A **left \mathcal{D} -extension** of X along u is an initial object in the comma category $(X \downarrow u^*)$.
- A **right \mathcal{D} -extension** of X along u is a terminal object in the comma category $(u^* \downarrow X)$.
- We say \mathcal{D} has **left extensions** along u if the functor $u^* : \mathcal{D}^B \rightarrow \mathcal{D}^A$ has a left adjoint, which we denote by $u_! : \mathcal{D}^A \rightarrow \mathcal{D}^B$.
- We say \mathcal{D} has **right extensions** along u if the functor $u^* : \mathcal{D}^B \rightarrow \mathcal{D}^A$ has a right adjoint, which we denote by $u_* : \mathcal{D}^A \rightarrow \mathcal{D}^B$.

We may refer to left and right \mathcal{D} -extensions generically as **homotopy Kan extensions** in \mathcal{D} .

REMARK 4.1.13. It is straightforward to check that \mathcal{D} has left (resp. right) extensions along u if and only if, for every object X in \mathcal{D}^A , there exists a left (resp. right) \mathcal{D} -extension of X along u .

Example 4.1.14. If \mathfrak{K} is a 2-full 2-subcategory of \mathfrak{Cat} and \mathcal{D} is the prederivator represented by an object in \mathfrak{K} , then \mathcal{D} -extensions are exactly the same thing as Kan extensions in the usual sense.

As we saw in [theorem A.5.15](#), pointwise left (resp. right) Kan extensions can be computed as colimits (resp. limits) of certain diagrams whose shapes are comma categories. We shall shortly see that more is true.

Definition 4.1.15. Let \mathcal{D} be a prederivator on \mathfrak{K} and suppose we have a diagram in \mathfrak{K} of the following form:

$$\begin{array}{ccc} D & \xrightarrow{q} & B \\ p \downarrow & \theta \nearrow & \downarrow v \\ A & \xrightarrow{u} & C \end{array}$$

- We say the square is a **left \mathcal{D} -exact square** if \mathcal{D} has left extensions along $u : A \rightarrow C$ and $q : D \rightarrow B$ and the induced diagram shown below satisfies the left Beck–Chevalley condition:

$$\begin{array}{ccc} \mathcal{D}^C & \xrightarrow{v^*} & \mathcal{D}^B \\ u^* \downarrow & \theta^* \nearrow & \downarrow q^* \\ \mathcal{D}^A & \xrightarrow{p^*} & \mathcal{D}^D \end{array}$$

- We say the square is a **right \mathcal{D} -exact square** if \mathcal{D} has right extensions along $v : B \rightarrow C$ and $p : D \rightarrow A$ and the induced diagram shown below satisfies the right Beck–Chevalley condition:

$$\begin{array}{ccc} \mathcal{D}^C & \xrightarrow{u^*} & \mathcal{D}^A \\ v^* \downarrow & \theta^* \searrow & \downarrow p^* \\ \mathcal{D}^B & \xrightarrow{q^*} & \mathcal{D}^D \end{array}$$

A \mathcal{D} -**exact square** in \mathfrak{K} is a diagram in \mathfrak{K} that is both left \mathcal{D} -exact and right \mathcal{D} -exact.

Proposition 4.1.16. Let \mathcal{D} be a prederivator on \mathfrak{K} . Given the following diagram in \mathfrak{K} ,

$$\begin{array}{ccc} D & \xrightarrow{q} & B \\ p \downarrow & \theta \nearrow & \downarrow v \\ A & \xrightarrow{u} & C \end{array}$$

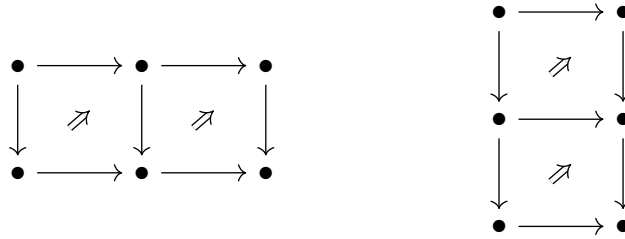
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if \mathcal{D} has left extensions along $u : A \rightarrow C$ and $q : D \rightarrow B$, and \mathcal{D} has right extensions along $v : B \rightarrow C$ and $p : D \rightarrow A$, then the following are equivalent:

- (i) The diagram is a \mathcal{D} -exact square.
- (ii) The diagram is a left \mathcal{D} -exact square.
- (iii) The diagram is a right \mathcal{D} -exact square.

Proof. Statement (i) is just the conjunction of statements (ii) and (iii), and when the required left and right adjoints exist, [proposition A.I.9](#) implies that statements (ii) and (iii) are equivalent. ■

Lemma 4.I.17 (Pasting exact squares). *Let \mathcal{D} be a prederivator, and consider pasting diagrams of the following forms in \mathfrak{K} :*



In either diagram, if both squares are left (resp. right) \mathcal{D} -exact squares, then the rectangle obtained by pasting the two squares is also a left (resp. right) \mathcal{D} -exact square.

Proof. Apply [lemma A.I.8](#). ■

Lemma 4.I.18. *Let \mathbf{Set} be the category of \mathbf{U} -sets. If \mathcal{D} is the prederivator of \mathbf{Set} restricted to the subdomain \mathfrak{Cat} , then every comma square in \mathfrak{Cat} is a right \mathcal{D} -exact square.*

Proof. Suppose we have the following comma square in \mathfrak{Cat} :

$$\begin{array}{ccc}
 (u \downarrow v) & \xrightarrow{q} & \mathbb{B} \\
 p \downarrow & \theta \nearrow & \downarrow v \\
 \mathbb{A} & \xrightarrow{u} & \mathbb{C}
 \end{array}$$

Let $Y : \mathbb{B} \rightarrow \mathbf{Set}$ be a functor and let (Z, ε) be a right Kan extension of Y along v , i.e. a terminal object in the comma category $(v^* \downarrow Y)$. In view of [lemma A.1.7](#), to deduce the claim, it is enough to show that $(u^*(Z), q^*(\varepsilon) \bullet \theta_Z^*)$ is a terminal object in the comma category $(p^* \downarrow q^*(Y))$, i.e. a right Kan extension of Yq along p ; but this was done in [lemma A.5.8](#). \blacksquare

Proposition 4.1.19. *Let \mathcal{M} be a locally \mathbf{U} -small category, and let \mathcal{D} be the prederivator of \mathcal{M} restricted to \mathfrak{Cat} .*

- *If \mathcal{M} has colimits for all \mathbf{U} -small diagrams, then every comma square in \mathfrak{Cat} is a left \mathcal{D} -exact square.*
- *If \mathcal{M} has limits for all \mathbf{U} -small diagrams, then every comma square in \mathfrak{Cat} is a right \mathcal{D} -exact square.*

Proof. The two claims are formally dual; we will prove the first version.

Consider a comma square in \mathfrak{Cat} :

$$\begin{array}{ccc} \mathbb{D} & \xrightarrow{q} & \mathbb{B} \\ p \downarrow & \theta \nearrow & \downarrow v \\ \mathbb{A} & \xrightarrow{u} & \mathbb{C} \end{array}$$

If \mathcal{M} has colimits for all \mathbf{U} -small diagrams, then [theorem A.5.15](#) implies that, for any functor $X : \mathbb{A} \rightarrow \mathcal{M}$, the left Kan extension of X along u exists and is pointwise, and same is true for the left Kan extension of $p^*(X)$ along q . Thus, for any object M in \mathcal{M} , if $h_M : \mathcal{M}^{\text{op}} \rightarrow \mathbf{Set}^+$ is the representable functor $\mathcal{M}(-, M)$, we may use [lemma A.1.7](#) to deduce that the following (commutative!) diagrams satisfy the *right Beck–Chevalley condition*:

$$\begin{array}{ccc} [\mathbb{C}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{[\mathbb{C}^{\text{op}}, h_M]} & [\mathbb{C}^{\text{op}}, \mathbf{Set}^+] \\ (u^{\text{op}})^* \downarrow & & \downarrow (u^{\text{op}})^* \\ [\mathbb{A}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{[\mathbb{A}^{\text{op}}, h_M]} & [\mathbb{A}^{\text{op}}, \mathbf{Set}^+] \end{array} \quad \begin{array}{ccc} [\mathbb{B}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{[\mathbb{A}^{\text{op}}, h_M]} & [\mathbb{B}^{\text{op}}, \mathbf{Set}^+] \\ (q^{\text{op}})^* \downarrow & & \downarrow (q^{\text{op}})^* \\ [\mathbb{D}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{[\mathbb{B}^{\text{op}}, h_M]} & [\mathbb{D}^{\text{op}}, \mathbf{Set}^+] \end{array}$$

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On the other hand, lemma 4.1.18 says the diagram below satisfies the right Beck–Chevalley condition,

$$\begin{array}{ccc}
 [\mathbb{C}^{\text{op}}, \mathbf{Set}^+] & \xrightarrow{(v^{\text{op}})^*} & [\mathbb{B}^{\text{op}}, \mathbf{Set}^+] \\
 (u^{\text{op}})^* \downarrow & \swarrow_{(\theta^{\text{op}})^*} & \downarrow (q^{\text{op}})^* \\
 [\mathbb{A}^{\text{op}}, \mathbf{Set}^+] & \xrightarrow{(p^{\text{op}})^*} & [\mathbb{D}^{\text{op}}, \mathbf{Set}^+]
 \end{array}$$

and the family $\{h_M : \mathcal{M}^{\text{op}} \rightarrow \mathbf{Set}^+ \mid M \in \text{ob } \mathcal{M}\}$ is jointly conservative, so we deduce that the right Beck–Chevalley condition for the following diagram is satisfied,

$$\begin{array}{ccc}
 [\mathbb{C}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{(v^{\text{op}})^*} & [\mathbb{B}^{\text{op}}, \mathcal{M}^{\text{op}}] \\
 (u^{\text{op}})^* \downarrow & \swarrow_{(\theta^{\text{op}})^*} & \downarrow (q^{\text{op}})^* \\
 [\mathbb{A}^{\text{op}}, \mathcal{M}^{\text{op}}] & \xrightarrow{(p^{\text{op}})^*} & [\mathbb{D}^{\text{op}}, \mathcal{M}^{\text{op}}]
 \end{array}$$

and therefore this diagram satisfies the *left* Beck–Chevalley condition:

$$\begin{array}{ccc}
 [\mathbb{C}, \mathcal{M}] & \xrightarrow{v^*} & [\mathbb{B}, \mathcal{M}] \\
 u^* \downarrow & \nearrow_{\theta^*} & \downarrow q^* \\
 [\mathbb{A}, \mathcal{M}] & \xrightarrow{p^*} & [\mathbb{D}, \mathcal{M}]
 \end{array}$$

We then conclude that every comma square in \mathfrak{Cat} is a left \mathcal{D} -exact square. \blacksquare

Definition 4.1.20. A \mathfrak{K} -cocomplete semiderivator is a semiderivator \mathcal{D} on \mathfrak{K} satisfying these additional axioms:

Der3L. \mathcal{D} has left extensions along every morphism $u : A \rightarrow B$ in \mathfrak{K} .

Der4L. Every comma square in \mathfrak{K} of the form below is a left \mathcal{D} -exact square:

$$\begin{array}{ccc}
 (u \downarrow c) & \longrightarrow & 1 \\
 \downarrow & \nearrow & \downarrow c \\
 A & \xrightarrow{u} & C
 \end{array}$$

Dually, a \mathfrak{K} -complete U-semiderivator is one satisfying these axioms:

Der3R. \mathcal{D} has right extensions along every morphism $u : A \rightarrow B$ in \mathfrak{K} .

Der4R. Every comma square in \mathfrak{K} of the form below is a right \mathcal{D} -exact square:

$$\begin{array}{ccc} (c \downarrow v) & \longrightarrow & B \\ \downarrow & \nearrow & \downarrow v \\ 1 & \xrightarrow{c} & C \end{array}$$

Theorem 4.1.21. Let \mathbf{U}^+ be a universe with $\mathbf{U} \subseteq \mathbf{U}^+$, let \mathcal{M} be a \mathbf{U}^+ -small category, and let \mathcal{D} be the prederivator of \mathcal{M} restricted to \mathfrak{Cat} .

- (i) \mathcal{D} is a strong semiderivator.
- (ii) \mathcal{D} is \mathfrak{Cat} -cocomplete (resp. \mathfrak{Cat} -complete) if and only if \mathcal{M} is \mathbf{U} -complete (resp. \mathbf{U} -complete).

Proof. (i). This can be shown using the same arguments as [remark 4.1.9](#).

(ii). This is the content of [proposition 4.1.19](#). ■

Finally, we come to the definition of the subject of this chapter:

Definition 4.1.22. A **derivator** on \mathfrak{K} is a semiderivator that is \mathfrak{K} -cocomplete and \mathfrak{K} -complete, and a **strong derivator** is one that satisfies axiom Der5.

REMARK 4.1.23. The definition of ‘subdomain’ ensures that the restriction of any derivator (resp. semiderivator, complete semiderivator, cocomplete semiderivator) on \mathfrak{K} to any subdomain of \mathfrak{K} is again a derivator (resp. semiderivator, complete semiderivator, cocomplete semiderivator).

Proposition 4.1.24. Let \mathcal{D} be a prederivator on \mathfrak{K} , and let $u \dashv v : B \rightarrow A$ be an adjunction in \mathfrak{K} , with unit $\eta : \text{id}_A \Rightarrow v \circ u$ and counit $\varepsilon : u \circ v \Rightarrow \text{id}_B$.

- (i) We have an adjunction $v^* \dashv u^* : \mathcal{D}^B \rightarrow \mathcal{D}^A$, with unit $\eta^* : \text{id}_{\mathcal{D}^A} \Rightarrow u^* \circ v^*$ and counit $\varepsilon^* : v^* \circ u^* \Rightarrow \text{id}_{\mathcal{D}^B}$; in particular, \mathcal{D} has left extensions along $u : A \rightarrow B$ and right extensions along $v : B \rightarrow A$.
- (ii) Consider the following commutative diagrams in \mathfrak{K} :

$$\begin{array}{ccc} A & \xrightarrow{u} & B \\ p \downarrow & & \downarrow q \\ 1 & \xrightarrow{\text{id}} & 1 \end{array} \qquad \begin{array}{ccc} B & \xrightarrow{q} & 1 \\ v \downarrow & & \downarrow \text{id} \\ A & \xrightarrow{p} & 1 \end{array}$$

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The diagram on the left is a left \mathcal{D} -exact square, and the diagram on the right is a right \mathcal{D} -exact square.

- (iii) Moreover, if \mathcal{D} has left extensions along $p : A \rightarrow 1$ and $q : A \rightarrow 1$, then the diagram on the left is a left \mathcal{D} -exact square; and if \mathcal{D} has right extensions along $p : A \rightarrow 1$ and $q : A \rightarrow 1$, then the diagram on the right is a right \mathcal{D} -exact square.

Proof. (i). Since \mathcal{D} is a 2-functor, it preserves the triangle identities; thus $v^* \dashv u^*$ is indeed an adjunction. (The left and right adjoints are exchanged because \mathcal{D} is contravariant.)

(ii). The two halves of the claim are formally dual; we will prove the first version. By lemma A.I.7, it is enough to show that, for any object M in \mathcal{D}^1 , the object $(p^*(M), \text{id}_{q^*(M)})$ is an initial object in the comma category $(q^*(M) \downarrow u^*)$.

Noting that $v^*q^*(M) = p^*(M)$ and $u^*p^*(M) = q^*(M)$, we see that, for any object (X, α) in $(q^*(M) \downarrow u^*)$, the diagram in \mathcal{D}^B shown below commutes,

$$\begin{array}{ccc} u^*v^*q^*(M) & \xrightarrow{u^*v^*(\alpha)} & u^*v^*u^*(X) \\ \parallel & & \downarrow u^*(\epsilon_X^*) \\ q^*(M) & \xrightarrow{\alpha} & u^*(X) \end{array}$$

and a similar calculation shows that any $\varphi : (p^*(M), \text{id}_{q^*(M)}) \rightarrow (X, \alpha)$ in $(q^*(M) \downarrow u^*)$ must be equal to $\epsilon_X^* \circ v^*(\alpha)$. This shows that $(p^*(M), \text{id}_{q^*(M)})$ is an initial object in $(q^*(M) \downarrow u^*)$.

- (iii). This is a special case of proposition 4.I.16. ■

Theorem 4.I.25. *Let \mathcal{D} be a semiderivator on \mathfrak{K} that satisfies axioms Der3L and Der3R, and let 1 be a terminal object in \mathfrak{K} . The following are equivalent:*

- (i) \mathcal{D} is a derivator.
- (ii) \mathcal{D} satisfies axiom Der4L.
- (iii) Every comma square in \mathfrak{K} is left \mathcal{D} -exact.
- (iv) \mathcal{D} satisfies axiom Der4R.
- (v) Every comma square in \mathfrak{K} is right \mathcal{D} -exact.

Proof. Obviously, statement (i) implies statements (ii)–(v), and the conjunction of statements (iii) and (v) implies statement (i). We are assuming that \mathcal{D} has left and right extensions along all morphisms in \mathfrak{K} , so the equivalence of statements (iii) and (v) is just [proposition 4.1.16](#). It remains to be shown that (ii) \Rightarrow (iii) and (iv) \Rightarrow (v), but the two implications are formally dual, so it is enough to prove just one; we prove the former.

Consider a general comma square in \mathfrak{K} :

$$\begin{array}{ccc} (u \downarrow v) & \xrightarrow{q} & B \\ p \downarrow & \theta \nearrow & \downarrow v \\ A & \xrightarrow{u} & C \end{array}$$

Let $b : 1 \rightarrow B$ be a morphism in \mathfrak{K} , and let $c = v \circ b$, and consider the following pasting diagrams,

$$\begin{array}{ccc} (q \downarrow b) & \xrightarrow{t} & 1 \\ s \downarrow & \tau \nearrow & \downarrow b \\ (u \downarrow v) & \xrightarrow{q} & B \\ p \downarrow & \theta \nearrow & \downarrow v \\ A & \xrightarrow{u} & C \end{array} \qquad \begin{array}{ccc} (u \downarrow c) & \xrightarrow{r} & 1 \\ j \downarrow & // & \downarrow b \\ (u \downarrow v) & \xrightarrow{q} & B \\ p \downarrow & \theta \nearrow & \downarrow v \\ A & \xrightarrow{u} & C \end{array}$$

where the upper square of the diagram on the left is a comma square, and the upper square of the diagram on the right is a 2-pullback square; note that the pasting lemma for comma squares implies that the outer rectangle of the diagram on the right is also a comma square.

Let $\pi = p \circ j$ and let $\lambda = \theta \circ \text{id}_j$. By the universal property of comma objects, there is a unique morphism $f : (q \downarrow b) \rightarrow (u \downarrow c)$ such that $\pi \circ f = p \circ s$, $r \circ f = t$, and $\lambda \circ \text{id}_f = (\text{id}_v \circ \tau) \cdot (\theta \circ \text{id}_s)$; and similarly there is a unique morphism $g : (u \downarrow c) \rightarrow (q \downarrow b)$ such that $s \circ g = j$, $t \circ g = r$, and $\tau \circ \text{id}_g = \text{id}_{q \circ j} = \text{id}_{b \circ r}$. Then,

$$\pi \circ (f \circ g) = p \circ s \circ g = p \circ j = \pi \qquad r \circ (f \circ g) = r$$

so $f \circ g = \text{id}_{(u \downarrow c)}$; and since $p \circ s = p \circ s \circ g \circ f$, we may think of $\text{id}_{p \circ s}$ as a 2-cell $\beta : p \circ s \Rightarrow p \circ s \circ g \circ f$, whereas $b \circ t = q \circ s \circ g \circ f$, so $\tau : q \circ s \Rightarrow b \circ t$ is also a

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2-cell $\gamma : q \circ s \Rightarrow q \circ s \circ g \circ f$, but then

$$(\theta \circ \text{id}_{s \circ g \circ f}) \bullet (\text{id}_u \circ \beta) = \theta \circ \text{id}_j \circ \text{id}_f = \lambda \circ \text{id}_f = (\text{id}_v \circ \gamma) \bullet (\theta \circ \text{id}_s)$$

so by the 2-universal property of $(u \downarrow v)$, there is a unique 2-cell $\alpha : s \Rightarrow s \circ g \circ f$ such that $\text{id}_p \circ \alpha = \beta$ and $\text{id}_q \circ \alpha = \gamma$; and furthermore,

TODO: Justify this more carefully...

$$(\tau \circ \text{id}_{g \circ f}) \bullet (\text{id}_q \circ \alpha) = (\text{id}_b \circ \text{id}_{t \circ g \circ f}) \bullet \tau$$

therefore there is a unique 2-cell $\eta : \text{id}_{(q \downarrow b)} \Rightarrow g \circ f$ such that $\text{id}_s \circ \eta = \alpha$ and $\text{id}_t \circ \eta = \text{id}_{t \circ g \circ f}$.

We will now show that we have an adjunction $f \dashv g : (u \downarrow c) \rightarrow (q \downarrow b)$ in \mathfrak{K} ; since $f \circ g = \text{id}_{(u \downarrow c)}$, it is enough to check that $\text{id}_f \circ \eta = \text{id}_f$ and $\eta \circ \text{id}_g = \text{id}_g$. By construction, $\text{id}_\pi \circ (\text{id}_f \circ \eta) = \text{id}_p \circ \text{id}_s \circ \eta = \text{id}_{p \circ s}$, and $\text{id}_r \circ (\text{id}_f \circ \eta) = \text{id}_t$, so indeed $\text{id}_f \circ \eta = \text{id}_f$; and $\text{id}_s \circ (\eta \circ \text{id}_g) = \text{id}_s$ and $\text{id}_t \circ (\eta \circ \text{id}_g) = \text{id}_t$, so $\eta \circ \text{id}_g = \text{id}_g$ as well. Thus, by [proposition 4.1.24](#), the commutative diagram in \mathfrak{K} shown below on the left is a left \mathcal{D} -exact square,

$$\begin{array}{ccc} (u \downarrow c) & \xrightarrow{r} & 1 \\ g \downarrow & & \downarrow \text{id} \\ (q \downarrow b) & \longrightarrow & 1 \end{array} \qquad \begin{array}{ccc} (q \downarrow b) & \xrightarrow{t} & 1 \\ s \downarrow & \nearrow \tau & \downarrow b \\ (u \downarrow v) & \xrightarrow{q} & B \end{array}$$

and the diagram on the right is a left \mathcal{D} -exact square by hypothesis, so by the pasting lemma ([4.1.17](#)), the following commutative diagram is also a left \mathcal{D} -exact square:

$$\begin{array}{ccc} (u \downarrow c) & \longrightarrow & 1 \\ j \downarrow & & \downarrow b \\ (u \downarrow v) & \xrightarrow{q} & B \end{array}$$

The hypothesis also implies that this diagram satisfies the left Beck–Chevalley condition,

$$\begin{array}{ccc} \mathcal{D}^C & \xrightarrow{c^*} & \mathcal{D}^1 \\ u^* \downarrow & \nearrow \lambda^* & \downarrow r^* \\ \mathcal{D}^A & \xrightarrow{\pi^*} & \mathcal{D}^{(u \downarrow c)} \end{array}$$

but the pasting lemma (A.1.8) says that the left Beck–Chevalley transformation $r_! \pi^* \Rightarrow c^* u_!$ is obtained by pasting together the left Beck–Chevalley transformations of the squares in the diagram below,

$$\begin{array}{ccccc}
 \mathcal{D}^C & \xrightarrow{v^*} & \mathcal{D}^B & \xrightarrow{b^*} & \mathcal{D}^1 \\
 u^* \downarrow & \swarrow \theta^* & \downarrow q^* & \parallel & \downarrow r^* \\
 \mathcal{D}^A & \xrightarrow{p^*} & \mathcal{D}^{(u \downarrow v)} & \xrightarrow{j^*} & \mathcal{D}^{(u \downarrow e)}
 \end{array}$$

and so, allowing $b : 1 \rightarrow B$ to vary, we deduce that every component of the left Beck–Chevalley transformation $v^* u_! \Rightarrow q_! p^*$ is an isomorphism in \mathcal{D}^1 . We may then apply axiom Der2 to conclude that the comma square we started with is a left \mathcal{D} -exact square. ■

GENERALITIES

A.I Adjoints and mates

We begin by recalling a standard definition:

Definition A.I.1. An **adjunction of categories** consists of the following data:

- A functor $F : C \rightarrow D$, called the **left adjoint**.
- A functor $G : D \rightarrow C$, called the **right adjoint**.
- A natural transformation $\eta : \text{id}_C \Rightarrow GF$, called the **unit**.
- A natural transformation $\varepsilon : FG \Rightarrow \text{id}_D$, called the **counit**.

These are moreover required to satisfy the **triangle identities**:

$$\varepsilon F \bullet F \eta = \text{id}_F \qquad G \varepsilon \bullet \eta G = \text{id}_G$$

If such data exist, we write

$$F \dashv G : D \rightarrow C$$

and say that F is a left adjoint of G , and G is a right adjoint of F .

Proposition A.I.2. *Let $F \dashv G : D \rightarrow C$ and $F' \dashv G' : D' \rightarrow C'$ be adjunctions, let $\eta : \text{id}_C \Rightarrow GF$ and $\eta' : \text{id}_{C'} \Rightarrow G'F'$ be the respective units, and let $\varepsilon : FG \Rightarrow \text{id}_D$ and $\varepsilon' : F'G' \Rightarrow \text{id}_{D'}$ be the respective counits. Let $H : C \rightarrow C'$ and*

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$K : \mathcal{D} \rightarrow \mathcal{D}'$ be functors, and let φ and ψ be natural transformations as in the diagrams below:

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \\ F \downarrow & \swarrow \varphi & \downarrow F' \\ \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \end{array} \qquad \begin{array}{ccc} \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \\ G \downarrow & \nearrow \psi & \downarrow G' \\ \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \end{array}$$

Then, the following are equivalent:

- (i) $\varepsilon' KF \cdot F' \psi F \cdot F' H \eta = \varphi$.
- (ii) $\psi F \cdot H \eta = G' \varphi \cdot \eta' H$.
- (iii) $\psi = G' K \varepsilon \cdot G' \varphi G \cdot \eta' HG$.
- (iv) $\varepsilon' K \cdot F' \psi = K \varepsilon \cdot \varphi G$.

Proof.

$$\begin{aligned} \text{(i)} \Rightarrow \text{(ii).} \quad G' \varphi \cdot \eta' H &= G' \varepsilon' KF \cdot G' F' \psi F \cdot G' F' H \eta \cdot \eta' H \\ &= G' \varepsilon' KF \cdot \eta' G' KF \cdot \psi F \cdot H \eta \\ &= \psi F \cdot H \eta \end{aligned}$$

$$\begin{aligned} \text{(ii)} \Rightarrow \text{(iii).} \quad G' K \varepsilon \cdot G' \varphi G \cdot \eta' HG &= G' K \varepsilon \cdot \psi FG \cdot H \eta G \\ &= \psi \cdot HG \varepsilon \cdot H \eta G \\ &= \psi \end{aligned}$$

$$\begin{aligned} \text{(iii)} \Rightarrow \text{(iv).} \quad \varepsilon' K \cdot F' \psi &= \varepsilon' K \cdot F' G' K \varepsilon \cdot F' G' \varphi G \cdot F' \eta' HG \\ &= K \varepsilon \cdot \varphi G \cdot \varepsilon' HG \cdot F' \eta' HG \\ &= K \varepsilon \cdot \varphi G \end{aligned}$$

$$\begin{aligned} \text{(iv)} \Rightarrow \text{(i).} \quad \varepsilon' KF \cdot F' \psi F \cdot F' H \eta &= K \varepsilon F \cdot \varphi GF \cdot F' H \eta \\ &= K \varepsilon F \cdot KF \eta \cdot \varphi \\ &= \varphi \end{aligned} \quad \blacksquare$$

Definition A.I.3. A **conjugate pair of natural transformations** is a pair (φ, ψ) satisfying the equivalent conditions of the above proposition. Given such, we say φ is the **left mate** of ψ , and ψ is the **right mate** of φ .

Definition A.I.4. Let $F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$ and $F' \dashv G' : \mathcal{D}' \rightarrow \mathcal{C}'$ be adjunctions, let $H : \mathcal{C} \rightarrow \mathcal{C}'$ and $K : \mathcal{D} \rightarrow \mathcal{D}'$ be functors, and let φ and ψ be a conjugate pair of natural transformations as in the diagrams below:

$$\begin{array}{ccc}
 \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \\
 F \downarrow & \swarrow \varphi & \downarrow F' \\
 \mathcal{D} & \xrightarrow{K} & \mathcal{D}'
 \end{array}
 \qquad
 \begin{array}{ccc}
 \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \\
 G \downarrow & \nearrow \psi & \downarrow G' \\
 \mathcal{C} & \xrightarrow{H} & \mathcal{C}'
 \end{array}$$

We say the diagram on the right satisfies the **left Beck–Chevalley condition** if the left mate φ is a natural isomorphism, and we say the diagram on the left satisfies the **right Beck–Chevalley condition** if the right mate ψ is a natural isomorphism. More generally, the **local left Beck–Chevalley condition** is satisfied at an object C in \mathcal{C} if the component $\varphi_C : F'HC \rightarrow KFC$ is an isomorphism, and the **local right Beck–Chevalley condition** is satisfied at an object D in \mathcal{D} if the component $\psi_D : HGD \rightarrow G'KD$ is an isomorphism.

REMARK A.I.5. Unfortunately, the Beck–Chevalley conditions are not vacuous. For example, consider the following (strictly!) commutative diagram of forgetful functors:

$$\begin{array}{ccc}
 \mathbf{CRing} & \longrightarrow & \mathbf{Ab} \\
 \downarrow & & \downarrow \\
 \mathbf{Set} & \xrightarrow{\text{id}} & \mathbf{Set}
 \end{array}$$

The left mate of the trivial natural transformation in the above diagram is the group homomorphism $\mathbb{Z}X \rightarrow \mathbb{Z}[X]$ that sends a generator in $\mathbb{Z}X$ to the corresponding generator in $\mathbb{Z}[X]$; clearly, this is never an isomorphism. However, this is unsurprising: we do not expect the free abelian group generated by X to be naturally isomorphic to the additive group of free commutative ring generated by X .

Example A.1.6. Let \mathcal{C} be a category with pullbacks, and suppose the following diagram is a pullback square in \mathcal{C} :

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

Let $\Sigma_f : \mathcal{C}_{/X} \rightarrow \mathcal{C}_{/Y}$ etc. be the functor that sends an object $p : E \rightarrow X$ in $\mathcal{C}_{/X}$ to the object $f \circ p : E \rightarrow Y$ in $\mathcal{C}_{/Y}$, and consider the induced (strictly!) commutative diagram of functors:

$$\begin{array}{ccc} \mathcal{C}_{/Z} & \xrightarrow{\Sigma_z} & \mathcal{C}_{/X} \\ \Sigma_g \downarrow & & \downarrow \Sigma_f \\ \mathcal{C}_{/W} & \xrightarrow{\Sigma_w} & \mathcal{C}_{/Y} \end{array}$$

Since \mathcal{C} has pullbacks, Σ_g and Σ_f have right adjoints,^[1] and the pullback pasting lemma then implies that the above square satisfies the right Beck–Chevalley condition.

Lemma A.1.7. *Given a diagram of functors and natural transformations of the form below,*

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \\ G \downarrow & \psi \nearrow & \downarrow G' \\ \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \end{array}$$

where $\psi : HG \Rightarrow G'K$ is a natural isomorphism, $F \dashv G$, and $F' \dashv G'$, for each object C in \mathcal{C} , the following are equivalent:

- (i) *The diagram satisfies the local left Beck–Chevalley condition at \mathcal{C} .*
- (ii) *The functor $(\mathcal{C} \downarrow G) \rightarrow (\mathcal{H}\mathcal{C} \downarrow G')$ sending an object (D, f) in the comma category $(\mathcal{C} \downarrow G)$ to the object $(KD, \psi_D \circ Hf)$ in $(\mathcal{H}\mathcal{C} \downarrow G')$ preserves initial objects.*

[1] See lemma A.2.17.

Proof. We know (FC, η_C) is an initial object of $(C \downarrow G)$ and $(F'HC, \eta'_{HC})$ is an initial object of $(HC \downarrow G')$, so there is a unique morphism $\varphi_C : F'HC \rightarrow KFC$ such that $G'\varphi_C \circ \eta'_{HC} = \psi_{FC} \circ H\eta_C$. However, we observe that

$$\begin{aligned}\varphi_C &= \varphi_C \circ \varepsilon'_{F'HC} \circ F'\eta'_{HC} \\ &= \varepsilon'_{KFC} \circ F'G'\varphi_C \circ F'\eta'_{HC} \\ &= \varepsilon'_{KFC} \circ F'\psi_{FC} \circ F'H\eta_C\end{aligned}$$

so φ_C is precisely the component at C of the left mate of ψ . ■

Lemma A.I.8 (Pasting conjugate pairs).

- (i) Let $F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$, $F' \dashv G' : \mathcal{D}' \rightarrow \mathcal{C}'$, and $F'' \dashv G'' : \mathcal{D}'' \rightarrow \mathcal{C}''$ be adjunctions, let $H : \mathcal{C} \rightarrow \mathcal{C}'$, $H' : \mathcal{C}' \rightarrow \mathcal{C}''$, $K : \mathcal{D} \rightarrow \mathcal{D}'$, and $K' : \mathcal{D}' \rightarrow \mathcal{D}''$ be functors, and let $\varphi, \varphi', \psi, \psi'$ be natural transformations as in the following pasting diagrams:

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{H} & \mathcal{C}' & \xrightarrow{H'} & \mathcal{C}'' \\ F \downarrow & \swarrow \varphi & \downarrow F' & \swarrow \varphi' & \downarrow F'' \\ \mathcal{D} & \xrightarrow{K} & \mathcal{D}' & \xrightarrow{K'} & \mathcal{D}'' \end{array} \quad \begin{array}{ccc} \mathcal{D} & \xrightarrow{K} & \mathcal{D}' & \xrightarrow{K'} & \mathcal{D}'' \\ G \downarrow & \swarrow \psi & \downarrow G' & \swarrow \psi' & \downarrow G'' \\ \mathcal{C} & \xrightarrow{H} & \mathcal{C}' & \xrightarrow{H'} & \mathcal{C}'' \end{array}$$

Let $\bar{\varphi} = K'\varphi \circ \varphi'H$ and $\bar{\psi} = \psi'K \circ H'\psi$. If (φ, ψ) and (φ', ψ') are conjugate pairs, then $(\bar{\varphi}, \bar{\psi})$ is also a conjugate pair.

- (ii) Let $F_1 \dashv G_1 : \mathcal{D} \rightarrow \mathcal{C}$, $F_2 \dashv G_2 : \mathcal{E} \rightarrow \mathcal{D}$, $F'_1 \dashv G'_1 : \mathcal{D}' \rightarrow \mathcal{C}'$, and $F'_2 \dashv G'_2 : \mathcal{E}' \rightarrow \mathcal{D}'$ be adjunctions, let $H : \mathcal{C} \rightarrow \mathcal{C}'$, $K : \mathcal{D} \rightarrow \mathcal{D}'$, and $L : \mathcal{E} \rightarrow \mathcal{E}'$ be functors, and let $\varphi_1, \varphi_2, \psi_1, \psi_2$ be natural transformations as in the following pasting diagrams:

$$\begin{array}{ccc} \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \\ F_1 \downarrow & \swarrow \varphi_1 & \downarrow F'_1 \\ \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \\ F_2 \downarrow & \swarrow \varphi_2 & \downarrow F'_2 \\ \mathcal{E} & \xrightarrow{L} & \mathcal{E}' \end{array} \quad \begin{array}{ccc} \mathcal{E} & \xrightarrow{L} & \mathcal{E}' \\ G_2 \downarrow & \swarrow \psi_2 & \downarrow G'_2 \\ \mathcal{D} & \xrightarrow{K} & \mathcal{D}' \\ G_1 \downarrow & \swarrow \psi_1 & \downarrow G'_1 \\ \mathcal{C} & \xrightarrow{H} & \mathcal{C}' \end{array}$$

Let $\varphi = \varphi_2 F_1 \circ F'_2 \varphi_1$ and $\psi = G'_1 \psi_2 \circ \psi_1 G_2$. If (φ_1, ψ_1) and (φ_2, ψ_2) are conjugate pairs, then (φ, ψ) is also a conjugate pair.

Proof. These are straightforward exercises in using the triangle identities. \blacklozenge

Proposition A.1.9. *Let $u_! \dashv u^* : C \rightarrow \mathcal{A}$, $q_! \dashv q^* : B \rightarrow \mathcal{D}$, $v^* \dashv v_* : B \rightarrow C$, and $p^* \dashv p_* : \mathcal{D} \rightarrow \mathcal{A}$ be adjunctions, and let $\theta : u^* p^* \Rightarrow v^* q^*$ be a natural transformation.*

$$\begin{array}{ccc} C & \xrightarrow{v^*} & B \\ u^* \downarrow & \theta \nearrow & \downarrow q^* \\ \mathcal{A} & \xrightarrow{p^*} & \mathcal{D} \end{array} \qquad \begin{array}{ccc} C & \xrightarrow{u^*} & \mathcal{A} \\ v^* \downarrow & \theta \swarrow & \downarrow p^* \\ B & \xrightarrow{q^*} & \mathcal{D} \end{array}$$

The following are equivalent:

- (i) The diagram on the left satisfies the left Beck–Chevalley condition.
- (ii) The diagram on the right satisfies the right Beck–Chevalley condition.

Proof. Let $\varphi : q_! p^* \Rightarrow v^* u_!$ be the left mate of θ , and let $\psi : u^* v_* \Rightarrow p_* q^*$ be the right mate of θ . Then, by [proposition A.1.2](#),

$$\begin{aligned} \theta u_! \cdot p^* \eta^u &= q^* \varphi \cdot \eta^q p^* & \varepsilon^q v^* \cdot q_! \theta &= v^* \varepsilon^u \cdot \varphi u^* \\ \psi v^* \cdot u^* \eta^v &= p_* \theta \cdot \eta^p u^* & \varepsilon^p q^* \cdot p^* \psi &= q^* \varepsilon^v \cdot \theta v_* \end{aligned}$$

where the η denote the various adjunction units and the ε denote the various adjunction counits, thus:

$$\begin{aligned} \psi v^* u_! \cdot (u^* \eta^v u_! \cdot \eta^u) &= p_* \theta u_! \cdot \eta^p u^* u_! \cdot \eta^u \\ p_* \theta u_! \cdot p_* p^* \eta^u \cdot \eta^p &= p_* q^* \varphi \cdot (p_* \eta^q p^* \cdot \eta^p) \\ (\varepsilon^q \cdot q_! \varepsilon^p q^*) \cdot q_! p^* \psi &= \varepsilon^q \cdot q_! q^* \varepsilon^v \cdot q_! \theta v_* \\ \varepsilon^v \cdot \varepsilon^q v^* v_* \cdot q_! \theta v_* &= (\varepsilon^v \cdot v^* \varepsilon^u v_*) \cdot \varphi u^* v_* \end{aligned}$$

Thus, (φ, ψ) is a conjugate pair of natural transformations between the adjunctions $v^* u_! \dashv u^* v_*$ and $q_! p^* \dashv p_* q^*$. It follows (by [lemma A.1.8](#)) that φ is a natural isomorphism if and only if ψ is a natural isomorphism. \blacksquare

A.2 Cartesian closed categories

Definition A.2.1. Let C be a category with binary products, and let Y and Z be objects in C . An **exponential object** for Y and Z is an object $[Y, Z]_C$ in C and a morphism $\text{ev}_{Y,Z} : [Y, Z]_C \times Y \rightarrow Z$ with the following universal property:

- For all morphisms $f : X \times Y \rightarrow Z$ in C , there exists a unique morphism $\bar{f} : X \rightarrow [Y, Z]_C$ such that $\text{ev}_{Y,Z} \circ (\bar{f} \times \text{id}_Y) = f$.

An **exponentiable object** in C is an object Y such that, for all objects Z in C , the exponential object $[Y, Z]_C$ exists. We may write $[Y, Z]$ or Z^Y instead of $[Y, Z]_C$ if there is no risk of confusion.

Lemma A.2.2. *Let Y be an object in a category C with binary products. The following are equivalent:*

- Y is an exponentiable object in C .
- The functor $- \times Y : C \rightarrow C$ has a right adjoint $[Y, -]_C : C \rightarrow C$, and the counit of this adjunction is $\text{ev}_{Y,-}$.

Proof. Immediate from the definitions. ◆

Definition A.2.3. A **cartesian closed category** is a category with finite products, in which every object is exponentiable. A **locally cartesian closed category** is a category C such that, for every object I , the slice category $C_{/I}$ is a cartesian closed category.

Example A.2.4. **Set** is cartesian closed category; in fact, it is even a locally cartesian closed category.

Proposition A.2.5. *Let C be a cartesian closed category.*

- The assignment $(Y, Z) \mapsto [Y, Z]_C$ extends to a functor $C^{\text{op}} \times C \rightarrow C$.
- For each object Z , the functor $[-, Z]_C : C^{\text{op}} \rightarrow C$ is a contravariant right adjoint for itself.

Proof. (i). This is an instance of the parametrised adjunction theorem.^[2]

(ii). We have the following natural bijections:

$$\begin{aligned} C(X, [Y, Z]) &\cong C(X \times Y, Z) \\ &\cong C(Y \times X, Z) \\ &\cong C(Y, [X, Z]) \end{aligned}$$
■

[2] See Theorem 3 in [CWM, Ch. IV, § 7].

Lemma A.2.6. *Let \mathcal{C} and \mathcal{D} be cartesian closed categories. If $F : \mathcal{C} \rightarrow \mathcal{D}$ is a functor that preserves binary products, then:*

- (i) *For any two objects X and Y in \mathcal{C} , there is a unique morphism $\varphi_{Y,Z} : F[X, Y]_{\mathcal{C}} \rightarrow [FX, FY]_{\mathcal{D}}$ such that the following diagram commutes:*

$$\begin{array}{ccc} F[X, Y]_{\mathcal{C}} \times FX & \xrightarrow{\cong} & F([X, Y]_{\mathcal{C}} \times X) \\ \varphi_{X,Y} \times \text{id} \downarrow & & \downarrow F\text{ev}_{X,Y} \\ [FX, FY]_{\mathcal{D}} \times FX & \xrightarrow{\text{ev}_{FX,FY}} & FY \end{array}$$

- (ii) *The morphism $\varphi_{Y,Z}$ is natural in both Y and Z .*

Proof. The existence and uniqueness of $\varphi_{X,Y}$ follows from the universal property of $[FX, FY]_{\mathcal{D}}$ as an exponential object, and a standard argument proves naturality. \blacklozenge

Definition A.2.7. A **cartesian closed functor** is a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ between cartesian closed categories such that the canonical comparison morphisms $\varphi_{X,Y} : F[X, Y]_{\mathcal{C}} \rightarrow [FX, FY]_{\mathcal{D}}$ described above are isomorphisms.

Proposition A.2.8. *Let \mathcal{C} and \mathcal{D} be cartesian closed categories, and let Y be an object in \mathcal{C} and let Z be an object in \mathcal{D} . Suppose we have an adjunction $F \dashv G : \mathcal{D} \rightarrow \mathcal{C}$ with unit $\eta : \text{id}_{\mathcal{C}} \Rightarrow GF$ and counit $\varepsilon : \text{id}_{\mathcal{D}} \Rightarrow FG$; then:*

- (i) *If $\psi_{FY,Z} : G[FY, Z]_{\mathcal{D}} \rightarrow [GFY, GZ]_{\mathcal{C}}$ is the canonical comparison morphism, then $\theta_{Y,Z} = [\eta_Y, GZ]_{\mathcal{C}} \circ \psi_{FY,Z}$ is the unique morphism in \mathcal{C} making the following diagram commute:*

$$\begin{array}{ccc} G[FY, Z]_{\mathcal{D}} \times Y & \xrightarrow{\text{id} \times \eta_Y} & G[FY, Z]_{\mathcal{D}} \times GFY \\ \theta_{Y,Z} \times \text{id} \downarrow & & \downarrow \cong \\ [Y, GZ]_{\mathcal{C}} \times Y & & G([FY, Z]_{\mathcal{D}} \times FY) \\ \text{ev}_{Y,GZ} \downarrow & & \downarrow G\text{ev}_{FY,Z} \\ GZ & \xlongequal{\quad\quad\quad} & GZ \end{array}$$

- (ii) If the canonical comparison morphism $F(X \times Y) \rightarrow FX \times FY$ is an isomorphism for all objects X in \mathcal{C} , and $\varphi_{Y,GZ} : F[Y, GZ]_{\mathcal{C}} \rightarrow [FY, FGZ]_{\mathcal{D}}$ is the canonical comparison morphism, then $\chi_{Y,Z} = [FY, \varepsilon_Z]_{\mathcal{D}} \circ \varphi_{Y,GZ}$ is the unique morphism in \mathcal{D} making the following diagram commute:

$$\begin{array}{ccc}
 F[Y, GZ]_{\mathcal{C}} \times FY & \xrightarrow{\cong} & F([Y, GZ]_{\mathcal{C}} \times Y) \\
 \chi_{Y,Z} \times \text{id} \downarrow & & \downarrow \text{ev}_{Y,GZ} \\
 [FY, Z]_{\mathcal{D}} \times FY & & FGZ \\
 \text{ev}_{FY,Z} \downarrow & & \downarrow \varepsilon_Z \\
 Z & \xlongequal{\quad\quad\quad} & Z
 \end{array}$$

Moreover, under this hypothesis, $G\chi_{Y,Z} \circ \eta_{[Y,GZ]_{\mathcal{C}}}$ is a two-sided inverse for $\theta_{Y,Z}$.

- (iii) If $\theta_{Y,Z}$ is an isomorphism for all objects Z in \mathcal{D} , then for all objects X in \mathcal{C} , the canonical comparison morphism $F(X \times Y) \rightarrow FX \times FY$ is an isomorphism.

Proof. (i). The claim is proven by the commutativity of the following diagram:

$$\begin{array}{ccccc}
 G[FY, Z]_{\mathcal{D}} \times Y & \xrightarrow{\text{id} \times \eta_Y} & G[FY, Z]_{\mathcal{D}} \times GFY & \xrightarrow{\cong} & G([FY, Z]_{\mathcal{D}} \times FY) \\
 \psi_{FY,Z} \times \text{id} \downarrow & & \downarrow \psi_{FY,Z} \times \text{id} & & \downarrow G\text{ev}_{FY,Z} \\
 [GFY, GZ]_{\mathcal{C}} \times Y & \xrightarrow{\text{id} \times \eta_Y} & [GFY, GZ]_{\mathcal{C}} \times GFY & & \\
 [\eta_Y, Z]_{\mathcal{C}} \downarrow & & \searrow \text{ev}_{GFY,GZ} & & \downarrow \\
 [Y, GZ]_{\mathcal{C}} & \xrightarrow{\quad\quad\quad} & & & GZ \\
 & \text{ev}_{Y,GZ} & & &
 \end{array}$$

- (ii). To show that $\chi_{Y,Z}$ makes the diagram commute, one uses the fact that $\text{ev}_{FY,Z} : [FY, Z]_{\mathcal{D}} \times FY \rightarrow Z$ is natural in Z . Since F preserves products with Y , we have the following natural bijections:

$$\begin{aligned}
 \mathcal{C}(X, G[FY, Z]_{\mathcal{D}}) &\cong \mathcal{D}(FX, [FY, Z]_{\mathcal{D}}) \cong \mathcal{D}(FX \times FY, Z) \\
 &\cong \mathcal{D}(F(X \times Y), Z) \cong \mathcal{C}(X \times Y, GZ) \cong \mathcal{C}(X, [Y, GZ]_{\mathcal{C}})
 \end{aligned}$$

One obtains explicit isomorphisms by chasing id_X in both directions. Taking $X = [Y, GZ]_C$, we find that the isomorphism $[Y, GZ]_C \rightarrow G[FY, Z]_D$ is precisely $G\chi_{Y,Z} \circ \eta_{[Y, GZ]_C}$, and taking $X = G[FY, Z]_D$, we find that the inverse is the right exponential transpose of

$$G(\text{ev}_{FY,Z} \circ (\varepsilon_{[FY,Z]_D} \times \text{id}_Y)) \circ \eta_{G[FY,Z]_D \times Y}$$

where we have suppressed the comparison isomorphism $F(G[FY, Z]_D \times Y) \cong FG[FY, Z]_D \times FY$; but naturality of the comparison morphisms for binary products gives us the commutative diagram below,

$$\begin{array}{ccccc}
G[FY, Z]_D \times Y & \xrightarrow{\eta} & GF(G[FY, Z]_D \times Y) & & \\
\parallel & & \cong \downarrow & & \\
& & G(FG[FY, Z]_D \times FY) & \xrightarrow{G(\varepsilon \times \text{id})} & G([FY, Z]_D \times FY) \\
& & \cong \downarrow & & \downarrow \cong \\
G[FY, Z]_D \times Y & \xrightarrow{\eta \times \eta} & GFG[FY, Z]_D \times GFY & \xrightarrow{G\varepsilon \times \text{id}} & G[FY, Z]_D \times GFY \\
& \searrow \text{id} \times \eta & & \nearrow & \\
& & & &
\end{array}$$

so, suppressing the comparison isomorphisms, we obtain the following equation:

$$G(\varepsilon_{[FY,Z]_D} \times \text{id}_{FY}) \circ \eta_{G[FY,Z]_D \times Y} = \text{id}_{G[FY,Z]_D} \times \eta_Y$$

Thus, the isomorphism $G[FY, Z]_D \rightarrow [GY, Z]_C$ is indeed $\theta_{Y,Z}$, as claimed.

(iii). Now, suppose $\theta_{Y,Z} : G[FY, Z]_D \rightarrow [GY, Z]_C$ is an isomorphism for all Z . Then, we have the natural bijections

$$\begin{aligned}
D(FX \times FY, Z) &\cong D(FX, [FY, Z]_D) \cong C(X, G[FY, Z]_D) \\
&\cong C(X, [Y, GZ]_C) \cong C(X \times Y, GZ) \cong D(F(X \times Y), Z)
\end{aligned}$$

and by chasing id_Z for $Z = FX \times FY$, we conclude that the *canonical* comparison morphism $F(X \times Y) \rightarrow FX \times FY$ is an isomorphism. \blacksquare

Definition A.2.9. A **Frobenius adjunction of cartesian closed categories** is an adjunction $F \dashv G : D \rightarrow C$ where C and D are cartesian closed categories, such that the natural morphisms $\theta_{Y,Z} : G[FY, Z]_D \rightarrow [Y, GZ]_C$ described above are isomorphisms, or equivalently, such that the left adjoint $F : C \rightarrow D$ preserves binary products.

REMARK A.2.10. If \mathcal{C} and \mathcal{D} are cartesian closed categories and $G : \mathcal{D} \rightarrow \mathcal{C}$ is any functor that preserves finite products, then G induces a \mathcal{D} -enrichment of \mathcal{C} from the cartesian closed structure of \mathcal{C} , and the exponential comparison morphisms $\psi_{Y,Z} : G[Y, Z]_{\mathcal{C}} \rightarrow [GY, GZ]_{\mathcal{D}}$ makes $G : \mathcal{D} \rightarrow \mathcal{C}$ into a \mathcal{D} -enriched functor.

Now, suppose G has a left adjoint $F : \mathcal{C} \rightarrow \mathcal{D}$. The adjunction $F \dashv G$ is a Frobenius adjunction precisely when it is compatible with the \mathcal{D} -enrichments of \mathcal{C} and \mathcal{D} . (Of course, this means F is also a \mathcal{D} -enriched functor.)

However, not all enriched adjunctions between cartesian closed categories are of the above form.

Proposition A.2.11. *Let $X, Y,$ and Z be any three objects in a cartesian closed category \mathcal{C} .*

- (i) *There is a unique morphism $\lambda_{X,Y,Z} : [X \times Y, Z] \rightarrow [X, [Y, Z]]$ making the following diagram commute:*

$$\begin{array}{ccc}
 ([X \times Y, Z] \times X) \times Y & \xrightarrow{\cong} & [X \times Y, Z] \times (X \times Y) \\
 (\lambda_{X,Y,Z} \times \text{id}_X) \times \text{id}_Y \downarrow & & \downarrow \text{ev}_{X \times Y, Z} \\
 ([X, [Y, Z]] \times X) \times Y & & \\
 \text{ev}_{X, [Y, Z]} \times \text{id}_X \downarrow & & \downarrow \\
 [Y, Z] \times Y & \xrightarrow{\text{ev}_{Y, Z}} & Z
 \end{array}$$

- (ii) *The morphisms $\lambda_{X,Y,Z} : [X \times Y, Z] \rightarrow [X, [Y, Z]]$ constitute a natural isomorphism.*

Proof. The existence and uniqueness of $\lambda_{X,Y,Z}$ follows from the universal property of $[X, [Y, Z]]$ and $[Y, Z]$ as exponential objects, and a standard argument shows that $\lambda_{X,Y,Z}$ is natural in $X, Y,$ and Z . By the associativity of cartesian products, we have the following natural bijections:

$$\begin{aligned}
 \mathcal{C}(T, [X \times Y, Z]) &\cong \mathcal{C}(T \times (X \times Y), Z) \\
 &\cong \mathcal{C}((T \times X) \times Y, Z) \cong \mathcal{C}(T \times X, [Y, Z]) \cong \mathcal{C}(T, [X, [Y, Z]])
 \end{aligned}$$

Chasing id_T for $T = [X \times Y, Z]$, we find that $\lambda_{X,Y,Z}$ is an isomorphism. ■

Definition A.2.12. Let \mathcal{C} be a cartesian closed category. An **exponential ideal** of \mathcal{C} is a full subcategory $\mathcal{D} \subseteq \mathcal{C}$ such that, for all objects Y in \mathcal{C} , if Z is in \mathcal{D} , then

the exponential object $[Y, Z]_C$ is (isomorphic to) an object in \mathcal{D} . A **reflective exponential ideal** of C is an exponential ideal \mathcal{D} such that the inclusion $\mathcal{D} \hookrightarrow C$ has a left adjoint.

Proposition A.2.13. *Let C be a cartesian closed category, let $G : \mathcal{D} \rightarrow C$ be the inclusion of a full subcategory, and suppose G has a left adjoint $F : C \rightarrow \mathcal{D}$. The following are equivalent:*

- (i) *F preserves finite products.*
- (ii) *F preserves binary products.*
- (iii) *\mathcal{D} is a reflective exponential ideal of C .*
- (iv) *\mathcal{D} is a cartesian closed category, $G : \mathcal{D} \rightarrow C$ is a cartesian closed functor, and the canonical morphisms $G[FY, Z]_{\mathcal{D}} \rightarrow [Y, GZ]_C$ are isomorphisms.*

Proof. (i) \Rightarrow (ii). Immediate.

(ii) \Rightarrow (iii). Under our hypotheses, the product of two objects X and Y in \mathcal{D} can be computed as $F(GX \times GY)$. Let $\eta : \text{id}_C \rightarrow GF$ be the unit of the adjunction. We have the following natural bijections:

$$\begin{aligned} C(X, [Y, GZ]_C) &\cong C(X \times Y, GZ) \\ &\cong \mathcal{D}(FX \times FY, Z) \\ &\cong \mathcal{D}(FGFX \times FY, Z) \\ &\cong C(GFX \times Y, GZ) \\ &\cong C(GFX, [Y, GZ]_C) \end{aligned}$$

By chasing these maps explicitly, we find that every morphism $X \rightarrow [Y, GZ]_C$ factors through $\eta_X : X \rightarrow GFX$ in a unique way. In particular, we have

$$\text{id}_{[Y, GZ]_C} = r_{Y, Z} \circ \eta_{[Y, GZ]_C}$$

for a unique $r_{Y, Z} : GF[Y, GZ]_C \rightarrow [Y, GZ]_C$. The triangle identity then implies $Fr_{Y, Z} = \varepsilon_{F[Y, GZ]_C}$, thus,

$$\eta_{[Y, GZ]_C} \circ r_{Y, Z} = GF r_{Y, Z} \circ \eta_{GF[Y, GZ]_C} = G\varepsilon_{F[Y, GZ]_C} \circ \eta_{GF[Y, GZ]_C} = \text{id}_{GF[Y, GZ]_C}$$

and therefore $r_{Y, Z}$ is an isomorphism.

(iii) \Rightarrow (iv). It is a standard fact that a reflective subcategory is closed under all limits that exist in \mathcal{C} , so \mathcal{D} must have finite products and $G : \mathcal{D} \rightarrow \mathcal{C}$ preserves them. If \mathcal{D} is an exponential ideal, then $\eta_{[Y, GZ]_{\mathcal{C}}} : [Y, GZ]_{\mathcal{C}} \rightarrow GF[Y, GZ]_{\mathcal{C}}$ must be an isomorphism, so we obtain natural bijections

$$\begin{aligned} \mathcal{D}(X \times Y, Z) &\cong \mathcal{C}(GX \times GY, GZ) \\ &\cong \mathcal{C}(GX, [GY, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(GX, GF[GY, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{D}(FGX, F[GY, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{D}(X, F[GY, GZ]_{\mathcal{C}}) \end{aligned}$$

and therefore we may take $[Y, Z]_{\mathcal{D}} = F[GY, GZ]_{\mathcal{C}}$. Obviously, this makes $G : \mathcal{D} \rightarrow \mathcal{C}$ into a cartesian closed functor. We also have

$$\begin{aligned} \mathcal{C}(X, G[FY, Z]_{\mathcal{D}}) &= \mathcal{C}(X, GF[GFY, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(X, [GFY, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(GFY, [X, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(GFY, GF[X, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(Y, GF[X, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(Y, [X, GZ]_{\mathcal{C}}) \\ &\cong \mathcal{C}(X, [Y, GZ]_{\mathcal{C}}) \end{aligned}$$

and so the canonical morphism $G[FY, Z]_{\mathcal{D}} \rightarrow [Y, GZ]_{\mathcal{C}}$ is an isomorphism.

(iv) \Rightarrow (i). It is not hard to show that $\eta_1 : 1 \rightarrow GF1$ is an isomorphism for any adjunction whatsoever; but G is fully faithful, so this implies $F1$ is a terminal object in \mathcal{D} . Now apply [proposition A.2.8](#). ■

Corollary A.2.14. *If \mathcal{E} is a reflective exponential ideal of \mathcal{D} , and \mathcal{D} is a reflective exponential ideal of \mathcal{C} , then \mathcal{E} is also a reflective exponential ideal of \mathcal{C} .* ■

Proposition A.2.15. *Let \mathbf{Cat} be the category of small categories, and let \mathbf{Grpd} be the full subcategory of groupoids.*

(i) *There exist adjunctions*

$$\pi_0 \dashv \text{disc} \dashv \text{ob} \dashv \text{codisc} : \mathbf{Set} \rightarrow \mathbf{Cat}$$

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where $\text{ob } \mathbb{C}$ is the set of objects in a category \mathbb{C} , $\text{disc } X$ is the category with $\text{ob } \text{disc } X = X$ and all arrows trivial, and $\text{codisc } X$ is the category with $\text{ob } \text{codisc } X = X$ and a unique arrow between any two objects.

- (ii) The functor $\text{disc} : \mathbf{Set} \rightarrow \mathbf{Cat}$ is fully faithful and exhibits \mathbf{Set} as a reflective exponential ideal of \mathbf{Cat} .
- (iii) The functor $\pi_0 : \mathbf{Cat} \rightarrow \mathbf{Set}$ preserves finite products.
- (iv) There exist adjunctions

$$\mathbf{I} \dashv \text{und} \dashv \text{iso} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$$

where $\text{und} : \mathbf{Grpd} \rightarrow \mathbf{Cat}$ is the inclusion and $\text{iso } \mathbb{C}$ is the maximal subgroupoid of a category \mathbb{C} .

- (v) \mathbf{Grpd} is a reflective exponential ideal of \mathbf{Cat} .
- (vi) The functor $\mathbf{I} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$ preserves finite products.
- (vii) The adjunctions in (i) factor through \mathbf{Grpd} , yielding adjunctions

$$\pi_0 \dashv \text{disc} \dashv \text{ob} \dashv \text{codisc} : \mathbf{Set} \rightarrow \mathbf{Grpd}$$

where $\pi_0 : \mathbf{Grpd} \rightarrow \mathbf{Set}$ again preserves finite products.

- (viii) The functor $\mathbf{Cat} \rightarrow \mathbf{Set}$ that sends a category \mathbb{C} to the set of isomorphism classes in \mathbb{C} preserves finite products.

Proof. (i). The functor $\text{disc} : \mathbf{Set} \rightarrow \mathbf{Cat}$ obviously satisfies the solution set condition, so the general adjoint functor theorem gives us a left adjoint $\pi_0 : \mathbf{Cat} \rightarrow \mathbf{Set}$; the existence of the other adjunctions is obvious.

(ii). It is clear that $\text{disc} : \mathbf{Set} \rightarrow \mathbf{Cat}$ is fully faithful, and direct computation shows that $[\mathbb{C}, \text{disc } X]$ is a discrete category for any \mathbb{C} , so \mathbf{Set} is indeed a reflective exponential ideal of \mathbf{Cat} .

(iii). Thus, by [proposition A.2.13](#), $\pi_0 : \mathbf{Cat} \rightarrow \mathbf{Set}$ must preserve finite products.

(iv). It is not hard to check that the inclusion $\mathbf{Grpd} \rightarrow \mathbf{Cat}$ satisfies the solution set condition, so the general adjoint functor theorem gives us a left adjoint $\mathbf{I} :$

Cat \rightarrow **Grpd**; the fact that $\text{iso} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$ is right adjoint to the inclusion is obvious.

(v). Direct computation shows that $[\mathbb{C}, \mathbb{G}]$ is a groupoid whenever \mathbb{G} is, so **Grpd** is indeed a reflective exponential ideal of **Cat**.

(vi). Thus, $\mathbf{I} : \mathbf{Cat} \rightarrow \mathbf{Grpd}$ must preserve finite products.

(vii). Clearly, $\text{disc } X$ and $\text{codisc } X$ are already groupoids, so the adjunctions do indeed factor through **Grpd**.

(viii). The set of isomorphism classes of objects in \mathbb{C} is precisely $\pi_0 \text{ iso } \mathbb{C}$. \blacksquare

Definition A.2.16. The **dependent sum** of an object $p : X \rightarrow I$ in $\mathcal{C}_{/I}$ along a morphism $j : I \rightarrow J$ in \mathcal{C} is the object $j \circ p : X \rightarrow J$ in $\mathcal{C}_{/J}$, and we write $\Sigma_j : \mathcal{C}_{/I} \rightarrow \mathcal{C}_{/J}$ for the functor sending an object to its dependent sum along j .

Lemma A.2.17. *Let $j : I \rightarrow J$ be a morphism in a category \mathcal{C} . The following are equivalent:*

- (i) \mathcal{C} has pullbacks along j .
- (ii) There exists an adjunction

$$\Sigma_j \dashv j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$$

where Σ_j is the dependent sum functor, and the right adjoint $j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$ is the pullback functor.

Proof. This is just a matter of unwinding the definitions. \blacklozenge

Definition A.2.18. Let \mathcal{C} be a category with pullbacks. A **dependent product** of an object $p : X \rightarrow I$ in $\mathcal{C}_{/I}$ along a morphism $j : I \rightarrow J$ in \mathcal{C} is an object $\Pi_j p$ in $\mathcal{C}_{/J}$ and a morphism $\text{ev}_{j,p} : j^* \Pi_j p \rightarrow p$ in $\mathcal{C}_{/I}$ with the following universal property:

- For all morphisms $f : j^* q \rightarrow p$ in $\mathcal{C}_{/I}$, there exists a unique morphism $\tilde{f} : q \rightarrow \Pi_j p$ in $\mathcal{C}_{/J}$ such that $\text{ev}_{j,p} \circ j^* \tilde{f} = f$.

A $\Sigma\Pi$ -category is a category \mathcal{C} with finite limits such that, for every morphism $j : I \rightarrow J$ in \mathcal{C} , dependent products along j exist.

Lemma A.2.19. *Let $j : I \rightarrow J$ be a morphism in a category \mathcal{C} with pullbacks. The following are equivalent:*

- (i) *For all objects $p : X \rightarrow I$ in \mathcal{C} , a dependent product of p along j exists.*
- (ii) *The pullback functor $j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$ has a right adjoint $\Pi_j : \mathcal{C}_{/I} \rightarrow \mathcal{C}_{/J}$ that sends an object to its dependent product along j , and the counit of this adjunction is $\text{ev}_{j,-}$.*

Proof. This is just a matter of unwinding the definitions. ◆

Corollary A.2.20. *If $j : I \rightarrow J$ is a morphism in a $\Sigma\Pi$ -category \mathcal{C} , then the pullback functor $j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$ preserves all limits and colimits.* ■

Proposition A.2.21. *Let \mathcal{C} be a category with a terminal object. The following are equivalent:*

- (i) *\mathcal{C} is a $\Sigma\Pi$ -category.*
- (ii) *\mathcal{C} is a locally cartesian closed category.*

Proof. See Proposition 9.20 in [Awodey, 2010]. □

Theorem A.2.22. *Let \mathbb{D} be a small category, and let $\mathcal{C} = [\mathbb{D}^{\text{op}}, \mathbf{Set}]$. Then:*

- (i) *\mathcal{C} has limits and colimits for all small diagrams, and these can be constructed componentwise in \mathbf{Set} : a cone (resp. cocone) in \mathcal{C} over (resp. under) a diagram in \mathcal{C} is a limiting cone (resp. colimiting cocone) if and only if it is so in every component.*
- (ii) *Every internal equivalence relation in \mathcal{C} is the kernel pair of its coequaliser.*
- (iii) *For all morphisms $j : I \rightarrow J$ in \mathcal{C} , the pullback functor $j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$ preserves all limits and colimits.*
- (iv) *The Yoneda embedding $\mathfrak{h}_\bullet : \mathbb{D} \rightarrow \mathcal{C}$ is a dense functor, i.e. for every presheaf $X : \mathbb{D}^{\text{op}} \rightarrow \mathbf{Set}$, the tautological cocone^[3] from the canonical diagram $(\mathfrak{h}_\bullet \downarrow X) \rightarrow \mathcal{C}$ to X is a colimiting cocone.*
- (v) *\mathcal{C} is a locally finitely presentable category.*

[3] See definition A.5.7.

(vi) \mathcal{C} is a $\Sigma\Pi$ -category.

Proof. (i). This is a standard fact about presheaf categories.

(ii) and (iii). The claims are true for **Set**, and hence for \mathcal{C} by claim (i).

(iv). See [proposition A.5.24](#).

(v). See [theorem 0.2.26](#).

(vi). Apply [theorem 0.2.35](#) to construct a right adjoint for $j^* : \mathcal{C}_{/J} \rightarrow \mathcal{C}_{/I}$. ■

REMARK A.2.23. The Yoneda lemma gives us an explicit description of the exponential objects in $[\mathbb{D}^{\text{op}}, \mathbf{Set}]$: given two presheaves $Y, Z : \mathbb{D}^{\text{op}} \rightarrow \mathbf{Set}$, if Z^Y is their exponential object, then we must have

$$Z^Y(d) \cong [\mathbb{D}^{\text{op}}, \mathbf{Set}](h_d, Z^Y) \cong [\mathbb{D}^{\text{op}}, \mathbf{Set}](h_d \times Y, Z)$$

and so we may define Y^Z by $Y^Z(d) = [\mathbb{D}^{\text{op}}, \mathbf{Set}](h_d \times Y, Z)$.

Definition A.2.24. Let Y and Z be topological spaces, and let $[Y, Z]$ be the set of all *continuous* maps $Y \rightarrow Z$. The **compact–open topology** on $[Y, Z]$ is the coarsest topology such that the subsets

$$V(K, U) = \{f \in [Y, Z] \mid K \subseteq f^{-1}U\}$$

are open in $[Y, Z]$ for all compact subsets $K \subseteq X$ and all open subsets $U \subseteq Y$.

REMARK A.2.25. If Y is a discrete space, then the compact–open topology on $[Y, Z]$ coincides with the product topology on Z^Y .

Definition A.2.26. A **compactly-generated Hausdorff space** is a Hausdorff topological space X such that a subset $U \subseteq X$ is open if and only if, for every continuous map $f : K \rightarrow X$ where K is a compact Hausdorff space, $f^{-1}U$ is an open subset of K . We write **CGHaus** for the category of compactly-generated Hausdorff spaces and continuous maps.

Proposition A.2.27.

- (i) *If Y is a locally compact Hausdorff space, then for all topological spaces Z , the set of all continuous maps $Y \rightarrow Z$, equipped with the compact–open topology, is an exponential object $[Y, Z]$ in **Top**.*

(ii) **Top** is not a cartesian closed category.

(iii) **CGHaus** is a cartesian closed category.

Proof. Claim (i) follows from Theorems 46.10 and 46.11 in [Munkres, 2000], and claim (ii) is Proposition 7.1.2 in [Borceux, 1994a], and claim (iii) is proved in [GZ, Ch. III, § 2]. \square

A.3 Factorisation systems

Definition A.3.1. Let $f : X \rightarrow Y$ and $g : Z \rightarrow W$ be morphisms in a category \mathcal{C} . Given a commutative square in \mathcal{C} ,

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

a **lift** is a morphism $h : W \rightarrow X$ such that $f \circ h = w$ and $h \circ g = z$.

We say g has the **left lifting property** with respect to f and f has the **right lifting property** with respect to g , and we write $g \sqsupseteq f$, if every commutative square in \mathcal{C} of the form above has a lift. We say f is **left orthogonal** to g and g is **right orthogonal** to f , and we write $g \perp f$ if lifts exist *and* are unique.

Given $\mathcal{I} \subseteq \text{mor } \mathcal{C}$, we define the following subensembles of $\text{mor } \mathcal{C}$:

$$\begin{aligned} \sqsupseteq \mathcal{I} &= \{f \in \text{mor } \mathcal{C} \mid \forall g \in \mathcal{I}. f \sqsupseteq g\} \\ \mathcal{I} \sqsupseteq &= \{g \in \text{mor } \mathcal{C} \mid \forall f \in \mathcal{I}. f \sqsupseteq g\} \\ \perp \mathcal{I} &= \{f \in \text{mor } \mathcal{C} \mid \forall g \in \mathcal{I}. f \perp g\} \\ \mathcal{I}^\perp &= \{g \in \text{mor } \mathcal{C} \mid \forall f \in \mathcal{I}. f \perp g\} \end{aligned}$$

Lemma A.3.2. Let $f : X \rightarrow Y$ and $g : Z \rightarrow W$ be morphisms in a locally small category \mathcal{C} . Consider the commutative diagram in **Set** shown below,

$$\begin{array}{ccc} \mathcal{C}(W, X) & \xrightarrow{f_*} & \mathcal{C}(W, Y) \\ \downarrow g^* & \dashrightarrow & \downarrow g^* \\ \bullet & \longrightarrow & \mathcal{C}(Z, Y) \\ \downarrow & & \downarrow f_* \\ \mathcal{C}(Z, X) & \xrightarrow{f_*} & \mathcal{C}(Z, Y) \end{array}$$

where the inner square is a pullback diagram.

- (i) The dashed arrow is a surjection if and only if $g \sqsupseteq f$.
- (ii) The dashed arrow is a bijection if and only if $g \perp f$.

Proof. This is just a restatement of the definition. ■

Proposition A.3.3. *Let \mathcal{C} be a category.*

- (i) If $\mathcal{R} \subseteq \text{mor } \mathcal{C}$, then ${}^\perp\mathcal{R} \subseteq \sqsupseteq\mathcal{R}$.
- (ii) If $\mathcal{R}' \subseteq \mathcal{R} \subseteq \text{mor } \mathcal{C}$, then $\sqsupseteq\mathcal{R}' \supseteq \sqsupseteq\mathcal{R}$.
- (iii) If $\mathcal{R}' \subseteq \mathcal{R} \subseteq \text{mor } \mathcal{C}$, then ${}^\perp\mathcal{R}' \supseteq {}^\perp\mathcal{R}$.

Dually:

- (i') If $\mathcal{L} \subseteq \text{mor } \mathcal{C}$, then $\mathcal{L}^\perp \subseteq \mathcal{L}^\sqsupseteq$.
- (ii') If $\mathcal{L}' \subseteq \mathcal{L} \subseteq \text{mor } \mathcal{C}$, then $\mathcal{L}'^\sqsupseteq \supseteq \mathcal{L}^\sqsupseteq$.
- (iii') If $\mathcal{L}' \subseteq \mathcal{L} \subseteq \text{mor } \mathcal{C}$, then $\mathcal{L}'^\perp \supseteq \mathcal{L}^\perp$.

Moreover, we have the following antitone Galois connections:

$$\begin{aligned} \mathcal{L} \subseteq \sqsupseteq\mathcal{R} \text{ if and only if } \mathcal{R} \subseteq \mathcal{L}^\sqsupseteq \\ \mathcal{L} \subseteq {}^\perp\mathcal{R} \text{ if and only if } \mathcal{R} \subseteq \mathcal{L}^\perp \end{aligned}$$

Proof. Obvious. ◆

Corollary A.3.4. *We have the following identities:*

$$\begin{aligned} \sqsupseteq((\sqsupseteq\mathcal{R})^\sqsupseteq) &= \sqsupseteq\mathcal{R} & {}^\perp(({}^\perp\mathcal{R})^\perp) &= {}^\perp\mathcal{R} \\ (\sqsupseteq(\mathcal{L}^\sqsupseteq))^\sqsupseteq &= \mathcal{L}^\sqsupseteq & ({}^\perp(\mathcal{L}^\perp))^\perp &= \mathcal{L}^\perp \end{aligned}$$

Proof. This is a standard fact about (antitone) Galois connections. ■

Lemma A.3.5. *Let $f : X \rightarrow Y$ be a morphism in a category \mathcal{C} . The following are equivalent:*

- (i) f is an isomorphism.
- (ii) f is right orthogonal to any morphism in \mathcal{C} .

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(iii) f has the right lifting property with respect to any morphism in \mathcal{C} .

(iv) f has the right lifting property with respect to itself.

Dually, the following are equivalent:

(i') f is an isomorphism.

(ii') f is left orthogonal to any morphism in \mathcal{C} .

(iii') f has the left lifting property with respect to any morphism in \mathcal{C} .

(iv') f has the left lifting property with respect to itself.

Proof. (i) \Rightarrow (ii). Suppose $r : Y \rightarrow X$ is a morphism such that $r \circ f = \text{id}_X$. Then, for any commutative square as below,

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

we have $(r \circ w) \circ g = r \circ f \circ z = z$; but if $f \circ r = \text{id}_Y$ as well, then $f \circ (r \circ w) = w$; thus $r \circ w : W \rightarrow X$ is the required lift. It is clearly unique, as f is monic.

(ii) \Rightarrow (iii), (iii) \Rightarrow (iv). Obvious.

(iv) \Rightarrow (i). Consider the following commutative square:

$$\begin{array}{ccc} X & \xrightarrow{\text{id}} & X \\ f \downarrow & & \downarrow f \\ Y & \xrightarrow{\text{id}} & Y \end{array}$$

Since f has the right lifting property with respect to itself, there exists a morphism $h : Y \rightarrow X$ such that $h \circ f = \text{id}_X$ and $f \circ h = \text{id}_Y$. \blacksquare

Definition A.3.6. A **weak factorisation system** for a category \mathcal{C} is a pair $(\mathcal{L}, \mathcal{R})$ of subclasses of $\text{mor } \mathcal{C}$ satisfying these conditions:

- For each morphism f in \mathcal{C} there exists a pair (g, h) with $g \in \mathcal{L}$ and $h \in \mathcal{R}$ such that $f = h \circ g$. Such a pair is a **$(\mathcal{L}, \mathcal{R})$ -factorisation** of f .

- A morphism is in \mathcal{L} if and only if it has the left lifting property with respect to every morphism in \mathcal{R} , i.e. $\mathcal{L} = \square\mathcal{R}$.
- A morphism is in \mathcal{R} if and only if it has the right lifting property with respect to every morphism in \mathcal{L} , i.e. $\mathcal{R} = \mathcal{L}\square$.

An **orthogonal factorisation system** is defined like a weak factorisation system, except for replacing ‘... has the left/right lifting property with respect to ...’ with ‘... is left/right orthogonal to ...’.

REMARK A.3.7. Obviously, $(\mathcal{L}, \mathcal{R})$ is a weak (resp. orthogonal) factorisation system for \mathcal{C} if and only if $(\mathcal{R}^{\text{op}}, \mathcal{L}^{\text{op}})$ is a weak (resp. orthogonal) factorisation system for \mathcal{C}^{op} .

Proposition A.3.8. *Let $(\mathcal{L}, \mathcal{R})$ be a weak factorisation system on \mathcal{C} . If either*

- *every morphism in \mathcal{R} is a monomorphism in \mathcal{C} , or*
- *every morphism in \mathcal{L} is an epimorphism in \mathcal{C} ,*

then $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system.

Proof. The two hypotheses are formally dual, so it is enough to check the first case. Observe that, given a commutative diagram

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & \nearrow h & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

where $f : X \rightarrow Y$ is a monomorphism, for any $h' : W \rightarrow X$ such that $f \circ h' = w$, we must have $h = h'$. Thus, for any monomorphism $f : X \rightarrow Y$, $g \square f$ if and only if $g \perp f$. Hence, $\mathcal{L} = \square\mathcal{R} = {}^\perp\mathcal{R}$. On the other hand, $\mathcal{L}^\perp \subseteq \mathcal{L}\square = \mathcal{R}$, so $\mathcal{R} = \mathcal{L}^\perp$ as well. ■

Definition A.3.9. A **proper factorisation system** on a category \mathcal{C} is an orthogonal factorisation system $(\mathcal{E}, \mathcal{M})$ on \mathcal{C} such that every morphism in \mathcal{E} is an epimorphism *and* every morphism in \mathcal{M} is a monomorphism.

Example A.3.10. In **Set**, if \mathcal{E} is the class of surjective maps and \mathcal{M} is the class of injective maps, then $(\mathcal{E}, \mathcal{M})$ is a proper factorisation system.

Lemma A.3.11. *Let A be an object in a category \mathcal{C} with a weak (resp. orthogonal) factorisation system $(\mathcal{L}, \mathcal{R})$. Then the slice category $\mathcal{C}_{/A}$ has a weak (resp. orthogonal) factorisation system where a morphism is in the left or right class if and only if it is so in \mathcal{C} .*

Proof. The projection $\mathcal{C}_{/A} \rightarrow \mathcal{C}$ induces a bijection between solutions for lifting problems in $\mathcal{C}_{/A}$ and solutions for the corresponding lifting problems in \mathcal{C} . ■

Proposition A.3.12 (Closure properties). *Let $\mathcal{R} \subseteq \text{mor } \mathcal{C}$ and suppose either $\mathcal{L} = \square\mathcal{R}$ or $\mathcal{L} = {}^\perp\mathcal{R}$.*

(i) *Given a pushout diagram in \mathcal{C} as below,*

$$\begin{array}{ccc} Z' & \xrightarrow{i_Z} & Z \\ g' \downarrow & & \downarrow g \\ W' & \xrightarrow{i_W} & W \end{array}$$

if the morphism g' is in \mathcal{L} , then g is also in \mathcal{L} .

(ii) *Let I be a set. If $g_i : Z_i \rightarrow W_i$ is a morphism in \mathcal{L} for all i in I and the coproduct $\coprod_i g_i : \coprod_i Z_i \rightarrow \coprod_i W_i$ exists in \mathcal{C} , then $\coprod_i g_i$ is also in \mathcal{L} .*

(iii) *Given a commutative diagram of the form*

$$\begin{array}{ccccc} & & \text{id} & & \\ & & \curvearrowright & & \\ Z & \xrightarrow{i_Z} & Z' & \xrightarrow{r_Z} & Z \\ g' \downarrow & & \downarrow g & & \downarrow g' \\ W' & \xrightarrow{i_W} & W & \xrightarrow{r_W} & W' \\ & & \text{id} & & \\ & & \curvearrowleft & & \end{array}$$

if g is in \mathcal{L} , then so is g' ; in other words, \mathcal{L} is closed under retracts.

(iv) \mathcal{L} is closed under composition.

(v) *Let γ be an ordinal and let $Z : \gamma \rightarrow \mathcal{C}$ be a colimit-preserving functor. We write Z_α for $Z(\alpha)$, where $\alpha < \gamma$, and $g_{\alpha,\beta} : Z_\alpha \rightarrow Z_\beta$ for the morphism $Z(\alpha \rightarrow \beta)$, where $\alpha < \beta < \gamma$. If λ is a colimiting cocone from Z to W and each $g_{\alpha,\beta}$ is in \mathcal{L} , then each component $\lambda_\alpha : Z_\alpha \rightarrow W$ is also in \mathcal{L} .*

Proof. (i). Suppose f is in \mathcal{R} , and consider the following commutative diagram:

$$\begin{array}{ccccc} Z' & \xrightarrow{i_Z} & Z & \xrightarrow{z} & X \\ g' \downarrow & & \downarrow g & & \downarrow f \\ W' & \xrightarrow{i_W} & W & \xrightarrow{w} & Y \end{array}$$

There exists $h' : W' \rightarrow X$ such that $h' \circ g' = z \circ i_Z$ and $f \circ h' = w \circ i_W$. In particular, there exists a unique morphism $h : W \rightarrow X$ such that $h \circ g = z$ and $h \circ i_W = h'$, by the universal property of pullbacks. Thus $f \circ h \circ i_W = f \circ h' = w \circ i_W$ and $f \circ h \circ g = f \circ z = w \circ g$, but i_W and g are jointly epic, so $f \circ h = w$. This shows h is the required lift, and h is unique if h' is.

(ii). We may construct the required lift componentwise.

(iii). Suppose f is in \mathcal{R} , and consider the following commutative diagram:

$$\begin{array}{ccccccc} Z & \xrightarrow{i_Z} & Z' & \xrightarrow{r_Z} & Z & \xrightarrow{z} & X \\ g' \downarrow & & \downarrow g & & \downarrow g' & & \downarrow f \\ W' & \xrightarrow{i_W} & W & \xrightarrow{r_W} & W' & \xrightarrow{w} & Y \end{array}$$

There exists $h : W \rightarrow X$ such that $h \circ g = z \circ r_Z$ and $f \circ h = w \circ r_W$, and so for $h' = h \circ i_W$:

$$\begin{aligned} f \circ h' &= f \circ h \circ i_W = w \circ r_W \circ i_W = w \\ h' \circ g' &= h \circ i_W \circ g' = h \circ g \circ i_Z = z \circ r_Z \circ i_Z = z \end{aligned}$$

Thus $h' : W' \rightarrow X$ is the required lift, and h' is unique if h is (because r_W is split epic).

(iv). Suppose $g' : Z' \rightarrow Z$ and $g : Z \rightarrow W$ are in \mathcal{L} and $f : X \rightarrow Y$ is in \mathcal{R} . Consider the following commutative diagram:

$$\begin{array}{ccc} Z' & \xrightarrow{z'} & X \\ g' \downarrow & & \downarrow f \\ Z & & \\ g \downarrow & & \downarrow \\ W & \xrightarrow{w} & Y \end{array}$$

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There must exist a morphism $z : Z \rightarrow X$ such that $z \circ g' = z'$ and $f \circ z' = w \circ g$, and hence a morphism $h : W \rightarrow X$ such that $h \circ g = z$ and $f \circ h = w$. Obviously, $h \circ (g' \circ g) = z'$, so h is the required lift. Moreover, h unique if $\mathcal{L} = {}^\perp\mathcal{R}$.

(v). We may assume without loss of generality that $\alpha = 0$, since any non-empty terminal segment of γ is cofinal in γ . Suppose $f : X \rightarrow Y$ is in \mathcal{R} and consider the following commutative diagram:

$$\begin{array}{ccc} Z_0 & \xrightarrow{z_0} & X \\ \lambda_0 \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

For each $\alpha < \gamma$, given z_α making the following diagram commute,

$$\begin{array}{ccccc} Z_\alpha & \xrightarrow{z_\alpha} & & X & \\ g_{\alpha,\alpha+1} \downarrow & & & \downarrow f & \\ Z_{\alpha+1} & \xrightarrow{\lambda_{\alpha+1}} & W & \xrightarrow{w} & Y \end{array}$$

choose a lift $z_{\alpha+1} : Z_{\alpha+1} \rightarrow X$; for each limit ordinal $\beta < \gamma$, let $z_\beta : Z_\beta \rightarrow X$ be the unique morphism such that $z_\beta \circ g_{\alpha,\beta} = z_\alpha$ for all $\alpha < \beta$. (Such z_β exist and are unique because $Z_\beta = \varinjlim_{\alpha < \beta} Z_\alpha$.) Note that the universal property of W then guarantees that $w \circ \lambda_\beta = f \circ z_\beta$.

Having constructed morphisms $z_\alpha : Z_\alpha \rightarrow X$ for all $\alpha < \gamma$ as above, we may now obtain $h : W \rightarrow X$ as the unique morphism such that $h \circ \lambda_\alpha = z_\alpha$ for all $\alpha < \gamma$, and again we automatically have $f \circ h = w$. It is also clear that h is unique if $\mathcal{L} = {}^\perp\mathcal{R}$. ■

Proposition A.3.13 (Cancellation properties). *Let $\mathcal{R} \subseteq \text{mor } C$.*

- (i) *Let \mathcal{L} be either $\square\mathcal{R}$ or ${}^\perp\mathcal{R}$, let $e : A \rightarrow Z$ be an epimorphism in C , and let $g : Z \rightarrow W$ be a morphism in C . If $g \circ e$ is in \mathcal{L} , then so is g .*
- (ii) *Suppose $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on \mathcal{R} , and let $e : A \rightarrow Z$ be in \mathcal{L} . Then, a morphism $g : Z \rightarrow W$ is in \mathcal{L} if and only if $g \circ e$ is in \mathcal{L} .*

Dually, let $\mathcal{L} \subseteq \text{mor } C$.

- (i') Let \mathcal{R} be either \mathcal{L}^{\square} or \mathcal{L}^{\perp} , let $m : Y \rightarrow B$ be an monomorphism in \mathcal{C} , and let $f : X \rightarrow Y$ be a morphism in \mathcal{C} . If $m \circ f$ is in \mathcal{R} , then so is f .
- (ii') Suppose $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on \mathcal{R} , and let $m : Y \rightarrow B$ be in \mathcal{L} . Then, a morphism $f : X \rightarrow Y$ is in \mathcal{L} if and only $g \circ e$ is in \mathcal{L} .

Proof. (i). The epimorphism $e : A \rightarrow Z$ induces a bijection between solutions of lifting problems in \mathcal{C} of the form

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

and lifting problems of the form

$$\begin{array}{ccc} A & \xrightarrow{z \circ e} & X \\ g \circ e \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

so $g \square f$ (resp. $g \perp f$) if and only if $g \circ e \square f$ (resp. $g \circ e \perp f$).

(ii). By [proposition A.3.12](#), we know $g \circ e$ is in \mathcal{L} if both g and e are in \mathcal{L} ; the converse remains to be shown. Let $r \circ l$ be an $(\mathcal{L}, \mathcal{R})$ -factorisation of g . If $g \circ e$ is in \mathcal{L} , then there exists a unique s making the diagram below commute,

$$\begin{array}{ccc} A & \xrightarrow{l \circ e} & M \\ g \circ e \downarrow & \nearrow s & \downarrow r \\ W & \xrightarrow{\text{id}} & W \end{array}$$

so $r \circ s = \text{id}_W$, but then we also have

$$\begin{aligned} r \circ (s \circ r) &= r \\ (s \circ r) \circ (l \circ e) &= s \circ (g \circ e) = l \circ e \end{aligned}$$

and $l \circ e \perp r$, so we must have $s \circ r = \text{id}_M$. Hence, g is also in \mathcal{L} . ■

Proposition A.3.14. *Every orthogonal factorisation system is also a weak factorisation system.*

Proof. Let $(\mathcal{L}, \mathcal{R})$ be an orthogonal factorisation system on a category \mathcal{C} . [Proposition A.3.3](#) implies $\mathcal{L} \subseteq \square\mathcal{R}$ and $\mathcal{R} \subseteq \mathcal{L}^\square$, so by duality it is enough to check that $\mathcal{L} \supseteq \square\mathcal{R}$.

Suppose $g : Z \rightarrow W$ is in $\square\mathcal{R}$, with $(\mathcal{L}, \mathcal{R})$ -factorisation $g = r \circ g'$. Then the diagram below commutes,

$$\begin{array}{ccc} Z & \xrightarrow{g'} & W' \\ g \downarrow & & \downarrow r \\ W & \xrightarrow{\text{id}} & W \end{array}$$

so there must exist $i : W \rightarrow W'$ such that $r \circ i = \text{id}_W$ and $i \circ g = g'$, and hence we have the following commutative diagram:

$$\begin{array}{ccccc} Z & \xrightarrow{\text{id}} & Z & \xrightarrow{\text{id}} & Z \\ g \downarrow & & g' \downarrow & & \downarrow g' \\ W & \xrightarrow{i} & W' & \xrightarrow{r} & W \\ & \searrow & \text{id} & \nearrow & \\ & & & & \end{array}$$

It follows from [proposition A.3.12](#) that g is also in \mathcal{L} , so $\mathcal{L} \supseteq \square\mathcal{R}$ as required. \blacksquare

Definition A.3.15. A weak factorisation system $(\mathcal{L}, \mathcal{R})$ on a category \mathcal{C} is **cofibrantly generated** by a subensemble $\mathcal{I} \subseteq \text{mor } \mathcal{C}$ if $\mathcal{R} = \mathcal{I}^\square$. Dually, $(\mathcal{L}, \mathcal{R})$ is **fibrantly generated** by a subensemble $\mathcal{F} \subseteq \text{mor } \mathcal{C}$ if $\mathcal{L} = \square\mathcal{F}$.

REMARK A.3.16. Of course, $(\mathcal{L}, \mathcal{R})$ is always cofibrantly generated by \mathcal{L} . The condition is most useful when $(\mathcal{L}, \mathcal{R})$ is cofibrantly generated by a (small) subset of \mathcal{L} , but it is convenient to have the more general definition available.

Definition A.3.17. Let $(\mathcal{L}, \mathcal{R})$ be a weak factorisation system on a category \mathcal{C} . An **extension** of $(\mathcal{L}, \mathcal{R})$ along a functor $i : \mathcal{C} \rightarrow \mathcal{C}^+$ is a weak factorisation system $(\mathcal{L}^+, \mathcal{R}^+)$ on \mathcal{C}^+ with the following properties:

- A morphism $f : X \rightarrow Y$ in \mathcal{C} is in \mathcal{R} if and only if $if : iX \rightarrow iY$ is in \mathcal{R}^+ .
- A morphism $g : Z \rightarrow W$ in \mathcal{C} is in \mathcal{L} if and only if $ig : iZ \rightarrow iW$ is in \mathcal{L}^+ .

Proposition A.3.18. *Let \mathcal{C} be a full subcategory of a category \mathcal{C}^+ , let $(\mathcal{L}, \mathcal{R})$ be a weak factorisation system on \mathcal{C} , and let $(\mathcal{L}^+, \mathcal{R}^+)$ be a weak factorisation system on \mathcal{C}^+ .*

- (i) *If $\mathcal{L} \subseteq \mathcal{L}^+$, then $\mathcal{R} \supseteq \mathcal{R}^+ \cap \text{mor } \mathcal{C}$.*
- (ii) *If $(\mathcal{L}, \mathcal{R})$ and $(\mathcal{L}^+, \mathcal{R}^+)$ are both cofibrantly generated by the same ensemble \mathcal{I} , then $\mathcal{R} = \mathcal{R}^+ \cap \text{mor } \mathcal{C}$.*

Dually:

- (i') *If $\mathcal{R} \subseteq \mathcal{R}^+$, then $\mathcal{L} \supseteq \mathcal{L}^+ \cap \text{mor } \mathcal{C}$.*
- (ii') *If $(\mathcal{L}, \mathcal{R})$ and $(\mathcal{L}^+, \mathcal{R}^+)$ are both fibrantly generated by the same ensemble \mathcal{F} , then $\mathcal{L} = \mathcal{L}^+ \cap \text{mor } \mathcal{C}$.*

Moreover:

- (iii) *If $\mathcal{L} \subseteq \mathcal{L}^+$ and $\mathcal{R} \subseteq \mathcal{R}^+$, then $(\mathcal{L}^+, \mathcal{R}^+)$ is an extension of $(\mathcal{L}, \mathcal{R})$.*

Proof. Since \mathcal{C} is a full subcategory of \mathcal{C}^+ , if $g : Z \rightarrow W$ and $f : X \rightarrow Y$ are morphisms in \mathcal{C} , then any lifting problem of the following form in \mathcal{C}^+ is already in \mathcal{C} ,

$$\begin{array}{ccc} Z & \longrightarrow & X \\ g \downarrow & & \downarrow f \\ W & \longrightarrow & Y \end{array}$$

and moreover any solution to the above lifting problem in \mathcal{C}^+ is also a solution in \mathcal{C} . Thus, $g \square f$ in \mathcal{C} if and only if $g \square f$ in \mathcal{C}^+ .

(i). Suppose f is in $\mathcal{R}^+ \cap \text{mor } \mathcal{C}$. Then f has the right lifting property in \mathcal{C}^+ with respect to every morphism in \mathcal{L}^+ , and in particular, f has the right lifting property in \mathcal{C} with respect to every morphism in \mathcal{L} ; hence f is in \mathcal{R} , and therefore $\mathcal{R} \supseteq \mathcal{R}^+ \cap \text{mor } \mathcal{C}$.

(ii). A morphism is in \mathcal{R} (resp. \mathcal{R}^+) if and only if it has the right lifting property in \mathcal{C} (resp. \mathcal{C}^+) with respect to every morphism in \mathcal{I} , so by our initial observation, we must have $\mathcal{R} = \mathcal{R}^+ \cap \text{mor } \mathcal{C}$.

(iii). Immediately follows from claims (i) and (i'). ■

Proposition A.3.19. *Let $(\mathcal{L}, \mathcal{R})$ be a weak (resp. orthogonal) factorisation system for a category \mathcal{C} , and let $(\mathcal{L}', \mathcal{R}')$ be a weak (resp. orthogonal) factorisation system for a category \mathcal{C}' . Given an adjunction*

$$F \dashv U : \mathcal{C}' \rightarrow \mathcal{C}$$

the following are equivalent:

- (i) *F sends morphisms in \mathcal{L} to morphisms in \mathcal{L}' .*
- (ii) *U sends morphisms in \mathcal{R}' to morphisms in \mathcal{R} .*

Proof. The adjunction induces a bijection between solutions to the two lifting problems shown below:

$$\begin{array}{ccc} FZ & \longrightarrow & X \\ Fg \downarrow & \nearrow ? & \downarrow f \\ FW & \longrightarrow & Y \end{array} \qquad \begin{array}{ccc} Z & \longrightarrow & UX \\ g \downarrow & \nearrow ? & \downarrow Uf \\ W & \longrightarrow & UY \end{array}$$

Thus, $Fg \boxtimes f$ (resp. $Fg \perp f$) if and only if $g \boxtimes Uf$ (resp. $g \perp Uf$). ■

¶ A.3.20. Let $\mathbb{2}$ be the category $\{0 \rightarrow 1\}$, and let $\mathbb{3}$ be $\{0 \rightarrow 1 \rightarrow 2\}$. Thus, given a category \mathcal{C} , the functor category $[\mathbb{2}, \mathcal{C}]$ is the category of arrows and commutative squares in \mathcal{C} . There are three embeddings $d^0, d^1, d^2 : \mathbb{2} \rightarrow \mathbb{3}$:

$$\begin{array}{lll} d^0(0) = 1 & d^1(0) = 0 & d^2(0) = 0 \\ d^0(1) = 2 & d^1(1) = 2 & d^2(1) = 1 \end{array}$$

These then induce (by precomposition) three functors $d_0, d_1, d_2 : [\mathbb{3}, \mathcal{C}] \rightarrow [\mathbb{2}, \mathcal{C}]$.

Definition A.3.21. A **functorial factorisation system** on a category \mathcal{C} is a pair of functors $L, R : [\mathbb{2}, \mathcal{C}] \rightarrow [\mathbb{2}, \mathcal{C}]$ for which there exists a (necessarily unique) functor $F : [\mathbb{2}, \mathcal{C}] \rightarrow [\mathbb{3}, \mathcal{C}]$ satisfying the following equations:

$$d_2 F = L \qquad d_1 F = \text{id}_{[\mathbb{2}, \mathcal{C}]} \qquad d_0 F = R$$

A **functorial weak** (resp. **orthogonal**) **factorisation system** on \mathcal{C} is a weak (resp. orthogonal) factorisation system $(\mathcal{L}, \mathcal{R})$ together with a functorial factorisation system (L, R) such that $Lf \in \mathcal{L}$ and $Rf \in \mathcal{R}$ for all morphisms f in \mathcal{C} .

Lemma A.3.22. *Let A be an object in a category C and let $\Sigma_A : C_{/A} \rightarrow C$ be the projection from the slice category.*

- (i) *For each functorial factorisation system (L, R) on C , there exists a unique functorial factorisation system (L_A, R_A) on $C_{/A}$ such that*

$$[2, \Sigma_A] \circ L_A = L \circ [2, \Sigma_A] \quad [2, \Sigma_A] \circ R_A = R \circ [2, \Sigma_A]$$

where $[2, \Sigma_A] : [2, C_{/A}] \rightarrow [2, C]$ is the evident induced functor.

- (ii) *If (L, R) is part of a functorial weak or orthogonal factorisation system on C , then (L_A, R_A) is compatible with the induced weak or orthogonal factorisation system on $C_{/A}$ as well.*

Proof. Obvious. ◆

Proposition A.3.23. *Any orthogonal factorisation system can be extended to a functorial one.*

Proof. For each morphism f in a category C with an orthogonal factorisation system $(\mathcal{L}, \mathcal{R})$, choose a factorisation $f = Rf \circ Lf$ with $Lf \in \mathcal{L}$ and $Rf \in \mathcal{R}$. Given a commutative square in C , say

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

the lifting property ensures that the dashed arrow in the diagram below exists,

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ Lg \downarrow & & \downarrow Lf \\ \bullet & \dashrightarrow & \bullet \\ Rg \downarrow & & \downarrow Rf \\ W & \xrightarrow{w} & Y \end{array}$$

and orthogonality ensures uniqueness and hence functoriality. ■

Corollary A.3.24. *If $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on a category \mathcal{C} , then, for any category \mathcal{J} , there exists an orthogonal factorisation system on the functor category $[\mathcal{J}, \mathcal{C}]$ where a natural transformation is in the left (resp. right) class if and only if all its components are in \mathcal{L} (resp. \mathcal{R}).*

Proof. Obviously, every morphism in $[\mathcal{J}, \mathcal{C}]$ admits such a factorisation, since $(\mathcal{L}, \mathcal{R})$ -factorisations in \mathcal{C} are functorial. By considering a commutative diagram in \mathcal{C} of the form below,

$$\begin{array}{ccccc}
Z' & \longrightarrow & X' & & \\
\downarrow g' & \searrow & \downarrow & \searrow & \\
& & Z & \longrightarrow & X \\
& & \downarrow g & \downarrow f' & \downarrow f \\
W' & \longrightarrow & Y' & & \\
& \searrow & \downarrow & \searrow & \\
& & W & \longrightarrow & Y
\end{array}$$

where f and f' are in \mathcal{R} while g and g' are in \mathcal{L} , using the fact that $(\mathcal{E}, \mathcal{M})$ is an *orthogonal* factorisation system, one may show that lifting problems in $[\mathcal{J}, \mathcal{C}]$ admit unique solutions, and that these solutions are moreover constructed componentwise. Thus, $(\mathcal{L}, \mathcal{R})$ induces an orthogonal factorisation system on $[\mathcal{J}, \mathcal{C}]$. ■

The following characterisation of functorial orthogonal factorisation systems is due to Grandis and Tholen [2006]:

Theorem A.3.25. *Let (L, R) be a functorial factorisation system on a category \mathcal{C} . The following are equivalent:*

- (i) L is the underlying endofunctor of an idempotent comonad on $[\mathbb{2}, \mathcal{C}]$ with counit given by $\varepsilon_k = (\text{id}_{\text{dom } k}, Rk)$, and R is the underlying endofunctor of an idempotent monad on $[\mathbb{2}, \mathcal{C}]$ with unit given by $\eta_h = (h, \text{id}_{\text{codom } h})$.
- (ii) For all morphisms h in \mathcal{C} , RLh and LRh are isomorphisms in \mathcal{C} .
- (iii) For any two morphisms in \mathcal{C} , say h and k , we have $Lk \perp Rh$.
- (iv) $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on \mathcal{C} extending (L, R) , where:

$$\begin{aligned}
\mathcal{L} &= \{g \in \text{mor } \mathcal{C} \mid Rg \text{ is an isomorphism in } \mathcal{C}\} \\
\mathcal{R} &= \{f \in \text{mor } \mathcal{C} \mid Lf \text{ is an isomorphism in } \mathcal{C}\}
\end{aligned}$$

(v) *There exists an orthogonal factorisation system $(\mathcal{L}, \mathcal{R})$ extending (L, R) .*

Proof. (i) \Leftrightarrow (ii). This is a standard fact about idempotent (co)monads.

(ii) \Rightarrow (iii). Now, consider the following lifting problem:

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

Since (L, R) is a functorial factorisation system, we get a commutative diagram of the form below,

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ Lg \downarrow & & \downarrow Lf \\ W' & \xrightarrow{t} & X' \\ Rg \downarrow & & \downarrow Rf \\ W & \xrightarrow{w} & Y \end{array}$$

but Rg and Lf are isomorphisms, so $(Lf)^{-1} \circ t \circ (Rg)^{-1}$ is the required lift $W \rightarrow X$. On the other hand, if $s : W \rightarrow X$ is any morphism such that $f \circ s = w$ and $s \circ g = z$, then by taking (L, R) -factorisations of the vertical arrows in the following diagram,

$$\begin{array}{ccccccc} Z & \xrightarrow{g} & W & \xrightarrow{s} & X & \xrightarrow{\text{id}} & X \\ g \downarrow & & \text{id} \downarrow & & \downarrow \text{id} & & \downarrow f \\ W & \xrightarrow{\text{id}} & W & \xrightarrow{s} & X & \xrightarrow{f} & Y \end{array}$$

we find it must be the case that $Lf \circ s \circ Rg = t$, so we indeed have $g \perp f$.

(iii) \Rightarrow (iv). In particular, $g \perp Rg$ and $Lf \perp f$, so there must exist morphisms i and r making the diagrams below commute:

$$\begin{array}{ccc} Z & \xrightarrow{Lg} & W' \\ g \downarrow & \nearrow i & \downarrow Rg \\ W & \xrightarrow{\text{id}} & W \end{array} \quad \begin{array}{ccc} X & \xrightarrow{\text{id}} & X \\ Lf \downarrow & \nearrow r & \downarrow f \\ X' & \xrightarrow{Rf} & Y \end{array}$$

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We then obtain the following equations,

$$\begin{aligned} (i \circ Rg) \circ Lg &= Lg & (Lf \circ r) \circ Lf &= Lf \\ Rg \circ (i \circ Rg) &= Rg & Rf \circ (Lf \circ r) &= Rf \end{aligned}$$

and since $Lg \perp Rg$ and $Lf \perp Rf$, we must have $i \circ Rg = \text{id}_{W'}$, and $Lf \circ r = \text{id}_{X'}$. Thus, $g \in \mathcal{L}$ and $f \in \mathcal{R}$, and the same argument now shows that ${}^\perp\mathcal{R} \subseteq \mathcal{L}$ and $\mathcal{L}^\perp \subseteq \mathcal{R}$.

It remains to be shown that $\mathcal{L} \subseteq {}^\perp\mathcal{R}$ and $\mathcal{R} \subseteq \mathcal{L}^\perp$. First, suppose $g \in \mathcal{L}$ and $f \in \mathcal{R}$, and consider the following lifting problem:

$$\begin{array}{ccc} Z & \xrightarrow{z} & X \\ g \downarrow & & \downarrow f \\ W & \xrightarrow{w} & Y \end{array}$$

With r and i as in the previous paragraph, we obtain a commutative diagram of the form below,

$$\begin{array}{ccccccc} Z & \xrightarrow{\text{id}} & Z & \xrightarrow{z} & X & \xrightarrow{\text{id}} & X \\ \downarrow g & & \downarrow Lg & & \downarrow Lf & \nearrow r & \downarrow f \\ & & W' & \xrightarrow{t} & X' & & \\ & \nearrow i & \downarrow Rg & & \downarrow Rf & & \\ W & \xrightarrow{\text{id}} & W & \xrightarrow{w} & Y & \xrightarrow{\text{id}} & Y \end{array}$$

where the arrow t is obtained by the functoriality of (L, R) -factorisations. Thus, $r \circ t \circ i$ is the required lift $W \rightarrow X$, and it is unique, since Rg and Lf are isomorphisms. (Recall the proof of (ii) \Rightarrow (iii).) We conclude that $\mathcal{L} = {}^\perp\mathcal{R}$ and $\mathcal{R} = \mathcal{L}^\perp$.

(iv) \Rightarrow (v). Immediate.

(v) \Rightarrow (iii). If $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on \mathcal{C} such that $Lf \in \mathcal{L}$ and $Rf \in \mathcal{R}$ for all morphisms f in \mathcal{C} , then we must have $Lk \perp Rh$ for all h and k in $\text{mor } \mathcal{C}$, as required.

(iv) \Rightarrow (ii). Immediate. ■

REMARK A.3.26. It is clear that a functorial factorisation system is associated with *at most one* orthogonal factorisation system: indeed, if $(\mathcal{L}', \mathcal{R}')$ is any orthogonal factorisation system extending a functorial factorisation system (L, R) , and $(\mathcal{L}, \mathcal{R})$ is the induced orthogonal factorisation system as in the theorem, then each morphism in \mathcal{L} (resp. \mathcal{R}) is a retract of some morphism in \mathcal{L}' (resp. \mathcal{R}'); but by [proposition A.3.12](#), this implies $\mathcal{L} \subseteq \mathcal{L}'$ and $\mathcal{R} \subseteq \mathcal{R}'$, and applying [proposition A.3.3](#), we also get $\mathcal{L} \supseteq \mathcal{L}'$ and $\mathcal{R} \supseteq \mathcal{R}'$.

Corollary A.3.27. *If $(\mathcal{L}, \mathcal{R})$ is an orthogonal factorisation system on a category \mathcal{C} , then:*

- (i) \mathcal{L} , considered as a full subcategory of $[2, \mathcal{C}]$, is replete and coreflective.
- (ii) \mathcal{L} is closed under all colimits in $[2, \mathcal{C}]$.
- (iii) If a diagram in \mathcal{L} has a limit in $[2, \mathcal{C}]$, then it also has a limit in \mathcal{L} .

Dually:

- (i') \mathcal{R} , considered as a full subcategory of $[2, \mathcal{C}]$, is replete and reflective.
- (ii') \mathcal{R} is closed under all limits in $[2, \mathcal{C}]$.
- (iii') If a diagram in \mathcal{R} has a colimit in $[2, \mathcal{C}]$, then it also has a colimit in \mathcal{R} .

Proof. Using [proposition A.3.23](#) and [theorem A.3.25](#), the above claims amount to standard facts about the Eilenberg–Moore category for idempotent (co)monads. ■

There is a similar characterisation of functorial weak factorisation systems:

Theorem A.3.28. *Let (L, R) be a functorial factorisation system on a category \mathcal{C} . The following are equivalent:*

- (i) For any two morphisms in \mathcal{C} , say h and k , $Lk \sqsupseteq Rh$.
- (ii) $(\mathcal{L}, \mathcal{R})$ is an weak factorisation system on \mathcal{C} extending (L, R) , where:

$$\begin{aligned} \mathcal{L} &= \{g \in \text{mor } \mathcal{C} \mid \exists i \in \text{mor } \mathcal{C}. i \circ g = Lg \wedge Rg \circ i = \text{id}_{\text{codom } g}\} \\ \mathcal{R} &= \{f \in \text{mor } \mathcal{C} \mid \exists r \in \text{mor } \mathcal{C}. f \circ r = Rf \wedge r \circ Lf = \text{id}_{\text{dom } f}\} \end{aligned}$$

- (iii) There exists a weak factorisation system $(\mathcal{L}, \mathcal{R})$ extending (L, R) .

Proof. The proof is essentially the same as that of [theorem A.3.25](#). ■

REMARK A.3.29. As with orthogonal factorisation systems, there is *at most one* weak factorisation system extending any functorial factorisation system.

The two theorems above motivate the following definition:

Definition A.3.30. A **algebraic factorisation system**^[4] on a category \mathcal{C} is a pair (\mathbf{L}, \mathbf{R}) satisfying the following conditions:

- $\mathbf{L} = (L, \varepsilon, \delta)$ is a comonad on $[2, \mathcal{C}]$, where $\varepsilon_k = (\text{id}_{\text{dom } k}, Rk)$.
- $\mathbf{R} = (R, \eta, \mu)$ is a monad on $[2, \mathcal{C}]$, where $\eta_h = (Lh, \text{id}_{\text{codom } h})$.
- (L, R) constitute a functorial factorisation system on \mathcal{C} .

Corollary A.3.31. *Any functorial orthogonal factorisation system extends to an algebraic factorisation system in a unique way; conversely, an algebraic factorisation system induces an orthogonal factorisation system if and only if the underlying comonad and monad are both idempotent.*

Proof. This follows from the definition above and [theorem A.3.25](#). ■

Proposition A.3.32. *Let (\mathbf{L}, \mathbf{R}) be an algebraic factorisation system on a category \mathcal{C} .*

- (i) *Let $f : X \rightarrow Y$ and $g : Z \rightarrow W$ be objects in $[2, \mathcal{C}]$. If $\alpha : Rf \rightarrow f$ is a \mathbf{R} -algebra structure and $\beta : g \rightarrow Lg$ is a \mathbf{L} -coalgebra structure, then $\alpha_1 : Y \rightarrow Y$ and $\beta_0 : Z \rightarrow Z$ are identity morphisms, and we have the following identities:*

$$\begin{array}{ll} \alpha_0 \circ Lf = \text{id}_X & Rg \circ \beta_1 = \text{id}_W \\ f \circ \alpha_0 = Rf & \beta_1 \circ g = Lg \end{array}$$

- (ii) *If f admits a \mathbf{L} -coalgebra structure and g admits an \mathbf{R} -algebra structure, then $f \boxtimes g$.*
- (iii) *There exists a (unique) weak factorisation system $(\mathcal{L}, \mathcal{R})$ on \mathcal{C} such that $Lk \in \mathcal{L}$ and $Rh \in \mathcal{R}$ for all h and k in $\text{mor } \mathcal{C}$.*

[4] — or **natural weak factorisation system** in the sense of Grandis and Tholen [2006] and Garner [2009].

Proof. (i). The claim follows from the **L**-coalgebra counitality axiom and the **R**-algebra unitality axiom:

$$\alpha \circ \eta_f = \text{id}_f \qquad \varepsilon_g \circ \beta = \text{id}_g$$

(ii). It then follows that the diagram below commutes,

$$\begin{array}{ccccccc}
Z & \xrightarrow{\text{id}} & Z & \xrightarrow{z} & X & \xrightarrow{\text{id}} & X \\
\downarrow g & & \downarrow Lg & & \downarrow Lf & \nearrow \alpha_0 & \downarrow f \\
& & W' & \xrightarrow{t} & X' & & \\
& \nearrow \beta_1 & \downarrow Rg & & \downarrow Rf & & \\
W & \xrightarrow{\text{id}} & W & \xrightarrow{w} & Y & \xrightarrow{\text{id}} & Y
\end{array}$$

where the arrow t is obtained by the functoriality of (L, R) -factorisations; clearly, $\alpha_0 \circ t \circ \beta_1$ is the required lift.

(iii). Finally, for any two morphisms in \mathcal{C} , say h and k , we simply note that $\delta_k : Lk \rightarrow LLk$ is an **L**-coalgebra structure and $\mu_h : RRh \rightarrow Rh$ is an **R**-algebra structure, so we may apply [theorem A.3.28](#) to obtain the conclusion. \blacksquare

A.4 Relative categories

Prerequisites. § 0.1.

In this section we use the explicit universe convention.

Definition A.4.1. A **relative category** \mathcal{C} consists of a category $\text{und } \mathcal{C}$ and a subcategory $\text{weq } \mathcal{C}$ such that $\text{ob und } \mathcal{C} = \text{ob weq } \mathcal{C}$. We say $\text{und } \mathcal{C}$ is the **underlying category** of \mathcal{C} , and that the morphisms in $\text{weq } \mathcal{C}$ are the **weak equivalences** in \mathcal{C} . A **relative subcategory** of a relative category \mathcal{C} is a relative category \mathcal{C}' such that $\text{und } \mathcal{C}'$ is a subcategory of $\text{und } \mathcal{C}$, and we further demand that $\text{weq } \mathcal{C}' = \text{weq } \mathcal{C} \cap \text{und } \mathcal{C}'$.

REMARK A.4.2. The subcategory $\text{weq } \mathcal{C}$ is entirely determined by $\text{mor weq } \mathcal{C}$, so a relative category may equivalently be defined as a category equipped with a distinguished subset of morphisms closed under composition and containing all the identity morphisms.

For brevity, we will write $\text{ob } \mathcal{C}$ for $\text{ob und } \mathcal{C}$, $\text{mor } \mathcal{C}$ for $\text{ob und } \mathcal{C}$, and we may occasionally abuse notation and write $\text{weq } \mathcal{C}$ instead of $\text{mor weq } \mathcal{C}$.

REMARK A.4.3. Every category C can be endowed with the structure of a relative category in two ways: we can make it into a **minimal relative category** $\min C$ by taking $\text{weq } \min C$ to be the set of identity morphisms in C ; or we could make it into a **maximal relative category** $\max C$ by taking $\text{weq } \max C = \text{mor } C$. We may also define the **minimal saturated relative category** $\min^+ C$ by taking $\text{weq } \min^+ C$ to be the set of all isomorphisms in C .

Definition A.4.4. Given a relative category C , the **opposite relative category** C^{op} is defined by $\text{und } C^{\text{op}} = (\text{und } C)^{\text{op}}$ and $\text{weq } C^{\text{op}} = (\text{weq } C)^{\text{op}}$.

Definition A.4.5. Let C and D be relative categories. A **relative functor** $C \rightarrow D$ is a functor $\text{und } C \rightarrow \text{und } D$ that sends weak equivalences in C to weak equivalences in D . The **relative functor category** $[C, D]_{\text{h}}$ is the full subcategory of $[\text{und } C, \text{und } D]$ spanned by the relative functors, and the weak equivalences in $[C, D]_{\text{h}}$ are defined to be the natural transformations that are componentwise weak equivalences in D .

Definition A.4.6. Let C be a category and let $\mathcal{W} \subseteq \text{mor } C$. A **localisation of C away from \mathcal{W}** is a category $C[\mathcal{W}^{-1}]$ equipped with a functor $\gamma : C \rightarrow C[\mathcal{W}^{-1}]$ with the following universal property:

- Given a functor $F : C \rightarrow D$ such that Ff is an isomorphism for all f in \mathcal{W} , there exists a unique functor $\bar{F} : C[\mathcal{W}^{-1}] \rightarrow D$ such that $\bar{F}\gamma = F$.

The functor $\gamma : C \rightarrow \text{Ho } C$ is called the **localising functor**.

REMARK A.4.7. The universal property in the above definition is strict; as such, $C[\mathcal{W}^{-1}]$ is unique up to unique isomorphism. Nonetheless, $C[\mathcal{W}^{-1}]$ automatically has a 2-universal property: if $F, G : C \rightarrow D$ both factor through $C[\mathcal{W}^{-1}]$, then so do all natural transformations $F \Rightarrow G$.

Proposition A.4.8. *If C is a \mathbf{U} -small category, then there exists a \mathbf{U} -small category with the universal property of $C[\mathcal{W}^{-1}]$.*

Proof. Use the general adjoint functor theorem. □

Definition A.4.9. The **homotopy category** of a relative category C is a localisation of $\text{und } C$ away from $\text{weq } C$ and is denoted $\text{Ho } C$. A **semi-saturated relative category** is a relative category in which every isomorphism is a weak equivalence. A **saturated relative category** is a relative category C such that the weak equivalences in C are precisely the ones that become isomorphisms in $\text{Ho } C$.

REMARK A.4.10. Obviously, there is no loss of generality in considering semi-saturated relative categories and their homotopy categories instead of localisations $C[\mathcal{W}^{-1}]$ for arbitrary subsets $\mathcal{W} \subseteq \text{mor } C$.

REMARK A.4.11. Clearly, every saturated relative category is semi-saturated, and a minimal saturated relative category is indeed saturated in the sense above.

Definition A.4.12. Let C be a category and let \mathcal{W} be a subset of $\text{mor } C$. The **2-out-of-3 property** for \mathcal{W} says:

- Given any two morphisms $f : X \rightarrow Y$, $g : Y \rightarrow Z$ in C , if any two of f , g , or $g \circ f$ are in \mathcal{W} , then all of them are.

The **2-out-of-6 property** for \mathcal{W} says:

- Given any three morphisms $f : X \rightarrow Y$, $g : Y \rightarrow Z$, $h : Y \rightarrow Z$ in C , if both $h \circ g$ and $g \circ f$ are in \mathcal{W} , then so too are f , g , h , and $h \circ g \circ f$.

Lemma A.4.13. Let C be a category and let $\mathcal{W} \subseteq \text{mor } C$.

- (i) If \mathcal{W} has the 2-out-of-6 property, then it also has the 2-out-of-3 property.
- (ii) The set of all isomorphisms in C has the 2-out-of-6 property.
- (iii) If $F : C' \rightarrow C$ is a functor and \mathcal{W} has either the 2-out-of-3 property or the 2-out-of-6 property, then $F^{-1}\mathcal{W}$ has the same property.

Proof. (i). Consider the three cases $f = \text{id}$, $g = \text{id}$, $h = \text{id}$ in turn.

(ii). If $h \circ g$ and $g \circ f$ are isomorphisms, then g must be split epic and split monic; thus g itself is an isomorphism, hence so too are f and h .

(iii). Obvious. ■

Corollary A.4.14. If C is a saturated relative category, then $\text{weq } C$ has the 2-out-of-6 property. ■

Proposition A.4.15. Let \mathbf{RelCat} be the category of \mathbf{U} -small relative categories and relative functors, let $\mathbf{SsRelCat}$ be the full subcategory of semi-saturated relative categories, and let \mathbf{Cat} be the category of \mathbf{U} -small categories and functors.

- (i) \mathbf{RelCat} is a cartesian closed category, where the product of C and D is the cartesian product $C \times D$ with weak equivalences taken componentwise, and the exponential of \mathcal{E} by D is the relative functor category $[D, \mathcal{E}]_{\mathbf{h}}$.

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- (ii) **RelCat** is a locally finitely presentable **U**-category,^[5] and the two functors $\text{und}, \text{weq} : \mathbf{RelCat} \rightarrow \mathbf{Cat}$ are \aleph_0 -accessible^[6] and jointly conservative.
- (iii) **SsRelCat** is a locally finitely presentable **U**-category, and the inclusion $\mathbf{SsRelCat} \hookrightarrow \mathbf{RelCat}$ is \aleph_0 -accessible and has a left adjoint.
- (iv) **SsRelCat** is an exponential ideal in **RelCat**.
- (v) The full subcategory spanned by the minimal relative categories is an exponential ideal in **RelCat**.
- (vi) The full subcategory spanned by the minimal saturated relative categories is an exponential ideal in **SsRelCat**.

Proof. (i). This is straightforward from the definitions.

(ii). Obviously, a relative functor $F : C \rightarrow D$ such that $\text{und } F : \text{und } C \rightarrow \text{und } D$ and $\text{weq } F : \text{weq } C \rightarrow \text{weq } D$ are both isomorphisms is itself an isomorphism, so $\text{und}, \text{weq} : \mathbf{RelCat} \rightarrow \mathbf{Cat}$ are indeed jointly conservative.

It is also not hard to check that limits for all **U**-small diagrams and colimits for **U**-small filtered diagrams in **RelCat** exist and can be computed componentwise in **Cat**, so (by [theorem 0.2.26](#)) it is enough to show that **RelCat** is a \aleph_0 -accessible **U**-category. Clearly, a relative category C such that $\text{und } C$ is finitely presentable in **Cat** and $\text{weq } C$ is a finitely-generated subcategory of $\text{und } C$ is itself finitely presentable in **RelCat**, so **RelCat** is indeed \aleph_0 -accessible.

(Alternatively, one may appeal to the sketchability theorem^[7] and the fact that a relative category is manifestly a model for a certain finite-limit sketch.)

(iii). It is clear that **SsRelCat** is closed in **RelCat** under limits for all **U**-small diagrams and colimits for all **U**-small filtered diagrams, and we know that **RelCat** is a locally finitely presentable category, so (by [proposition 0.2.21](#)) it is enough to construct a left adjoint for the inclusion $\mathbf{SsRelCat} \hookrightarrow \mathbf{RelCat}$. This may be done using the general adjoint functor theorem.

(iv) – (vi). All straightforward. ■

[5] See [definition 0.2.22](#).

[6] See [definition 0.2.18](#).

[7] See Proposition 1.51 in [LPAC] or Proposition 5.6.4 in [Borceux, 1994b].

Proposition A.4.16. *Let \mathbf{RelCat} be the category of \mathbf{U} -small relative categories and relative functors, let $\mathbf{SsRelCat}$ be the full subcategory of semi-saturated relative categories and relative functors, and let \mathbf{Cat} be the category of \mathbf{U} -small categories and functors. We have the following strings of adjoint functors:*

$$\begin{aligned} \min \dashv \text{und} \dashv \max \dashv \text{weq} &: \mathbf{RelCat} \rightarrow \mathbf{Cat} \\ \text{Ho} \dashv \min^+ \dashv \text{und} \dashv \max \dashv \text{weq} &: \mathbf{SsRelCat} \rightarrow \mathbf{Cat} \end{aligned}$$

The functors \min , \min^+ , and \max are moreover fully faithful, and Ho preserves finite products.

Proof. All but the last of the above claims are obvious; for the preservation of finite products under Ho , we refer to [proposition A.2.13](#). ■

Definition A.4.17. A **zigzag type** is a relative category T where $\text{und} T$ is the free category on an inhabited finite planar graph of the form



where the edges are arrows that point either left or right, and $\text{weq} T$ consists of all identities and all composites of left-pointing arrows. A **morphism of zigzag types** is a relative functor that maps the leftmost object to the leftmost object and the rightmost object to the rightmost object. We write \mathbf{T} for the category of zigzag types.^[8]

A **zigzag** of type T in a relative category C is a relative functor $T \rightarrow C$. Given objects X and Y in C , we denote by $C^T(X, Y)$ the category whose objects are the zigzags starting at X and ending at Y and whose morphisms are commutative diagrams in C of the form

$$\begin{array}{ccccccc} X & \text{---} & \bullet & \text{---} & \dots & \text{---} & \bullet & \text{---} & Y \\ \parallel & & \downarrow & & & & \downarrow & & \parallel \\ X & \text{---} & \bullet & \text{---} & \dots & \text{---} & \bullet & \text{---} & Y \end{array}$$

where the rows are zigzags of type T and the unmarked vertical arrows are weak equivalences.

[8] Warning: This is the *opposite* of the category \mathbf{T} defined in [DHKS, § 34].

Example A.4.18. If $f : X \rightarrow Y$ is a weak equivalence in a relative category \mathcal{C} , then we have commutative diagrams

$$\begin{array}{ccc}
 X & \xlongequal{\quad} & X & \xlongequal{\quad} & X \\
 \parallel & & \downarrow f & & \parallel \\
 X & \xrightarrow{f} & Y & \xleftarrow{f} & X
 \end{array}
 \qquad
 \begin{array}{ccc}
 Y & \xleftarrow{f} & X & \xrightarrow{f} & Y \\
 \parallel & & \downarrow f & & \parallel \\
 Y & \xlongequal{\quad} & Y & \xlongequal{\quad} & Y
 \end{array}$$

and these correspond to morphisms of zigzags in \mathcal{C} .

REMARK A.4.19. It is clear that $\mathcal{C}^T(X, Y)$ is a subcategory of the relative functor category $[T, \mathcal{C}]_{\text{h}}$. Thus, if \mathcal{C} is a \mathbf{U} -small relative category, precomposition makes the assignment $T \mapsto \mathcal{C}^T(X, Y)$ into a functor $\mathbf{T}^{\text{op}} \rightarrow \mathbf{Cat}$, which we denote by $\mathcal{C}^*(X, Y)$. A Grothendieck construction applied to this functor yields the following \mathbf{U} -small category $\mathcal{C}^{(\mathbf{T})}(X, Y)$:

- Its objects are pairs (T, f) , where T is a zigzag type and f is a zigzag of type T in \mathcal{C} .
- A morphism $(T', f') \rightarrow (T, f)$ is a pair (α, β) where $\alpha : T \rightarrow T'$ is a morphism in \mathbf{T} and $\beta : \alpha^* f' \rightarrow f$ is a morphism in $\mathcal{C}^T(X, Y)$.
- The composite of a pair of morphisms $(\alpha', \beta') : (T'', f'') \rightarrow (T', f')$ and $(\alpha, \beta) : (T', f') \rightarrow (T, f)$ is given by $(\alpha' \circ \alpha, \beta \circ \alpha^* \beta')$.

There is an evident projection functor $\mathcal{C}^{(\mathbf{T})}(X, Y) \rightarrow \mathbf{T}^{\text{op}}$, and by construction it is a Grothendieck opfibration with a canonical splitting.

Lemma A.4.20. *Given a commutative diagram of the form below in a relative category \mathcal{C} ,*

$$\begin{array}{ccc}
 X & \xrightarrow{f} & Y \\
 a \downarrow & & \downarrow b \\
 X' & \xrightarrow{f'} & Y'
 \end{array}$$

if a and b are weak equivalences in \mathcal{C} , then we obtain the following morphisms of zigzags:

$$\begin{array}{ccccccc} X & \xlongequal{\quad} & X & \xrightarrow{f} & Y & \xrightarrow{b} & Y' \\ \parallel & & \downarrow a & & \downarrow b & & \parallel \\ X' & \xrightarrow{a} & X' & \xrightarrow{f'} & Y' & \xlongequal{\quad} & Y' \end{array}$$

$$\begin{array}{ccccccc} X' & \xleftarrow{a} & X & \xrightarrow{f} & Y & \xlongequal{\quad} & Y \\ \parallel & & \downarrow a & & \downarrow b & & \parallel \\ X' & \xlongequal{\quad} & X' & \xrightarrow{f'} & Y' & \xleftarrow{b} & Y \end{array}$$

In particular, $X \xrightarrow{f} Y \xrightarrow{b} Y'$ and $X \xrightarrow{a} X' \xrightarrow{f'} Y'$ are in the same connected component of $\mathcal{C}^{(\mathbf{T})}(X, Y')$; and $X' \xleftarrow{a} X \xrightarrow{f} Y$ and $X' \xrightarrow{f'} Y' \xleftarrow{b} Y$ are in the same connected component of $\mathcal{C}^{(\mathbf{T})}(X', Y)$. ■

Theorem A.4.21. *Let X and Y be objects in a relative category \mathcal{C} .*

- (i) *For each zigzag type T , the map that sends an object in $\mathcal{C}^T(X, Y)$ to the corresponding composite in $\text{Ho } \mathcal{C}(X, Y)$ is a functor when the latter is regarded as a discrete category.*
- (ii) *The functors described above constitute a jointly surjective cocone from the diagram $\mathcal{C}^*(X, Y)$ to $\text{Ho } \mathcal{C}(X, Y)$.*
- (iii) *The induced functor $\mathcal{C}^{(\mathbf{T})}(X, Y) \rightarrow \text{Ho } \mathcal{C}(X, Y)$ is surjective, and moreover two objects in $\mathcal{C}^{(\mathbf{T})}(X, Y)$ become equal in $\text{Ho } \mathcal{C}$ if and only if they are in the same connected component.*

Proof. All obvious except for the last part of claim (iii), for which we refer to paragraphs 33.8 and 33.10 in [DHKS]. ■

A.5 Kan extensions

Prerequisites. §§ 0.I, A.I.

In this section we use the explicit universe convention.

Definition A.5.1. Let $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{C} \rightarrow \mathcal{E}$ be two functors. A **left Kan extension** (resp. **right Kan extension**) of G along F is an initial (resp. terminal) object of the category $(G \downarrow F^*)$ (resp. $(F^* \downarrow G)$) described below:

- The objects are pairs (H, α) where H is a functor $\mathcal{D} \rightarrow \mathcal{E}$ and α is a natural transformation of type $G \Rightarrow HF$ (resp. $HF \Rightarrow G$).
- The morphisms $(H', \alpha') \rightarrow (H, \alpha)$ are those natural transformations $\beta : H' \Rightarrow H$ such that $\beta F \bullet \alpha' = \alpha$ (resp. $\alpha \bullet \beta F = \alpha'$).

REMARK A.5.2. Clearly, Kan extensions are unique up to unique isomorphism if they exist. We write $(\text{Lan}_F G, \eta)$ for the left Kan extension of G along F and say η is the **unit** of $\text{Lan}_F G$; dually, we write $(\text{Ran}_F G, \varepsilon)$ for the right Kan extension of G along F and say ε is the **counit** of $\text{Ran}_F G$.

Lemma A.5.3. Let \mathbf{U} be a pre-universe and let \mathbf{Set} be the category of \mathbf{U} -sets. Let \mathcal{B} be a \mathbf{U} -small category and let \mathcal{C} be a locally \mathbf{U} -small category. Given functors $F : \mathcal{B} \rightarrow \mathcal{C}$ and $G : \mathcal{B} \rightarrow \mathbf{Set}$, if $H : \mathcal{C} \rightarrow \mathbf{Set}$ is the functor defined by the formula below,

$$H(C) = [\mathcal{B}, \mathbf{Set}](\mathcal{C}(C, F-), G-)$$

and $\varepsilon_B : H(FB) \rightarrow G(B)$ is defined by evaluation at id_{FB} , then (H, ε) is the right Kan extension of G along F .

Proof. Note that $H(C)$ so defined is indeed a \mathbf{U} -set, because \mathcal{B} is \mathbf{U} -small and \mathcal{C} is locally \mathbf{U} -small. The claim amounts to saying that (H, ε) is a terminal object in the comma category $(F^* \downarrow G)$, so that is what we must show.

Let $\varphi : (X, \alpha) \rightarrow (H, \varepsilon)$ be a morphism in $(F^* \downarrow G)$, i.e. a natural transformation $\varphi : X \Rightarrow H$ such that $\varepsilon \bullet \varphi F = \alpha$. Let C be an object in \mathcal{C} , let x be an element of $X(C)$, and consider the element $\varphi_C(x)$ of $H(C)$. By definition, this is a natural transformation $\mathcal{C}(C, F) \Rightarrow G$, so we may consider its component at an object B in \mathcal{B} , which will be a map $\mathcal{C}(C, FB) \rightarrow G(B)$. Let $f : C \rightarrow FB$ be an arrow in \mathcal{C} . By hypothesis,

$$\alpha_C(x) = \varepsilon_C(\varphi_C(x)_B \circ \mathcal{C}(f, FB)) = \varphi_C(x)_B(f)$$

thus the action of φ is entirely determined by α . Conversely, given any object (X, α) in the comma category $(F^* \downarrow G)$, it is easily verified that the above equation defines a morphism $\varphi : (X, \alpha) \rightarrow (H, \varepsilon)$, so (H, ε) is indeed a terminal object in $(F^* \downarrow G)$. ■

Corollary A.5.4. *For any two functors $F : \mathcal{B} \rightarrow \mathcal{C}$ and $G : \mathcal{B} \rightarrow \mathbf{Set}$, if \mathcal{B} is \mathbf{U} -small and \mathcal{C} is locally \mathbf{U} -small, then the following are equivalent:*

- (i) $(\mathbf{Ran}_F G, \varepsilon)$ is a right Kan extension of G along F .
- (ii) The maps $(\mathbf{Ran}_F G)(C) \rightarrow [\mathcal{B}, \mathbf{Set}](C(C, F), G)$ defined by $x \mapsto \varepsilon \bullet \theta_x F$, where $\theta_x : C(C, -) \Rightarrow G$ is the unique natural transformation such that $(\theta_x)_C(\text{id}_C) = x$, are bijections that are natural in C . ■

Definition A.5.5. Let $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{C} \rightarrow \mathcal{E}$ be two functors.

- A functor $L : \mathcal{E} \rightarrow \mathcal{F}$ **preserves** left Kan extensions of G along F if, given any left Kan extension (H, α) of G along F , $(LH, L\alpha)$ is a left Kan extension of LG along F .
- A functor $R : \mathcal{E} \rightarrow \mathcal{F}$ **preserves** right Kan extensions of G along F if, given any right Kan extension (H, α) of G along F , $(RH, R\alpha)$ is a right Kan extension of LG along F .

If a Kan extension is preserved by *all* functors, then it is said to be **absolute**.

Definition A.5.6. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -small sets, let \mathcal{E} be a locally \mathbf{U} -small category, and let $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{C} \rightarrow \mathcal{E}$ be two functors.

- A **pointwise left Kan extension** of G along F is one that is preserved by all functors of the form $\mathcal{E}(-, E) : \mathcal{E} \rightarrow \mathbf{Set}^{\text{op}}$.
- A **pointwise right Kan extension** of G along F is one that is preserved by all functors of the form $\mathcal{E}(E, -) : \mathcal{E} \rightarrow \mathbf{Set}$.

Definition A.5.7. Let $F : \mathcal{B} \rightarrow \mathcal{C}$ be a functor and let C be an object in \mathcal{C} .

- The **tautological cocone** to C induced by F is the cocone $\theta : FP_C \Rightarrow \Delta C$, where $P_C : (F \downarrow C) \rightarrow \mathcal{B}$ is the projection functor sending an object (B, f) in the comma category $(F \downarrow C)$ to the object B in \mathcal{B} , and $\theta_{(B, f)} = f$.
- The **tautological cone** from C induced by F is the cone $\theta : \Delta C \Rightarrow FP^C$, where $P^C : (C \downarrow F) \rightarrow \mathcal{C}$ is the projection functor sending an object (B, f) in the comma category $(C \downarrow F)$ to the object B in \mathcal{B} , and $\theta_{(B, f)} = f$.

Lemma A.5.8. *Let \mathcal{A} be any category, let \mathcal{B} be a \mathbf{U} -small category, let \mathcal{C} be locally \mathbf{U} -small category, and let $U : \mathcal{A} \rightarrow \mathcal{C}$, $V : \mathcal{B} \rightarrow \mathcal{C}$, and $Y : \mathcal{B} \rightarrow \mathbf{Set}$ be functors. Consider the following diagram of functors and natural transformations,*

$$\begin{array}{ccc} (U \downarrow V) & \xrightarrow{Q} & \mathcal{B} \\ P \downarrow & \theta \nearrow & \downarrow V \\ \mathcal{A} & \xrightarrow{U} & \mathcal{C} \end{array}$$

where $(U \downarrow V)$ is the comma category, $P : (U \downarrow V) \rightarrow \mathcal{A}$ and $Q : (U \downarrow V) \rightarrow \mathcal{B}$ are the two projections, and $\theta : UP \Rightarrow VQ$ is the tautological natural transformation defined by $\theta_{(A,B,f)} = f$. If (Z, ε) is a right Kan extension of Y along V , then $(ZU, \varepsilon Q \cdot Z\theta)$ is a right Kan extension of YQ along P .

Proof. By lemma A.5.3, we may take $Z : \mathcal{C} \rightarrow \mathbf{Set}$ to be the functor defined by the formula below,

$$Z(C) = [\mathbb{B}, \mathbf{Set}](C(C, F-), Y-)$$

with $\varepsilon : V^*(Z) \Rightarrow Y$ being the natural transformation obtained by evaluating elements of $Z(VB)$ at id_{VB} .

Let $\varphi : (X, \alpha) \rightarrow (ZU, \varepsilon Q \cdot Z\theta)$ be a morphism in $(P^* \downarrow YQ)$, i.e. a natural transformation $\varphi : X \Rightarrow ZU$ such that $\varepsilon Q \cdot Z\theta \cdot \varphi P = \alpha$. Let A be an object in \mathcal{A} , let x be an element of $X(A)$, and consider the element $\varphi_A(x)$ of $Z(UA)$. By definition, this is a natural transformation $N^V(C) \Rightarrow Y$, so we may consider its component at an object B in \mathcal{B} , which will be a map $C(UA, VB) \rightarrow Y(B)$. Let $f : UA \rightarrow VB$ be an arrow in \mathcal{C} ; then (A, B, f) is an object in the comma category $(U \downarrow V)$, and $\theta_{(A,B,f)} = f$ by definition. By hypothesis,

$$\alpha_{(A,B,f)}(x) = \varepsilon_B(\varphi_A(x)_B \circ C(f, VB)) = \varphi_A(x)_B(f)$$

thus the action of φ is entirely determined by α . Conversely, given any object (X, α) in the comma category $(P^* \downarrow YQ)$, it is easily verified that the above equation defines a morphism $\varphi : (X, \alpha) \rightarrow (ZU, \varepsilon Q \cdot Z\theta)$, so $(ZU, \varepsilon Q \cdot Z\theta)$ is indeed a terminal object in $(P^* \downarrow YQ)$. \blacksquare

Corollary A.5.9. *Let \mathcal{B} be a \mathbf{U} -small category and let \mathcal{C} be a locally \mathbf{U} -small category. Given functors $F : \mathcal{B} \rightarrow \mathcal{C}$ and $G : \mathcal{B} \rightarrow \mathbf{Set}$, if (H, ε) is a right Kan extension of G along F , then, for each object C in \mathcal{C} , the image under H of the tautological cone from C induced by F is a limiting cone in \mathbf{Set} .*

Proof. In the lemma, take \mathcal{A} to be the terminal category $\mathbb{1}$, take $U : \mathbb{1} \rightarrow C$ to be the functor sending the unique object in $\mathbb{1}$ to C , and take $V = F$; then $(HU, \varepsilon_Q \bullet H\theta)$ is a right Kan extension of $GQ : (C \downarrow F) \rightarrow \mathbf{Set}$ along the unique functor $P : (C \downarrow F) \rightarrow \mathbb{1}$, but it is clear that a right Kan extension of GQ along P amounts to a limit for the diagram GQ in \mathbf{Set} . ■

It is convenient at this juncture to introduce a concept borrowed from enriched category theory. The notation below follows [Kelly, 2005, § 3.1].

Definition A.5.10. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let C be a locally \mathbf{U} -small category. Given functors $W : \mathcal{J} \rightarrow \mathbf{Set}$ and $A : \mathcal{J} \rightarrow C$, a **W -weighted limit** of A is an object $\{W, A\}^{\mathcal{J}}$ in C together with bijections

$$C(C, \{W, A\}^{\mathcal{J}}) \cong [\mathcal{J}, \mathbf{Set}](W, C(C, A))$$

that are natural in C . We may also write $\lim_{\leftarrow j:\mathcal{J}}^{W_j} A_j$ instead of $\{W, A\}^{\mathcal{J}}$, if we wish to use an explicit variable j .

Dually, given functors $W : \mathcal{J}^{\text{op}} \rightarrow \mathbf{Set}$ and $A : \mathcal{J} \rightarrow C$, a **W -weighted colimit** of A is an object $W \star_{\mathcal{J}} A$ in C together with bijections

$$C(W \star_{\mathcal{J}} A, C) \cong [\mathcal{J}^{\text{op}}, \mathbf{Set}](W, C(A, C))$$

that are natural in C . We may also write $\lim_{\rightarrow j:\mathcal{J}}^{W_j} A_j$ instead of $W \star_{\mathcal{J}} A$, if we wish to use an explicit variable j .

REMARK A.5.11. Clearly, weighted limits and colimits are unique up to unique isomorphism if they exist.

It is also not hard to spell out the above definition in elementary terms; for example, one notes that to give a natural transformation $W \Rightarrow C(C, A)$, one must give a morphism $\lambda_{j,x} : C \rightarrow A_j$ for each object j in \mathcal{J} and each element x of W_j , and these are required to make various diagrams commute. This is a **W -weighted cone** from C to A , and $\{W, A\}^{\mathcal{J}}$ is an object equipped with a universal W -weighted cone to A . Similarly, one may define the notion of a **W -weighted cocone** from A to C , and then $W \star_{\mathcal{J}} A$ is an object equipped with a universal W -weighted cocone from A . In particular, if $W_j = 1$ for all j , then W -weighted limits and colimits reduce to ordinary limits and colimits.

The above discussion also shows that the concept of a weighted limit or colimit (within a fixed category!) does not depend on \mathbf{U} in any essential way.

Lemma A.5.12. *Let \mathcal{J} be a \mathbf{U} -small category. Given functors $F, G : \mathcal{J} \rightarrow \mathbf{Set}$, the F -weighted limit of G exists in \mathbf{Set} , and we have bijections*

$$\{F, G\}^{\mathcal{J}} \cong [\mathcal{J}, \mathbf{Set}](F, G)$$

that are natural in F and G .

Proof. One simply has to check that this works. ◆

Proposition A.5.13. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let $F : \mathcal{C} \rightarrow \mathcal{D}$ be any functor where \mathcal{C} and \mathcal{D} are locally \mathbf{U} -small categories.*

- (i) *For each weight $W : \mathcal{J} \rightarrow \mathbf{Set}$ and each diagram $A : \mathcal{J} \rightarrow \mathcal{C}$, if the weighted limits $\{W, A\}^{\mathcal{J}}$ and $\{W, FA\}^{\mathcal{J}}$ both exist, then there is a canonical comparison morphism*

$$F\{W, A\}^{\mathcal{J}} \rightarrow \{W, FA\}^{\mathcal{J}}$$

corresponding to the natural maps

$$[\mathcal{J}, \mathbf{Set}](W, \mathcal{C}(C, A)) \rightarrow [\mathcal{J}, \mathbf{Set}](W, \mathcal{D}(FC, FA))$$

induced by the functor F .

- (ii) *For any object C in \mathcal{C} , the functor $\mathcal{C}(C, -) : \mathcal{C} \rightarrow \mathbf{Set}$ preserves all weighted limits.*
- (iii) *The functors $\mathcal{C}(C, -) : \mathcal{C} \rightarrow \mathbf{Set}$ jointly reflect weighted limits.*
- (iv) *If F has a left adjoint, then F preserves weighted limits.*

Dually:

- (i') *For each weight $W : \mathcal{J}^{\text{op}} \rightarrow \mathbf{Set}$ and each diagram $A : \mathcal{J} \rightarrow \mathcal{C}$, if the weighted colimits $W \star_{\mathcal{J}} A$ and $W \star_{\mathcal{J}} FA$ both exist, then there is a canonical comparison morphism*

$$W \star_{\mathcal{J}} FA \rightarrow F(W \star_{\mathcal{J}} A)$$

corresponding to the natural maps

$$[\mathcal{J}, \mathbf{Set}](W, \mathcal{C}(A, C)) \rightarrow [\mathcal{J}, \mathbf{Set}](W, \mathcal{D}(FA, FC))$$

induced by the functor F .

- (ii') For any object C in \mathcal{C} , the functor $\mathcal{C}(-, C) : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$ sends any weighted colimit in \mathcal{C} to the corresponding weighted limit in \mathbf{Set} .
- (iii') The functors $\mathcal{C}(-, C) : \mathcal{C} \rightarrow \mathbf{Set}^{\text{op}}$ jointly reflect weighted colimits.
- (iv') If F has a right adjoint, then F preserves weighted colimits.

Proof. All straightforward. ◆

Definition A.5.14. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{D} be a locally \mathbf{U} -small category. Given a functor $F : \mathcal{C} \rightarrow \mathcal{D}$, the F -nerve functor $N^F : \mathcal{D} \rightarrow [\mathcal{C}^{\text{op}}, \mathbf{Set}]$ is defined by

$$N^F(\mathcal{D})(C) = \mathcal{D}(FC, D)$$

i.e. $N^F = F^*h_{\bullet}$, where $h_{\bullet} : \mathcal{D} \rightarrow [\mathcal{D}^{\text{op}}, \mathbf{Set}]$ is the usual Yoneda embedding.

Theorem A.5.15. Let \mathcal{C} , \mathcal{D} and \mathcal{E} be locally \mathbf{U} -small categories. Given functors $F : \mathcal{C} \rightarrow \mathcal{D}$ and $G : \mathcal{C} \rightarrow \mathcal{E}$, the following are equivalent:

- (i) (H, α) is a pointwise right Kan extension of G along F .
- (ii) For each object d in \mathcal{D} , the weighted limit $\{N^{F^{\text{op}}}(d), G\}^{\mathcal{C}}$ exists in \mathcal{E} , and there are isomorphisms

$$Hd \cong \{N^{F^{\text{op}}}(d), G\}^{\mathcal{C}}$$

natural in d , with $\alpha_c : HFc \rightarrow Gc$ corresponding to the element id_{Fc} of $N^{F^{\text{op}}}(Fc)(c) = \mathcal{D}(Fc, Fc)$.

- (iii) (Assuming \mathcal{C} is \mathbf{U} -small.) For each object d in \mathcal{D} , if $P^d : (d \downarrow F) \rightarrow \mathcal{C}$ is the projection sending (c, f) in the comma category $(d \downarrow F)$ to c , and $\varphi : \Delta d \Rightarrow FP^d$ is the tautological cone in \mathcal{D} , then the cone $\alpha P^d \bullet H\varphi : \Delta Hd \Rightarrow GP^d$ is limiting; and for each $g : d \rightarrow d'$ in \mathcal{D} , the morphism $Hg : Hd \rightarrow Hd'$ is the one induced by the functor $(d' \downarrow F) \rightarrow (d \downarrow F)$ sending (c', f') to $(c', f' \circ g)$. In particular, $\alpha_c : HFc \rightarrow Gc$ must be (equal to) the component of the limiting cone $\Delta Fc \Rightarrow GP^d$ at the object (c, id_{Fc}) of $(Fc \downarrow F)$.

In particular, if \mathcal{C} is a \mathbf{U} -small category and \mathcal{E} is \mathbf{U} -complete, then the right Kan extension of G along F exists and is pointwise.

Dually, the following are equivalent:

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- (i') (H, α) is a pointwise left Kan extension of G along F .
- (ii') For each object d in \mathcal{D} , the weighted colimit $N^F(d) \star_c G$ exists in \mathcal{E} , and there are isomorphisms

$$Hd \cong N^F(d) \star_c G$$

natural in d , with $\alpha_c : Gc \rightarrow HFc$ corresponding to the element id_{Fc} of $N^F(Fc)(c) = \mathcal{D}(Fc, Fc)$.

- (iii') (Assuming \mathcal{C} is \mathbf{U} -small.) For each object d in \mathcal{D} , if $P_d : (F \downarrow d) \rightarrow \mathcal{C}$ is the projection sending (c, f) in the comma category $(F \downarrow d)$ to c , and $\varphi : FP_d \Rightarrow \Delta d$ is the tautological cocone in \mathcal{D} , then the cocone $H\varphi \bullet \alpha P_d : GP_d \Rightarrow \Delta Hd$ is colimiting; and for each $g : d \rightarrow d'$ in \mathcal{D} , the morphism $Hg : Hd \rightarrow Hd'$ is the one induced by the functor $(F \downarrow d) \rightarrow (F \downarrow d')$ sending (c, f) to $(c, g \circ f)$. In particular, $\alpha_c : Gc \rightarrow HFc$ must be (equal to) the component of the colimiting cocone $GP_d \Rightarrow \Delta Fc$ at the object (c, id_{Fc}) of $(F \downarrow Fc)$.

In particular, if \mathcal{C} is a \mathbf{U} -small category and \mathcal{E} is \mathbf{U} -cocomplete, then the left Kan extension of G along F exists and is pointwise.

Proof. (i) \Leftrightarrow (ii). This is just a matter of unwinding the definitions.

(i) \Leftrightarrow (iii). [Corollary A.5.9](#) implies that the construction in (iii) does indeed define a right Kan extension in the special case $\mathcal{E} = \mathbf{Set}$, so we deduce that statements (i) and (iii) are equivalent by applying the Yoneda lemma; see also [\[CWM, Ch. X, §§ 3 and 5\]](#). \square

REMARK A.5.16. It is possible to extract an elementary characterisation of pointwise Kan extensions from the results above, thereby showing that the property of being pointwise does not depend on the choice of universe \mathbf{U} .

Corollary A.5.17. *Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be a functor. If \mathcal{C} is \mathbf{U} -small and \mathcal{D} is locally \mathbf{U} -small, then the functor $F^* : [\mathcal{D}, \mathbf{Set}] \rightarrow [\mathcal{C}, \mathbf{Set}]$ has both a left adjoint Lan_F and a right adjoint Ran_F . \blacksquare*

Corollary A.5.18. *If (H, α) is a pointwise right Kan extension of $G : \mathcal{C} \rightarrow \mathcal{E}$ along $F : \mathcal{C} \rightarrow \mathcal{D}$, and $R : \mathcal{E} \rightarrow \mathcal{F}$ is a functor, then $(RH, R\alpha)$ is a pointwise right Kan extension of RG along F , provided either:*

- (i) R preserves all weighted limits, or
- (ii) R preserves limits for \mathbf{U} -small diagrams and C is \mathbf{U} -small.

If (H, α) is a pointwise left Kan extension of $G : C \rightarrow \mathcal{E}$ along $F : C \rightarrow D$, and $L : \mathcal{E} \rightarrow \mathcal{F}$ is a functor, then $(LH, L\alpha)$ is a pointwise left Kan extension of LG along F , provided either:

- (i') L preserves all weighted colimits, or
- (ii') L preserves colimits for \mathbf{U} -small diagrams and C is \mathbf{U} -small. ■

Corollary A.5.19. *If (H, α) is a pointwise right (resp. left) Kan extension of $G : C \rightarrow \mathcal{E}$ along a fully faithful functor $F : C \rightarrow D$, then $\alpha : HF \Rightarrow G$ (resp. $\alpha : G \Rightarrow HF$) is a natural isomorphism.*

Proof. If F is fully faithful, then the comma category $(Fc \downarrow F)$ (resp. $(F \downarrow Fc)$) has an initial (resp. terminal) object, namely (c, id_{Fc}) , so the component $\alpha_c : HFc \rightarrow Gc$ (resp. $\alpha_c : Gc \rightarrow HFc$) must be an isomorphism. ■

Proposition A.5.20. *Let C and D be any two categories, and let $F : C \rightarrow D$ and $G : D \rightarrow C$ be any two functors. The following are equivalent:*

- (i) $F \dashv G$, with unit $\eta : \text{id}_C \Rightarrow GF$ and counit $\varepsilon : FG \Rightarrow \text{id}_D$.
- (ii) (F, ε) is an absolute right Kan extension of id_D along G .
- (iii) (F, ε) is a right Kan extension of id_D along G that is preserved by F .
- (iv) (G, η) is an absolute left Kan extension of id_C along F .
- (v) (G, η) is a left Kan extension of id_C along F that is preserved by G .

Proof. See [CWM, Ch. X, § 7]. □

Proposition A.5.21.

- Left adjoints preserve all left Kan extensions.
- Right adjoints preserve all right Kan extensions.

Proof. See Theorem I in [CWM, Ch. X, § 5]. □

Definition A.5.22. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{C} be a locally \mathbf{U} -small category. A **dense functor** is a functor $F : \mathcal{B} \rightarrow \mathcal{C}$ such that the F -nerve functor $N^F : \mathcal{C} \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}]$ is fully faithful. A **dense subcategory** of \mathcal{C} is a subcategory \mathcal{B} such that the inclusion $\mathcal{B} \hookrightarrow \mathcal{C}$ is a dense functor.

Dually, a **codense functor** is a functor $F : \mathcal{B} \rightarrow \mathcal{C}$ such that the opposite functor $F^{\text{op}} : \mathcal{B}^{\text{op}} \rightarrow \mathcal{C}^{\text{op}}$ is dense, and a **codense subcategory** of \mathcal{C} is a subcategory \mathcal{B} such that the inclusion $\mathcal{B} \hookrightarrow \mathcal{C}$ is a codense functor.

Example A.5.23. The Yoneda lemma implies $\text{id}_{\mathcal{C}} : \mathcal{C} \rightarrow \mathcal{C}$ is a dense and codense functor.

One may extract an elementary definition for ‘(co)dense functor’ from the following proposition.

Proposition A.5.24. *With notation as in definition A.5.22, the following are equivalent:*

- (i) $F : \mathcal{B} \rightarrow \mathcal{C}$ is a dense functor.
- (ii) For each object C in \mathcal{C} , the maps

$$\mathcal{C}(C, C') \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}](N^F(C), \mathcal{C}(F, C'))$$

induced by $N^F : \mathcal{C} \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}]$ are natural bijections, exhibiting C as a weighted colimit $N^F(C) \star_{\mathcal{B}} F$ in \mathcal{C} .

- (iii) For each object C in \mathcal{C} , the tautological cocone to C induced by F is a colimiting cocone.
- (iv) $(\text{id}_{\mathcal{C}}, \text{id}_F)$ is a pointwise left Kan extension of F along F .

Dually, the following are equivalent:

- (i') $F : \mathcal{B} \rightarrow \mathcal{C}$ is a codense functor.
- (ii') For each object C in \mathcal{C} , the maps

$$\mathcal{C}(C', C) \rightarrow [\mathcal{B}, \mathbf{Set}](N^{F^{\text{op}}}(C), \mathcal{C}(C', F))$$

induced by $N^{F^{\text{op}}} : \mathcal{C}^{\text{op}} \rightarrow [\mathcal{B}, \mathbf{Set}]$ are natural bijections, exhibiting C as a weighted limit $\{N^{F^{\text{op}}}(C), F\}^{\mathcal{B}}$ in \mathcal{C} .

(iii') For each object C in \mathcal{C} , the tautological cone from C induced by F is a limiting cone.

(iv') $(\text{id}_C, \text{id}_F)$ is a pointwise right Kan extension of F along F .

Proof. (i) \Leftrightarrow (ii). The indicated maps are bijections for all C and C' if and only if N^F is fully faithful, by definition.

(ii) \Leftrightarrow (iii) \Leftrightarrow (iv). This is an application of [theorem A.5.15](#). ■

Definition A.5.25. Let $G : \mathcal{D} \rightarrow \mathcal{C}$ be a functor. A **densely-defined partial left adjoint** for G is a triple (F, i, η) , where $F : \mathcal{B} \rightarrow \mathcal{D}$ is a functor, $i : \mathcal{B} \rightarrow \mathcal{C}$ is a dense functor, and $\eta : i \Rightarrow GF$ is a natural transformation such that the maps

$$\begin{aligned} \mathcal{D}(FB, D) &\rightarrow \mathcal{C}(iB, GD) \\ g &\mapsto Gg \circ \eta_B \end{aligned}$$

are bijections that are natural in B and D .

Dually, given a functor $F : \mathcal{C} \rightarrow \mathcal{D}$, a **codensely-defined partial right adjoint** for F is a triple (G, j, ε) , where $G : \mathcal{B} \rightarrow \mathcal{C}$ is a functor, $j : \mathcal{B} \rightarrow \mathcal{C}$ is a codense functor, and $\varepsilon : FG \Rightarrow j$ is a natural transformation such that the maps

$$\begin{aligned} \mathcal{C}(C, GB) &\rightarrow \mathcal{D}(FC, jB) \\ f &\mapsto \varepsilon_B \circ Ff \end{aligned}$$

are bijections that are natural in B and C .

Example A.5.26. The Yoneda embedding $\hat{h}_\bullet : \mathcal{B} \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}]$ has a densely-defined partial left adjoint, namely $(\text{id}_B, \hat{h}_\bullet, \text{id}_{\hat{h}_\bullet})$.

REMARK A.5.27. (F, id_C, η) is a densely-defined partial left adjoint for G if and only if F is a left adjoint for G in the usual sense, with η being the adjunction unit.

Proposition A.5.28. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{C} and \mathcal{D} be locally \mathbf{U} -small categories. Given functors $G : \mathcal{D} \rightarrow \mathcal{C}$, $F : \mathcal{B} \rightarrow \mathcal{D}$, and $i : \mathcal{B} \rightarrow \mathcal{C}$, the following are equivalent:

(i) (F, i, η) is a densely-defined partial left adjoint for G .

(ii) The functor $i : \mathcal{B} \rightarrow \mathcal{C}$ is dense, and there exists a diagram

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{h_\bullet} & [\mathcal{D}^{\text{op}}, \mathbf{Set}] \\ G \downarrow & \swarrow \alpha & \downarrow (F^{\text{op}})^* \\ \mathcal{C} & \xrightarrow{N^i} & [\mathcal{B}^{\text{op}}, \mathbf{Set}] \end{array}$$

where α factors through $\eta^* : \mathbf{N}^{GF} \Rightarrow \mathbf{N}^i$ and is a natural isomorphism.

(iii) The functor $i : \mathcal{B} \rightarrow \mathcal{C}$ is dense, and the diagram

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{h_\bullet} & [\mathcal{D}^{\text{op}}, \mathbf{Set}] \\ G \downarrow & & \downarrow (F^{\text{op}})^* \\ \mathcal{C} & \xrightarrow{N^i} & [\mathcal{B}^{\text{op}}, \mathbf{Set}] \end{array}$$

commutes up to natural isomorphism.

Dually, given functors $F : \mathcal{C} \rightarrow \mathcal{D}$, $G : \mathcal{B} \rightarrow \mathcal{C}$, and $j : \mathcal{B} \rightarrow \mathcal{D}$, the following are equivalent:

(i') (G, j, ε) is a codensely-defined partial right adjoint for F .

(ii') The functor $j : \mathcal{B} \rightarrow \mathcal{D}$ is codense, and there exists a diagram

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} & \xrightarrow{h^\bullet} & [\mathcal{C}, \mathbf{Set}] \\ F^{\text{op}} \downarrow & \swarrow \beta & \downarrow G^* \\ \mathcal{D}^{\text{op}} & \xrightarrow{N^{j^{\text{op}}}} & [\mathcal{B}, \mathbf{Set}] \end{array}$$

where β factors through $(\varepsilon^{\text{op}})^* : \mathbf{N}^{F^{\text{op}}G^{\text{op}}} \Rightarrow \mathbf{N}^{j^{\text{op}}}$ and is a natural isomorphism.

(iii') The functor $j : \mathcal{B} \rightarrow \mathcal{D}$ is codense, and the diagram

$$\begin{array}{ccc} \mathcal{C}^{\text{op}} & \xrightarrow{h^\bullet} & [\mathcal{C}, \mathbf{Set}] \\ F^{\text{op}} \downarrow & & \downarrow G^* \\ \mathcal{D}^{\text{op}} & \xrightarrow{N^{j^{\text{op}}}} & [\mathcal{B}, \mathbf{Set}] \end{array}$$

commutes up to natural isomorphism.

Proof. (i) \Rightarrow (ii). This immediately follows from the definition.

(ii) \Rightarrow (iii). Obvious.

(iii) \Rightarrow (i). The displayed diagram commutes up to natural isomorphism precisely when there are bijections

$$\alpha_{B,D} : D(FB, D) \rightarrow C(iB, GD)$$

that are natural in both B and D . Taking $D = FB$, let $\eta_B : iB \rightarrow GFB$ be the morphism corresponding to $\text{id}_{FB} : FB \rightarrow FB$. Applying the Yoneda lemma, we see that the natural bijection $\alpha_{B,D}$ must be the map $g \mapsto Gg \circ \eta_B$. ■

Corollary A.5.29. *Let C and D be any two categories. If a functor $G : D \rightarrow C$ has a densely-defined partial left adjoint, then G preserves:*

- (i) *limits for all diagrams in D ,*
- (ii) *weighted limits, and*
- (iii) *pointwise right Kan extensions.*

Dually, if a functor $F : C \rightarrow D$ has a codensely-defined partial right adjoint, then F preserves:

- (i') *colimits for all diagrams in C ,*
- (ii') *weighted colimits, and*
- (iii') *pointwise left Kan extensions.*

Proof. Choose a universe \mathbf{U} such that the domain of $i : \mathcal{B} \rightarrow \mathcal{C}$ is \mathbf{U} -small and both C and D are locally \mathbf{U} -small, and consider the following diagram:

$$\begin{array}{ccc} \mathcal{D} & \xrightarrow{h_\bullet} & [\mathcal{D}^{\text{op}}, \mathbf{Set}] \\ G \downarrow & & \downarrow (F^{\text{op}})^* \\ \mathcal{C} & \xrightarrow{N^i} & [\mathcal{B}^{\text{op}}, \mathbf{Set}] \end{array}$$

Since i is dense, the i -nerve functor $N^i : C \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}]$ is fully faithful. [Corollary A.5.17](#) implies $(F^{\text{op}})^* : [\mathcal{D}^{\text{op}}, \mathbf{Set}] \rightarrow [\mathcal{B}^{\text{op}}, \mathbf{Set}]$ is a right adjoint, and the Yoneda embedding $h_\bullet : \mathcal{D} \rightarrow [\mathcal{D}^{\text{op}}, \mathbf{Set}]$ preserves all limits and weighted limits (see [proposition A.5.13](#)), so we use the fact that N^i reflects limits and weighted limits to conclude that G preserves them. We then apply [corollary A.5.18](#). ■

A.6 Ends and coends

Prerequisites. §§ 0.1, A.5

In this section we use the explicit universe convention.

Definition A.6.1. Let $F, G : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D}$ be functors. A **dinatural transformation** $\alpha : F \overset{\diamond}{\rightarrow} G$ is a family $(\alpha_C : F(C, C) \rightarrow G(C, C) \mid C \in \text{ob } \mathcal{C})$ such that the diagram

$$\begin{array}{ccccc}
 & & F(C, C) & \xrightarrow{\alpha_C} & G(C, C) & & \\
 & & \nearrow^{F(\text{id}_C, f)} & & \searrow^{G(f, \text{id}_C)} & & \\
 & F(C, C') & & & & & G(C', C) \\
 & \searrow^{F(f, \text{id}_{C'})} & & & \nearrow^{G(\text{id}_{C'}, f)} & & \\
 & & F(C', C') & \xrightarrow{\alpha_{C'}} & G(C', C') & &
 \end{array}$$

commutes for all morphisms $f : C' \rightarrow C$ in \mathcal{C} .

Example A.6.2. Let \mathbf{U} be a pre-universe, let \mathcal{C} be a locally \mathbf{U} -small category, and let \mathbf{Set} be the category of \mathbf{U} -sets. Consider the functor $\text{Hom}_{\mathcal{C}} : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathbf{Set}$ that sends a pair of objects in \mathcal{C} to their hom-set. For each natural number n , we have an dinatural transformation $\text{Hom}_{\mathcal{C}} \overset{\diamond}{\rightarrow} \text{Hom}_{\mathcal{C}}$ defined by $e \mapsto e^n$, where e^n denotes the n -fold iterate of the endomorphism e .

Definition A.6.3. A **wedge** from an object D in \mathcal{D} to a functor $G : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D}$ is a dinatural transformation $\Delta D \overset{\diamond}{\rightarrow} G$, where $\Delta D : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D}$ is the constant functor with value D ; dually, a **cowedge** from a functor $F : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D}$ to an object D in \mathcal{D} is a dinatural transformation $F \overset{\diamond}{\rightarrow} \Delta D$.

Definition A.6.4. An **end** for a functor $G : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{D}$ is an object E and a wedge $\lambda : \Delta E \overset{\diamond}{\rightarrow} G$ with the following universal property:

- For each wedge $\varphi : \Delta D \overset{\diamond}{\rightarrow} G$, there is a unique morphism $f : D \rightarrow E$ in \mathcal{D} such that $\varphi_C = \lambda_C \circ f$ for all objects C in \mathcal{C} .

We write the following formula to mean that E is an end for G :

$$E = \int_{\mathcal{C}:\mathcal{C}} G(C, C)$$

Dually, a **coend** for a functor $F : C^{\text{op}} \times C \rightarrow D$ is an object E and a cowedge $\lambda : F \overset{\diamond}{\rightarrow} \Delta E$ with the following universal property:

- For each cowedge $\varphi : F \overset{\diamond}{\rightarrow} \Delta D$, there is a unique morphism $f : E \rightarrow D$ in D such that $\varphi_C = f \circ \lambda_C$ for all objects C in C .

We write the following formula to mean that E is a coend for F :

$$E = \int^{C:C} F(C, C)$$

REMARK A.6.5. Let \mathbf{U} be a pre-universe, let \mathbb{D} be a \mathbf{U} -small category, and let C be a locally \mathbf{U} -small category. Then, for all functors $F, G : \mathbb{D} \rightarrow C$, we have a bijection

$$[\mathbb{D}, C](F, G) \cong \int_{d:\mathbb{D}} C(Fd, Gd)$$

and this is natural in both F and G . (The size restriction ensures that the LHS is a \mathbf{U} -set.) See also [lemma A.5.12](#).

Proposition A.6.6. *Let \mathbf{U} be a pre-universe and let \mathbb{D} be a \mathbf{U} -small category. If C is a \mathbf{U} -complete category, then C has ends for all functors $A : \mathbb{D}^{\text{op}} \times \mathbb{D} \rightarrow C$. Dually, if C is a \mathbf{U} -cocomplete category, then C has coends for all functors $A : \mathbb{D}^{\text{op}} \times \mathbb{D} \rightarrow C$.*

Proof. It is clear from the definition that an end is a special kind of limit, and a coend is a special kind of colimit. To make this precise, one can use Mac Lane's subdivision category C^\S : see [[CWM](#), Ch. IX, § 5]. \square

Proposition A.6.7. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let $F : C \rightarrow D$ be any functor where C and D are locally \mathbf{U} -small categories.*

- (i) *For any functor $A : \mathcal{J}^{\text{op}} \times \mathcal{J} \rightarrow C$, if the ends $\int_{\mathcal{J}} A$ and $\int_{\mathcal{J}} FA$ both exist, with λ being the universal wedge in C , then there is a canonical comparison morphism*

$$F \int_{\mathcal{J}} A \rightarrow \int_{\mathcal{J}} FA$$

induced by the wedge $F\lambda$.

- (ii) *For any object C in C , the functor $C(C, -) : C \rightarrow \mathbf{Set}$ preserves all ends.*
(iii) *The functors $C(C, -)$ jointly reflect ends.*

(iv) If F has a left adjoint, then F preserves ends.

Dually:

(i') For any functor $A : \mathcal{J}^{\text{op}} \times \mathcal{J} \rightarrow \mathcal{C}$, if the coends $\int^{\mathcal{J}} A$ and $\int^{\mathcal{J}} FA$ both exist, with λ being the universal cowedge in \mathcal{C} , then there is a canonical comparison morphism

$$\int^{\mathcal{J}} FA \rightarrow F \int^{\mathcal{J}} A$$

induced by the cowedge $F\lambda$.

(ii') For any object C in \mathcal{C} , the functor $\mathcal{C}(-, C) : \mathcal{C} \rightarrow \mathbf{Set}$ sends any coend in \mathcal{C} to the corresponding end in \mathbf{Set} .

(iii') The functors $\mathcal{C}(-, C) : \mathcal{C} \rightarrow \mathbf{Set}^{\text{op}}$ jointly reflect coends.

(iv') If F has a right adjoint, then F preserves coends.

Proof. All straightforward. ◆

Definition A.6.8. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let $\mathbb{1}$ be the trivial category with $*$ as its only object. A **tensoring \mathbf{U} -category** is a locally \mathbf{U} -small category \mathcal{C} such that, for all weights $W : \mathbb{1} \rightarrow \mathbf{Set}$ and all diagrams $A : \mathbb{1} \rightarrow \mathbf{Set}$, a W -weighted colimit for A exists in \mathcal{C} ; if \mathcal{C} is a tensoring \mathbf{U} -category, then we write $X \odot C$ for the weighted colimit $W \star_{\mathbb{1}} A$, where $X = W(*)$ and $C = A(*)$.

Dually, a **cotensoring \mathbf{U} -category** is a locally \mathbf{U} -small category \mathcal{C} such that, for all weights $W : \mathbb{1} \rightarrow \mathbf{Set}$ and all diagrams $A : \mathbb{1} \rightarrow \mathbf{Set}$, a W -weighted limit for A exists in \mathcal{C} ; if \mathcal{C} is a cotensoring \mathbf{U} -category, then we write $X \pitchfork C$ for the weighted limit $\{W, A\}^{\mathbb{1}}$, where $X = W(*)$ and $C = A(*)$.

Proposition A.6.9 (Tensor–hom–cotensor adjunction). *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, let \mathcal{C} be a locally \mathbf{U} -small category.*

(i) *If \mathcal{C} is a tensoring \mathbf{U} -category, then the assignment $(X, C) \mapsto X \odot C$ can be extended to a functor $\mathbf{Set} \times \mathcal{C} \rightarrow \mathcal{C}$ such that, for each object C , we have the following adjunction:*

$$- \odot C \dashv \mathcal{C}(C, -) : \mathcal{C} \rightarrow \mathbf{Set}$$

- (ii) If \mathcal{C} is a cotensored \mathbf{U} -category, then the assignment $(X, C) \mapsto X \pitchfork C$ can be extended to a functor $\mathbf{Set}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$ such that, for each object C , the functors $- \pitchfork C : \mathbf{Set}^{\text{op}} \rightarrow \mathcal{C}$ and $\mathcal{C}(-, C) : \mathcal{C}^{\text{op}} \rightarrow \mathbf{Set}$ are contravariantly adjoint on the right.
- (iii) If \mathcal{C} is a tensored and cotensored \mathbf{U} -category, then for each set X , we have the following adjunction:

$$X \odot - \dashv X \pitchfork - : \mathcal{C} \rightarrow \mathcal{C}$$

Proof. Claims (i) and (ii) are formally dual and are straightforward applications of the parametrised adjunction theorem.^[9] For claim (iii), simply observe that we have bijections

$$\mathcal{C}(X \odot A, B) \cong \mathbf{Set}(X, \mathcal{C}(A, B)) \cong \mathcal{C}(A, X \pitchfork B)$$

and these are natural in A , B , and X . ■

Theorem A.6.10. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{C} be a locally \mathbf{U} -small category. The following are equivalent:*

- (i) \mathcal{C} is a \mathbf{U} -complete category.
- (ii) \mathcal{C} is a cotensored \mathbf{U} -category and, for all \mathbf{U} -small categories \mathbb{D} and all functors $B : \mathbb{D}^{\text{op}} \times \mathbb{D} \rightarrow \mathcal{C}$, an end for A exists in \mathcal{C} .
- (iii) For all weights $W : \mathbb{D}^{\text{op}} \rightarrow \mathbf{Set}$ and all diagrams $A : \mathbb{D} \rightarrow \mathbf{Set}$, \mathcal{C} has a W -weighted limit for A , provided \mathbb{D} is a \mathbf{U} -small category.

Dually, the following are equivalent:

- (i') \mathcal{C} is a \mathbf{U} -cocomplete category.
- (ii') \mathcal{C} is a tensored \mathbf{U} -category and, for all \mathbf{U} -small categories \mathbb{D} and all functors $B : \mathbb{D}^{\text{op}} \times \mathbb{D} \rightarrow \mathcal{C}$, a coend for A exists in \mathcal{C} .
- (iii') For all weights $W : \mathbb{D}^{\text{op}} \rightarrow \mathbf{Set}$ and all diagrams $A : \mathbb{D} \rightarrow \mathbf{Set}$, \mathcal{C} has a W -weighted colimit for A , provided \mathbb{D} is a \mathbf{U} -small category.

[9] See Theorem 3 in [CWM, Ch. IV, § 7].

Proof. (i) \Rightarrow (ii). It is clear that $X \pitchfork C$ is nothing more than an X -fold product of copies of C , so C is certainly \mathbf{U} -cotensored if it is \mathbf{U} -complete, and [proposition A.6.6](#) says C also has the required ends in that case.

(ii) \Rightarrow (iii). We have the following natural bijections:

$$\begin{aligned} \mathcal{C}(C, \{W, A\}^{\mathbb{D}}) &\cong [\mathbb{D}, \mathbf{Set}](W, \mathcal{C}(C, A)) \\ &\cong \int_{d:\mathbb{D}} \mathbf{Set}(Wd, \mathcal{C}(C, Ad)) \\ &\cong \int_{d:\mathbb{D}} \mathcal{C}(C, Wd \pitchfork Ad) \\ &\cong \mathcal{C}\left(C, \int_{d:\mathbb{D}} Wd \pitchfork Ad\right) \end{aligned}$$

Thus, using the Yoneda lemma and *assuming C is a cotensored \mathbf{U} -category*, the weighted limit $\{W, A\}^{\mathbb{D}}$ exists if and only if the end $\int_{d:\mathbb{D}} Wd \pitchfork Ad$ exists.

(iii) \Rightarrow (i). Ordinary limits are a special case of weighted limits, as remarked in [A.5.II](#). ■

Proposition A.6.II. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, let \mathcal{C} be a locally \mathbf{U} -small category, and let \mathcal{J} be any category. If \mathcal{C} is a tensored \mathbf{U} -category and has weighted limits for all weights $W : \mathcal{J} \rightarrow \mathbf{Set}$ and diagrams $A : \mathcal{J} \rightarrow \mathcal{C}$, then:*

- (i) $(W, A) \mapsto \{W, A\}^{\mathcal{J}}$ extends to a functor $[\mathcal{J}, \mathbf{Set}]^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$.
- (ii) For each diagram $A : \mathcal{J} \rightarrow \mathcal{C}$, the functors $\{-, A\}^{\mathcal{J}} : [\mathcal{J}, \mathbf{Set}]^{\text{op}} \rightarrow \mathcal{C}$ and $\mathcal{C}(-, A) : \mathcal{C}^{\text{op}} \rightarrow [\mathcal{J}, \mathbf{Set}]$ are contravariantly adjoint on the right.
- (iii) For each weight $W : \mathcal{J} \rightarrow \mathbf{Set}$, we have the following adjunction:

$$W \odot - \dashv \{W, -\}^{\mathcal{J}} : [\mathcal{J}, \mathcal{C}] \rightarrow \mathcal{C}$$

Here, $W \odot C : \mathcal{J} \rightarrow \mathcal{C}$ is the diagram $j \mapsto Wj \odot C$.

Dually, if \mathcal{C} is a cotensored \mathbf{U} -category and has weighted colimits for all weights $W : \mathcal{J}^{\text{op}} \rightarrow \mathbf{Set}$ and diagrams $A : \mathcal{J} \rightarrow \mathcal{C}$, then:

- (i') $(W, A) \mapsto W \star_{\mathcal{J}} A$ extends to a functor $[\mathcal{J}^{\text{op}}, \mathbf{Set}] \times \mathcal{C} \rightarrow \mathcal{C}$.

(ii') For each diagram $A : \mathcal{J} \rightarrow \mathcal{C}$, we have the following adjunction:

$$- \star_{\mathcal{J}} A \dashv \mathcal{C}(A, -) : \mathcal{C} \rightarrow [\mathcal{J}^{\text{op}}, \mathbf{Set}]$$

(iii') For each weight $W : \mathcal{J}^{\text{op}} \rightarrow \mathbf{Set}$, we have the following adjunction:

$$W \star_{\mathcal{J}} - \dashv W \pitchfork - : \mathcal{C} \rightarrow [\mathcal{J}, \mathcal{C}]$$

Here, $W \pitchfork C : \mathcal{J} \rightarrow \mathcal{C}$ is the diagram $j \mapsto W j \pitchfork C$.

Proof. Claim (i) is straightforward, and for claims (ii) and (iii), observe that we have bijections

$$\begin{aligned} \mathcal{C}(C, \{W, A\}^{\mathcal{J}}) &\cong [\mathcal{J}, \mathbf{Set}](W, \mathcal{C}(C, A)) \\ &\cong \int_{j:\mathcal{J}} \mathbf{Set}(W j, \mathcal{C}(C, A j)) \\ &\cong \int_{j:\mathcal{J}} \mathcal{C}(W j \odot C, A j) \\ &\cong [\mathcal{J}, \mathcal{C}](W \odot C, A) \end{aligned}$$

and these are natural in W , A , and C . ■

Lemma A.6.12. Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathbb{I} and \mathbb{J} be \mathbf{U} -small categories. For all functors $A : \mathbb{I}^{\text{op}} \times \mathbb{J}^{\text{op}} \times \mathbb{I} \times \mathbb{J} \rightarrow \mathbf{Set}$:

- (i) The assignment $(i', i) \mapsto \int_{j:\mathbb{J}} A(i', j, i, j)$ extends to a functor $\mathbb{I}^{\text{op}} \times \mathbb{I} \rightarrow \mathbf{Set}$.
- (ii) There is a unique morphism θ making the diagram below commute for all i and j ,

$$\begin{array}{ccc} \int_{i':\mathbb{I}} \int_{j':\mathbb{J}} A(i', j', i', j') & \longrightarrow & \int_{j':\mathbb{J}} A(i, j', i, j') \\ \downarrow \theta & & \downarrow \\ \int_{(i', j'):\mathbb{I} \times \mathbb{J}} A(i', j', i', j') & \longrightarrow & A(i, j, i, j) \end{array}$$

where the unlabelled arrows are the components of the respective universal wedges, and θ is moreover an isomorphism.

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(iii) *There is a unique morphism σ making the diagram below commute for all i and j ,*

$$\begin{array}{ccc}
 \int_{i':\mathbb{I}} \int_{j':\mathbb{J}} A(i', j', i', j') & \longrightarrow & \int_{j':\mathbb{J}} A(i, j', i, j') \\
 \downarrow \sigma & & \searrow \\
 \int_{j':\mathbb{J}} \int_{i':\mathbb{I}} A(i', j', i', j') & \longrightarrow & \int_{i':\mathbb{I}} A(i, j', i, j') \\
 & & \nearrow \\
 & & A(i, j, i, j)
 \end{array}$$

where the unmarked arrows are the components of the respective universal wedges, and σ is moreover an isomorphism.

Proof. See [CWM, Ch. IX, § 8]. □

Theorem A.6.13 (Interchange law for ends and coends). *Let \mathcal{C} be any category and let $A : \mathcal{I}^{\text{op}} \times \mathcal{J}^{\text{op}} \times \mathcal{I} \times \mathcal{J} \rightarrow \mathbf{Set}$ be any functor. If the end $\int_{i:\mathcal{I}} A(i, j', i, j)$ exists in \mathcal{C} for all j' and j in \mathcal{J} , and the end $\int_{j:\mathcal{J}} A(i', j, i, j)$ exists in \mathcal{C} for all i' and i in \mathcal{I} , then the following are equivalent:*

- (i) *The end $\int_{(i,j):\mathcal{I} \times \mathcal{J}} A(i, j, i, j)$ exists in \mathcal{C} .*
- (ii) *The iterated end $\int_{i:\mathcal{I}} \int_{j:\mathcal{J}} A(i, j, i, j)$ exists in \mathcal{C} .*
- (iii) *The iterated end $\int_{j:\mathcal{J}} \int_{i:\mathcal{I}} A(i, j, i, j)$ exists in \mathcal{C} .*

In this case, we have a canonical isomorphism in \mathcal{C} :

$$\int_{i:\mathcal{I}} \int_{j:\mathcal{J}} A(i, j, i, j) \cong \int_{j:\mathcal{J}} \int_{i:\mathcal{I}} A(i, j, i, j)$$

Dually, if the coend $\int^{i:\mathcal{I}} A(i, j', i, j)$ exists in \mathcal{C} for all j' and j in \mathcal{J} , and the coend $\int^{j:\mathcal{J}} A(i', j, i, j)$ exists in \mathcal{C} for all i' and i in \mathcal{I} , then the following are equivalent:

- (i') *The coend $\int^{(i,j):\mathcal{I} \times \mathcal{J}} A(i, j, i, j)$ exists in \mathcal{C} .*
- (ii') *The iterated coend $\int^{i:\mathcal{I}} \int^{j:\mathcal{J}} A(i, j, i, j)$ exists in \mathcal{C} .*
- (iii') *The iterated coend $\int^{j:\mathcal{J}} \int^{i:\mathcal{I}} A(i, j, i, j)$ exists in \mathcal{C} .*

In this case, we have a canonical isomorphism in \mathcal{C} :

$$\int^{i:\mathcal{I}} \int^{j:\mathcal{J}} A(i, j, i, j) \cong \int^{j:\mathcal{J}} \int^{i:\mathcal{I}} A(i, j, i, j)$$

Proof. Choose a pre-universe \mathbf{U} such that \mathcal{I} and \mathcal{J} are \mathbf{U} -small categories and \mathcal{C} is a locally \mathbf{U} -small category, and use the Yoneda lemma to reduce the claims to the previous lemma. \blacksquare

Proposition A.6.14. *Let \mathbf{U} be a pre-universe, let \mathbf{Set} be the category of \mathbf{U} -sets, and let \mathcal{C} and \mathcal{J} be locally \mathbf{U} -small categories.*

(i) *For all j in \mathcal{J} and all functors $A : \mathcal{J} \rightarrow \mathcal{C}$, the Yoneda bijection*

$$\mathcal{C}(\mathcal{C}, A_j) \cong [\mathcal{J}, \mathbf{Set}](h^j, \mathcal{C}(\mathcal{C}, A))$$

exhibits A_j as the weighted limit $\{h^j, A\}^{\mathcal{J}}$ in \mathcal{C} .

(ii) *If \mathcal{C} is a cotensored \mathbf{U} -category, then the end $\int_{j', j} \mathcal{J}(j, j') \pitchfork A_j'$ exists in \mathcal{C} and can be canonically identified with A_j .*

(iii) *For all functors $H : \mathcal{J}^{\text{op}} \times \mathcal{J} \rightarrow \mathcal{C}$, the weighted limit $\{\text{Hom}_{\mathcal{J}}, H\}^{\mathcal{J}^{\text{op}} \times \mathcal{J}}$ exists in \mathcal{C} if and only if the end $\int_{j, j} H(j, j)$ exists in \mathcal{C} , and there is a canonical identification of the two.*

Dually:

(i') *For all j in \mathcal{J} and all functors $A : \mathcal{J} \rightarrow \mathcal{C}$, the Yoneda bijection*

$$\mathcal{C}(A_j, \mathcal{C}) \cong [\mathcal{J}^{\text{op}}, \mathbf{Set}](h_j, \mathcal{C}(A, \mathcal{C}))$$

exhibits A_j as the weighted colimit $h_j \star_{\mathcal{J}} A$ in \mathcal{C} .

(ii') *If \mathcal{C} is a tensored \mathbf{U} -category, then the coend $\int^{j', j} \mathcal{J}(j', j) \odot A_j'$ exists in \mathcal{C} and can be canonically identified with A_j .*

(iii') *For all functors $H : \mathcal{J}^{\text{op}} \times \mathcal{J} \rightarrow \mathcal{C}$, the weighted colimit $\text{Hom}_{\mathcal{J}^{\text{op}}} \star_{\mathcal{J}^{\text{op}} \times \mathcal{J}} H$ exists in \mathcal{C} if and only if the coend $\int^{j, j} H(j, j)$ exists in \mathcal{C} , and there is a canonical identification of the two.*

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Proof. (i). This is an immediate consequence of the Yoneda lemma and the definition of weighted limit.

(ii). Use the identification constructed in the proof of [theorem A.6.10](#).

(iii). For all objects C in \mathcal{C} , using claim (ii) and the interchange law for ends ([theorem A.6.13](#)), there are bijections

$$\begin{aligned} [\mathcal{J}^{\text{op}} \times \mathcal{J}, \mathbf{Set}](\text{Hom}_{\mathcal{J}} C(C, H)) &\cong \int_{(j', j): \mathcal{J}^{\text{op}} \times \mathcal{J}} \mathbf{Set}(\mathcal{J}(j', j), C(H(j', j))) \\ &\cong \int_{j: \mathcal{J}} \int_{j': \mathcal{J}^{\text{op}}} \mathbf{Set}(\mathcal{J}(j', j), C(H(j', j))) \\ &\cong \int_{j: \mathcal{J}} C(C, H(j, j)) \end{aligned}$$

and these are natural in C ; now apply propositions [A.5.13](#) and [A.6.7](#). ■

HIGHER GENERALITIES

B.I Monoidal categories

Standard references for monoidal categories include [CWM, Ch. VII and Ch. XI] and [Kelly, 2005, Ch. 1]. To fix notation, we will quickly review the main definitions in the theory of monoidal categories.

Definition B.I.1. A **strict monoidal category** is a category \mathcal{C} together with an object I and a functor $\otimes : \mathcal{C} \times \mathcal{C} \rightarrow \mathcal{C}$ satisfying the following axioms:

- (Left unit). $I \otimes (-) = \text{id}_{\mathcal{C}}$.
- (Right unit). $(-) \otimes I = \text{id}_{\mathcal{C}}$.
- (Associativity). For all objects $X, Y,$ and Z in \mathcal{C} ,

$$(X \otimes Y) \otimes Z = X \otimes (Y \otimes Z)$$

and similarly for morphisms in \mathcal{C} .

I is called the **monoidal unit**, and \otimes is called the **monoidal product**.

In short, a strict monoidal category is an internal monoid in the metacategory of all categories.

Example B.I.2. For any category \mathcal{C} , the endofunctor category $[\mathcal{C}, \mathcal{C}]$ is a strict monoidal category with $\text{id}_{\mathcal{C}}$ as the monoidal unit and endofunctor composition as the monoidal product.

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Despite the above example, strict monoidal categories turn out to be less useful than one might hope: not even **Set** equipped with the usual cartesian product is a strict monoidal category.^[1] The problem is in the *equations* we have imposed in the axioms above: in naturally-occurring examples, we do not get *identities* but only natural isomorphisms. This observation led Bénabou [1963] to propose the following notion instead:

Definition B.I.3. A **monoidal category** is a category C together with an object I , a functor $(-) \otimes (-) : C \times C \rightarrow C$, and three natural isomorphisms λ , ρ , and α ,^[2] of type

$$\begin{aligned}\lambda_X &: I \otimes X \xrightarrow{\cong} X \\ \rho_X &: X \otimes I \xrightarrow{\cong} X \\ \alpha_{X,Y,Z} &: (X \otimes Y) \otimes Z \xrightarrow{\cong} X \otimes (Y \otimes Z)\end{aligned}$$

such that the following diagrams commute for all choices of objects in C :

$$\begin{array}{ccc} & (W \otimes (X \otimes Y)) \otimes Z & \\ \alpha_{W,X,Y} \otimes \text{id}_Z \nearrow & & \searrow \alpha_{W,X \otimes Y,Z} \\ ((W \otimes X) \otimes Y) \otimes Z & & (W \otimes ((X \otimes Y) \otimes Z)) \\ \alpha_{W \otimes X,Y,Z} \downarrow & & \downarrow \text{id}_W \otimes \alpha_{X,Y,Z} \\ (W \otimes X) \otimes (Y \otimes Z) & \xrightarrow{\alpha_{W,X,Y \otimes Z}} & W \otimes (X \otimes (Y \otimes Z)) \end{array}$$

$$\begin{array}{ccc} (X \otimes I) \otimes Y & \xrightarrow{\alpha_{X,I,Y}} & X \otimes (I \otimes Y) \\ \rho_X \otimes \text{id}_Y \searrow & & \swarrow \lambda_X \otimes \text{id}_Y \\ & X \otimes Y & \end{array}$$

The natural isomorphisms λ , ρ , and α are called, respectively, the **left unitor**, **right unitor**, and **associator** of the monoidal category C .

[1] In fact, even if we identify all isomorphic objects, there is still a problem: see the closing remarks in [CWM, Ch. VII, § 1].

[2] Beware: Mac Lane [CWM, Ch. VII] uses the opposite convention for α .

REMARK B.I.4. Since λ , ρ , and α are natural *isomorphisms*, a monoidal structure on \mathcal{C} induces a monoidal structure on \mathcal{C}^{op} . Less obviously, we can define a monoidal category \mathcal{C}^{rev} whose underlying category is the same as \mathcal{C} , but $X \otimes^{\text{rev}} Y = Y \otimes X$, $\lambda^{\text{rev}} = \rho$, $\rho^{\text{rev}} = \lambda$, and $\alpha^{\text{rev}} = \alpha^{-1}$.

¶ B.I.5. A fairly non-trivial theorem of Mac Lane [1963] and Kelly [1964] essentially states that these two axioms are enough to prove that “all diagrams involving only λ , ρ , and α commute”. For example, using the pentagon axiom and the triangle axiom, we may derive

$$\begin{array}{ccc} (I \otimes X) \otimes Y & \xrightarrow{\alpha_{I,X,Y}} & I \otimes (X \otimes Y) \\ & \searrow \lambda_X \otimes \text{id}_Y & \swarrow \lambda_{X \otimes Y} \\ & X \otimes Y & \end{array}$$

from which the equation (!) below can be obtained:

$$\lambda_I = \rho_I$$

Definition B.I.6. Let \mathcal{C} and \mathcal{D} be monoidal categories. A **lax monoidal functor** $\mathcal{C} \rightarrow \mathcal{D}$ consists of a functor $F : \mathcal{C} \rightarrow \mathcal{D}$ of the underlying categories, together with a morphism $\eta : I_{\mathcal{D}} \rightarrow FI_{\mathcal{C}}$ in \mathcal{D} and a natural transformation μ of type $F(-) \otimes_{\mathcal{D}} F(-) \rightarrow F(- \otimes_{\mathcal{C}} -)$ making these diagrams commute:

$$\begin{array}{ccc} I_{\mathcal{D}} \otimes_{\mathcal{D}} FX & \xrightarrow{\eta \otimes \text{id}_{FX}} & FI_{\mathcal{C}} \otimes_{\mathcal{D}} FX \\ \lambda_{FX} \downarrow & & \downarrow \mu_{I_{\mathcal{C}}, X} \\ FX & \xleftarrow{F\lambda_X} & F(I_{\mathcal{C}} \otimes_{\mathcal{C}} X) \end{array} \quad \begin{array}{ccc} FX \otimes_{\mathcal{D}} I_{\mathcal{D}} & \xrightarrow{\text{id}_{FX} \otimes \eta} & FX \otimes_{\mathcal{D}} FI_{\mathcal{C}} \\ \rho_{FX} \downarrow & & \downarrow \mu_{X, I_{\mathcal{C}}} \\ FX & \xleftarrow{F\rho_X} & F(X \otimes_{\mathcal{C}} I_{\mathcal{C}}) \end{array}$$

$$\begin{array}{ccc} (FX \otimes_{\mathcal{D}} FY) \otimes_{\mathcal{D}} FZ & \xrightarrow{\alpha_{FX, FY, FZ}} & FX \otimes_{\mathcal{D}} (FY \otimes_{\mathcal{D}} FZ) \\ \mu_{X, Y} \otimes \text{id}_{FZ} \downarrow & & \downarrow \text{id}_{FX} \otimes \mu_{Y, Z} \\ F(X \otimes_{\mathcal{C}} Y) \otimes_{\mathcal{D}} FZ & & FX \otimes_{\mathcal{D}} F(Y \otimes_{\mathcal{C}} Z) \\ \mu_{X \otimes_{\mathcal{C}} Y, Z} \downarrow & & \downarrow \mu_{X, Y \otimes_{\mathcal{C}} Z} \\ F((X \otimes_{\mathcal{C}} Y) \otimes_{\mathcal{C}} Z) & \xrightarrow{F\alpha_{X, Y, Z}} & F(X \otimes_{\mathcal{C}} (Y \otimes_{\mathcal{C}} Z)) \end{array}$$

An **oplax monoidal functor** $\mathcal{C} \rightarrow \mathcal{D}$ is a lax monoidal functor $\mathcal{C}^{\text{op}} \rightarrow \mathcal{D}^{\text{op}}$. A **strong monoidal functor** is a lax monoidal functor such that η and μ are

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isomorphisms. A **strict monoidal functor** is a lax monoidal functor such that η and μ are *identities*.

Definition B.I.7. Let C and D be monoidal categories and let $F, F' : C \rightarrow D$ be lax monoidal functors. A **monoidal natural transformation** $\varphi : F \Rightarrow F'$ is a natural transformation making the following diagrams commute:

$$\begin{array}{ccc}
 I_D & \xrightarrow{\eta} & FI_C \\
 \text{id} \downarrow & & \downarrow \varphi_{I_C} \\
 I_D & \xrightarrow{\eta'} & F'I_C
 \end{array}
 \qquad
 \begin{array}{ccc}
 FX \otimes_D FY & \xrightarrow{\mu_{X,Y}} & F(X \otimes_C Y) \\
 \varphi_X \otimes_D \varphi_Y \downarrow & & \downarrow \varphi_{X \otimes_C Y} \\
 F'X \otimes_D F'Y & \xrightarrow{\mu'_{X,Y}} & F'(X \otimes_C Y)
 \end{array}$$

REMARK B.I.8. Note that if C and D are both strict monoidal categories, then the diagrams above simplify to more familiar ones:

$$\begin{array}{ccc}
 FX & \xrightarrow{\eta \otimes_D \text{id}_{FX}} & FI_C \otimes_D FX \\
 \text{id}_{FX} \searrow & & \downarrow \mu_{I_C, X} \\
 & & FX
 \end{array}
 \qquad
 \begin{array}{ccc}
 FX & \xrightarrow{\text{id}_{FX} \otimes_D \eta} & FX \otimes_D FI_C \\
 \text{id}_{FX} \searrow & & \downarrow \mu_{X, I_C} \\
 & & FX
 \end{array}$$

$$\begin{array}{ccc}
 & FX \otimes_D FY \otimes_D FZ & \\
 \mu_{X,Y} \otimes_D \text{id}_{FZ} \swarrow & & \searrow \text{id}_{FX} \otimes_D \mu_{Y,Z} \\
 F(X \otimes_C Y) \otimes_D FZ & & FX \otimes_D F(Y \otimes_C Z) \\
 \mu_{X \otimes_C Y, Z} \searrow & & \swarrow \mu_{X, Y \otimes_C Z} \\
 & F(X \otimes_C Y \otimes_C Z) &
 \end{array}$$

Thus, we see one reason for defining lax monoidal functors as we have done: if $\mathbb{1}$ is the terminal category, then a lax monoidal functor $\mathbb{1} \rightarrow D$ is the same thing as an internal monoid^[3] in D , and a monoidal natural transformation of such lax monoidal functors is the same thing as a homomorphism of internal monoids.

Many natural examples of monoidal categories have a “commutative” monoidal product. For example, the cartesian product in any category satisfies $X \times Y \cong Y \times X$. As usual, to do anything useful, we must demand not only the

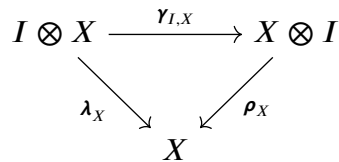
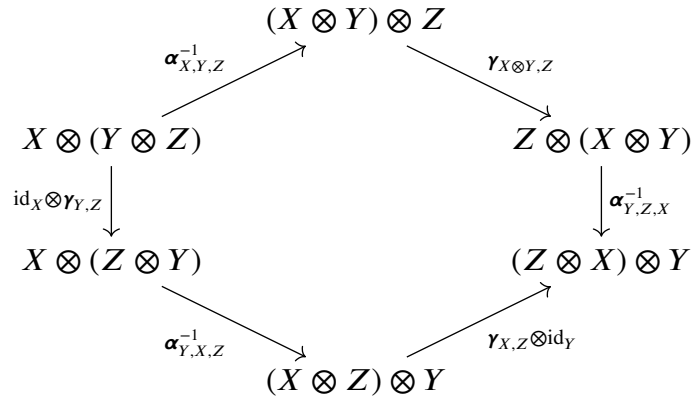
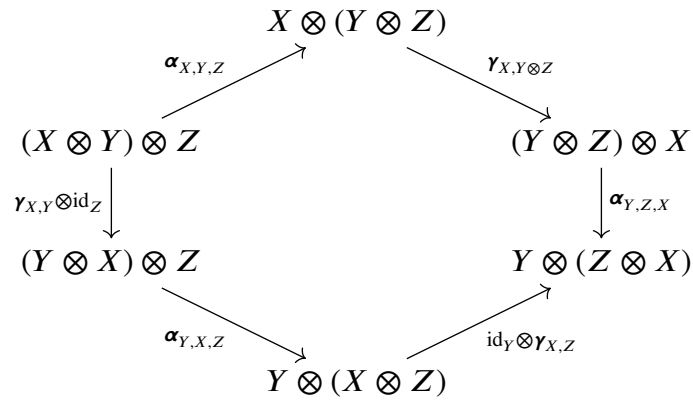
[3] — in the monoidal category sense, of course.

existence of such isomorphisms but also that they be natural and coherent in the following sense:

Definition B.I.9. A **braided monoidal category** is a monoidal category \mathcal{C} together with a natural isomorphism γ of type

$$\gamma_{X,Y} : X \otimes Y \xrightarrow{\cong} Y \otimes X$$

such that the following diagrams commute for all choices of objects in \mathcal{C} :



The natural isomorphism γ is called the **braiding** of \mathcal{C} . A **symmetric monoidal category** is a braided monoidal category \mathcal{C} satisfying the following additional

axiom:

$$\gamma \cdot \gamma = \text{id}_C$$

A **braided / symmetric strict monoidal category** is a braided / symmetric monoidal category that is strict as a monoidal category.

There is a coherence theorem for braided and symmetric monoidal categories as well, but in the braided case it is somewhat subtle compared to the coherence theorem for monoidal categories – we cannot be so cavalier as to say that “all diagrams commute” in a braided monoidal category. Instead, just as before, every braided / symmetric monoidal category is equivalent to a strict one via functors respecting the various structural isomorphisms.

Definition B.I.10. Let C and D be braided monoidal categories. A **lax / oplax / strong / strict braided monoidal functor** $C \rightarrow D$ is a lax / oplax / strong / strict monoidal functor $F : C \rightarrow D$ making the diagram below commute:

$$\begin{array}{ccc} FX \otimes_D FY & \xrightarrow{\mu_{X,Y}} & F(X \otimes_C Y) \\ \gamma_{FX,FY} \downarrow & & \downarrow F\gamma_{X,Y} \\ FY \otimes_D FX & \xrightarrow{\mu_{Y,X}} & F(Y \otimes_C X) \end{array}$$

REMARK B.I.11. The appropriate notion of natural transformation for lax braided monoidal functors is precisely that of a monoidal natural transformation: we need not impose any extra conditions.

Here is an example of an equation that does *not* necessarily hold in a braided monoidal category, even though they have the same domain and codomain:

$$\gamma_{X,Y} \stackrel{?}{=} \gamma_{Y,X}^{-1}$$

Indeed, if it were true, then every braided monoidal category would be a symmetric monoidal category! On the other hand, in a symmetric strict monoidal category, it is true that any two composites of braiding operations with the same domain and codomain are equal – provided each object is identified with a different letter, so that we do not get absurdities like this:

$$\gamma_{X,X} \stackrel{?}{=} \text{id}_{X \otimes X}$$

A similar restriction applies to our claim that “all diagrams commute” in a monoidal category, so it is not unreasonable to say the same is true in a symmetric monoidal category.

We pause briefly to indicate an important special case of a symmetric monoidal category.

Definition B.I.12. A **cartesian monoidal category** is a category with products for all finite families of objects, and a **cartesian monoidal functor** is a functor between cartesian monoidal categories that preserves all finite products.

Proposition B.I.13.

- (i) *A category with all finite products is automatically a symmetric monoidal category, with the terminal object 1 as its monoidal unit and the cartesian product \times as the monoidal product.*
- (ii) *If \mathcal{C} and \mathcal{D} are two categories with finite products regarded as symmetric monoidal categories, then every functor $\mathcal{C} \rightarrow \mathcal{D}$ can be equipped with a canonical oplax braided monoidal functor structure.*
- (iii) *A cartesian monoidal functor is canonically equipped with the structure of a strong braided monoidal functor.*

Proof. (i). The verification of the axioms is straightforward and left to the reader as an exercise.

(ii). Let $F : \mathcal{C} \rightarrow \mathcal{D}$ be a functor. The universal property of the terminal object gives a unique morphism $\epsilon : F1 \rightarrow 1$ in \mathcal{D} , and the universal property of binary products gives a canonical morphism $\delta_{X,Y} : F(X \times Y) \rightarrow FX \times FY$. It can be shown that the diagrams below commute,

$$\begin{array}{ccc}
 F(1_C \times_C X) & \xrightarrow{\delta_{1_C, X}} & F1_C \times_D FX \\
 F\lambda_X \downarrow & & \downarrow \epsilon \times_D \text{id}_{FX} \\
 FX & \xleftarrow{\lambda_{FX}} & 1_D \times_D FX
 \end{array}
 \qquad
 \begin{array}{ccc}
 F(X \times_C 1_C) & \xrightarrow{\delta_{X, 1_C}} & FX \times_D F1_C \\
 F\rho_X \downarrow & & \downarrow \text{id}_{FX} \times_D \epsilon \\
 FX & \xleftarrow{\rho_{FX}} & FX \times_D 1_D
 \end{array}$$

$$\begin{array}{ccc}
 F((X \times_C Y) \times_C Z) & \xrightarrow{F\alpha_{X,Y,Z}} & F(X \times_C (Y \times_C Z)) \\
 \delta_{X \times_C Y, Z} \downarrow & & \downarrow \delta_{X, Y \times_C Z} \\
 F(X \times_C Y) \times_D FZ & & FX \times_D F(Y \times_C Z) \\
 \delta_{X, Y \times_D \text{id}_{FZ}} \downarrow & & \downarrow \text{id}_{FX} \times_D \delta_{Y, Z} \\
 (FX \times_D FY) \times_D FZ & \xrightarrow{\alpha_{FX, FY, FZ}} & FX \times_D (FY \times_D FZ)
 \end{array}$$

$$\begin{array}{ccc}
 F(X \times_C Y) & \xrightarrow{\delta_{X,Y}} & FX \times_D FY \\
 F\gamma_{X,Y} \downarrow & & \downarrow \gamma_{FX, FY} \\
 F(Y \otimes_C X) & \xrightarrow{\delta_{Y,X}} & FY \otimes_D FX
 \end{array}$$

making F into an oplax braided monoidal functor $C \rightarrow D$.

(iii). A functor is cartesian monoidal precisely if ε and δ as defined above are isomorphisms. \diamond

Definition B.I.14. Let Y and Z be objects in a monoidal category C .

- A **right internal hom object** for Y and Z is an object $\mathcal{H}om(Y, Z)$ in C together with a morphism $\text{ev}_{Y,Z} : \mathcal{H}om(Y, Z) \otimes Y \rightarrow Z$ having the following universal property: for all morphisms $f : X \otimes Y \rightarrow Z$ in C , there is a unique morphism $\tilde{f} : X \rightarrow \mathcal{H}om(Y, Z)$ in C such that $\text{ev}_{Y,Z} \circ (\tilde{f} \otimes \text{id}_Y) = f$; equivalently, $\mathcal{H}om(Y, Z)$ is an object in C equipped with bijections

$$C(X \otimes Y, Z) \cong C(X, \mathcal{H}om(Y, Z))$$

that are natural for each object X in C . We may also write $[Y, Z]$ or $Y \multimap Z$ for a right internal hom object for Y and Z .

- A **left internal hom object** for Y and Z is a right internal hom object $Y \multimap Z$ in the reverse monoidal structure on C ; equivalently, $Y \multimap Z$ is an object equipped with bijections

$$C(Y \otimes X, Z) \cong C(X, Y \multimap Z)$$

that are natural for each object X in C . We may also write Z^Y or $Z \multimap Y$ for a left internal hom object for Y and Z .

- A **right-closed monoidal category** is a monoidal category that has right internal hom object for all pairs of objects.
- A **left-closed monoidal category** is a monoidal category that has left internal hom objects for all pairs of objects.
- A **biclosed monoidal category** is a monoidal category that is both left-closed and right-closed.

Note that in a symmetric monoidal category, $Y \pitchfork Z$ and $\mathcal{H}om(Y, Z)$ are naturally isomorphic if they exist; a **closed symmetric monoidal category** is a symmetric monoidal category that is biclosed.

Proposition B.I.15. *Let \mathcal{C} be a right-closed monoidal category.*

- (i) *The assignment $(Y, Z) \mapsto \mathcal{H}om(Y, Z)$ extends to a functor $\mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$ making the bijection*

$$\mathcal{C}(X \otimes Y, Z) \cong \mathcal{C}(X, \mathcal{H}om(Y, Z))$$

natural in $X, Y,$ and $Z.$

- (ii) *For each object $Y,$ we have an adjunction*

$$(-) \otimes Y \dashv \mathcal{H}om(Y, -) : \mathcal{C} \rightarrow \mathcal{C}$$

whose counit is $ev_{Y,-} : \mathcal{H}om(Y, -) \otimes Y \Rightarrow id_{\mathcal{C}}.$

- (iii) *If I is the monoidal unit of $\mathcal{C},$ then there is a bijection*

$$\mathcal{C}(Y, Z) \cong \mathcal{C}(I, \mathcal{H}om(Y, Z))$$

that is natural in Y and $Z.$

Proof. (i). This is a straightforward example of an adjunction with a parameter.^[4]

(ii). This is clear from the definition of $\mathcal{H}om(Y, Z)$ and $ev_{Y,-}.$

(iii). The left unitor $\lambda_Y : Y \xrightarrow{\cong} I \otimes Y$ induces the required bijection. ■

REMARK B.I.16. A cartesian monoidal category is a closed symmetric monoidal category if and only if it is a cartesian closed category ([definition A.2.3](#)).

[4] See [CWM, Ch. IV, § 7].

B.2 Categories with actions

Prerequisites. § B.1.

Definition B.2.1. Let \mathcal{V} be a monoidal category. A **left \mathcal{V} -action** on a category \mathcal{C} is a strong monoidal functor $\mathcal{V} \rightarrow [\mathcal{C}, \mathcal{C}]$, where $[\mathcal{C}, \mathcal{C}]$ is regarded as a strict monoidal category under composition. Similarly, a **right \mathcal{V} -action** on \mathcal{C} is a strong monoidal functor $\mathcal{V} \rightarrow [\mathcal{C}, \mathcal{C}]^{\text{rev}}$.

REMARK B.2.2. We can unfold the above definition somewhat by taking the left exponential transpose of the strong monoidal functor $\mathcal{V} \rightarrow [\mathcal{C}, \mathcal{C}]$: let \otimes be the corresponding functor $\mathcal{V} \times \mathcal{C} \rightarrow \mathcal{C}$. Since the original functor was strong monoidal, we get a natural isomorphism $\eta : \text{id}_{\mathcal{C}} \Rightarrow I \otimes (-)$ and a natural isomorphism $\mu_{X,Y} : X \otimes (Y \otimes (-)) \Rightarrow (X \otimes Y) \otimes (-)$ for each pair of objects X and Y in \mathcal{V} ; these moreover satisfy the following coherence laws:

$$\begin{array}{ccc}
 X \otimes (-) & \xrightarrow{\eta \circ (X \otimes (-))} & I \otimes (X \otimes (-)) \\
 \text{id} \downarrow & & \downarrow \mu_{I,X} \\
 X \otimes (-) & \xleftarrow{\lambda_X \otimes \text{id}} & (I \otimes X) \otimes (-)
 \end{array}
 \qquad
 \begin{array}{ccc}
 X \otimes (-) & \xrightarrow{(X \otimes (-)) \circ \eta} & X \otimes (I \otimes (-)) \\
 \text{id} \downarrow & & \downarrow \mu_{X,I} \\
 X \otimes (-) & \xleftarrow{\rho_X \otimes \text{id}} & (X \otimes I) \otimes (-)
 \end{array}$$

$$\begin{array}{ccc}
 & W \otimes (X \otimes (Y \otimes (-))) & \\
 \mu_{W,X} \circ (Y \otimes (-)) \swarrow & & \searrow \text{id}_W \circ \mu_{X,Y} \\
 (W \otimes X) \otimes (Y \otimes (-)) & & W \otimes ((X \otimes Y) \otimes (-)) \\
 \mu_{W \otimes X, Y} \downarrow & & \downarrow \mu_{W, X \otimes Y} \\
 ((W \otimes X) \otimes Y) \otimes (-) & \xrightarrow{\alpha_{W,X,Y} \otimes \text{id}} & (W \otimes (X \otimes Y)) \otimes (-)
 \end{array}$$

Conversely, any functor $\otimes : \mathcal{V} \times \mathcal{C} \rightarrow \mathcal{C}$ equipped with such a collection of natural isomorphisms defines a left \mathcal{V} -action on \mathcal{C} .

Proposition B.2.3 (Bénabou). *For any monoidal category \mathcal{C} , there is a faithful strong monoidal functor $F : \mathcal{C} \rightarrow [\mathcal{C}, \mathcal{C}]$ defined by the following data:*

$$\begin{aligned}
 FX &= X \otimes (-) \\
 \eta &= \lambda^{-1} \\
 (\mu_{X,Y})_Z &= \alpha_{X,Y,Z}^{-1}
 \end{aligned}$$

In particular, this defines a left \mathcal{C} -action on \mathcal{C} , called the **left regular representation** of \mathcal{C} .

Proof. F is clearly a faithful functor. In this case, the strong monoidal functor axioms become the following diagrams:

$$\begin{array}{ccc}
 X \otimes Y & \xrightarrow{\lambda_{X \otimes Y}^{-1}} & I \otimes (X \otimes Y) \\
 \text{id} \downarrow & & \downarrow \alpha_{I, X, Y}^{-1} \\
 X \otimes Y & \xleftarrow{\lambda_X \otimes \text{id}_Y} & (I \otimes X) \otimes Y
 \end{array}
 \qquad
 \begin{array}{ccc}
 X \otimes Y & \xrightarrow{\text{id}_X \otimes \lambda_Y^{-1}} & X \otimes (I \otimes Y) \\
 \text{id} \downarrow & & \downarrow \alpha_{X, I, Y}^{-1} \\
 X \otimes Y & \xleftarrow{\rho_X \otimes \text{id}_Y} & (X \otimes I) \otimes Y
 \end{array}$$

$$\begin{array}{ccc}
 & W \otimes (X \otimes (Y \otimes Z)) & \\
 \alpha_{W, X, Y \otimes Z}^{-1} \swarrow & & \searrow \text{id}_W \otimes \alpha_{X, Y, Z}^{-1} \\
 (W \otimes X) \otimes (Y \otimes Z) & & W \otimes ((X \otimes Y) \otimes Z) \\
 \alpha_{W \otimes X, Y, Z}^{-1} \downarrow & & \downarrow \alpha_{W, X \otimes Y, Z}^{-1} \\
 ((W \otimes X) \otimes Y) \otimes Z & \xrightarrow{\alpha_{W, X, Y} \otimes \text{id}_Z} & (W \otimes (X \otimes Y)) \otimes Z
 \end{array}$$

The left square commutes by the coherence theorem, while the right square and the pentagon are seen to be immediate consequences of the triangle and pentagon axioms, respectively. \blacksquare

Proposition B.2.4. *Let \mathcal{V} be a monoidal category and let \mathcal{C} be a category.*

- If $\otimes : \mathcal{V} \times \mathcal{C} \rightarrow \mathcal{C}$ defines a left \mathcal{V} -action on \mathcal{C} such that, for each object X in \mathcal{V} , the endofunctor $X \otimes (-)$ has a right adjoint $(-) \circlearrowleft X$, then the functor $\circlearrowleft : \mathcal{C} \times \mathcal{V}^{\text{op}} \rightarrow \mathcal{C}$ defines a right \mathcal{V}^{op} -action on \mathcal{C} .
- If $\otimes : \mathcal{C} \times \mathcal{V} \rightarrow \mathcal{C}$ defines a right \mathcal{V} -action on \mathcal{C} such that, for each object X in \mathcal{V} , the endofunctor $(-) \otimes X$ has a right adjoint $X \circlearrowright (-)$, then the functor $\circlearrowright : \mathcal{V}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$ defines a left \mathcal{V}^{op} -action on \mathcal{C} .
- If $\circlearrowleft : \mathcal{C} \times \mathcal{V}^{\text{op}} \rightarrow \mathcal{C}$ defines a right \mathcal{V}^{op} -action on \mathcal{C} such that, for each object X in \mathcal{V} , the endofunctor $X \circlearrowleft (-)$ has a left adjoint $X \otimes (-)$, then the functor $\otimes : \mathcal{V} \times \mathcal{C} \rightarrow \mathcal{C}$ defines a left \mathcal{V} -action on \mathcal{C} .
- If $\circlearrowright : \mathcal{V}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$ defines a left \mathcal{V}^{op} -action on \mathcal{C} such that, for each object X in \mathcal{V} , the endofunctor $X \circlearrowright (-)$ has a left adjoint $(-) \otimes X$, then the functor $\otimes : \mathcal{C} \times \mathcal{V} \rightarrow \mathcal{C}$ defines a right \mathcal{V} -action on \mathcal{C} .

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Proof. The four statements are related by applying $(-)^{\text{op}}$ and $(-)^{\text{rev}}$ at the appropriate points, so it suffices to prove the first claim.

First, note that $\circ-$ is indeed a functor $\mathcal{C} \times \mathcal{V}^{\text{op}} \rightarrow \mathcal{C}$, by the parameter theorem for adjunctions.^[5] Let $\text{ev}_{X,A} : X \otimes (A \circ X) \rightarrow A$ denote the component of the counit of the adjunction $X \otimes (-) \dashv (-) \circ X$ at an object A in \mathcal{C} . For each pair of objects X and Y in \mathcal{V} and each object A in \mathcal{C} , we define the morphism $(\delta_{X,Y})_A : A \circ (X \otimes Y) \rightarrow (A \circ X) \circ Y$ to be the right adjoint transpose of $\text{ev}_{X \otimes Y, A} \circ (\mu_{X,Y})_{(A \circ X) \circ Y}$, and for each A , we define $\varepsilon_A : A \circ I \rightarrow A$ to be the composite $\text{ev}_{I,A} \circ \eta_{A \circ I}$. These are clearly natural in A , and it is straightforward to check that $\delta_{X,Y}$ is also natural in X and Y . One may then use the calculus of mates to show that ε and $\delta_{X,Y}$ are natural isomorphisms and that they satisfy the axioms for making the right exponential transpose of $\circ- : \mathcal{C} \times \mathcal{V}^{\text{op}} \rightarrow \mathcal{C}$ into a strong monoidal functor $\mathcal{V}^{\text{op}} \rightarrow [\mathcal{C}, \mathcal{C}]^{\text{rev}}$, i.e. a right \mathcal{V}^{op} -action on \mathcal{C} . ■

Example B.2.5. \mathcal{V} is a left-closed (resp. right-closed) monoidal category if and only if the left (resp. right) self-action of \mathcal{V} has a parametrised right adjoint as in the proposition, and the right adjoint right (resp. left) \mathcal{V}^{op} -action so obtained is precisely a left (resp. right) internal hom functor.

Definition B.2.6. Let \mathcal{V} be a monoidal category and let \mathcal{C} be a category.

- A **right \mathcal{V} -hom system** for \mathcal{C} consists of a left \mathcal{V} -action $\otimes : \mathcal{V} \times \mathcal{C} \rightarrow \mathcal{C}$, a functor $\underline{C} : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{V}$, and a right \mathcal{V}^{op} -action $\circ- : \mathcal{C} \times \mathcal{V}^{\text{op}} \rightarrow \mathcal{C}$ together with natural bijections of the types below,

$$\begin{aligned} \mathcal{V}(X, \underline{C}(A, B)) &\cong \mathcal{C}(A, B \circ X) \\ \mathcal{C}(X \otimes A, B) &\cong \mathcal{C}(A, B \circ X) \\ \mathcal{C}(X \otimes A, B) &\cong \mathcal{V}(X, \underline{C}(A, B)) \end{aligned}$$

where X varies over the objects in \mathcal{V} , and A and B vary over the objects in \mathcal{C} , such that the cyclic composition of the three bijections is the identity.

- A **left \mathcal{V} -hom system** for \mathcal{C} consists of a right \mathcal{V} -action $\otimes : \mathcal{C} \times \mathcal{V} \rightarrow \mathcal{C}$, a functor $\underline{C} : \mathcal{C}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{V}$, and a left \mathcal{V}^{op} -action $\circ- : \mathcal{V}^{\text{op}} \times \mathcal{C} \rightarrow \mathcal{C}$, together

[5] See [CWM, Ch. IV, § 7].

with natural bijections of the types below,

$$\begin{aligned}\mathcal{V}(X, \underline{C}(A, B)) &\cong C(A, X \multimap B) \\ C(A \otimes X, B) &\cong C(A, X \multimap B) \\ C(A \otimes X, B) &\cong \mathcal{V}(X, \underline{C}(A, B))\end{aligned}$$

where X varies over the objects in \mathcal{V} , and A and B vary over the objects in C , such that the cyclic composition of the three bijections is the identity.

Example B.2.7. If \mathcal{V} is a biclosed monoidal category with right internal hom functor $\mathcal{H}om$ and left internal hom functor \multimap , then $(\otimes, \multimap, \mathcal{H}om)$ is a left \mathcal{V} -hom system for \mathcal{V} :

$$\begin{aligned}\mathcal{V}(Y, X \multimap Z) &\cong \mathcal{V}(X, \mathcal{H}om(Y, Z)) \\ \mathcal{V}(X \otimes Y, Z) &\cong \mathcal{V}(X, \mathcal{H}om(Y, Z)) \\ \mathcal{V}(X \otimes Y, Z) &\cong \mathcal{V}(Y, X \multimap Z)\end{aligned}$$

Example B.2.8. If C is a locally small category that has products and coproducts for all small families of objects, then C admits a left **Set**-action and a right **Set**^{op}-action that are related by the following adjunctions:

$$\begin{aligned}\mathbf{Set}(X, C(A, B)) &\cong C(A, B \multimap X) \\ C(X \otimes A, B) &\cong C(A, B \multimap X) \\ C(X \otimes A, B) &\cong \mathbf{Set}(X, C(A, B))\end{aligned}$$

(The adjointness claim was checked in [proposition A.6.9](#), and the coherence laws are straightforwardly verified.) Thus, (\otimes, C, \multimap) is a right **Set**-hom system for C .

Theorem B.2.9. *Let \mathcal{V} be a monoidal category and let C be a category.*

- (i) *If \otimes is a left \mathcal{V} -action on C and $\underline{C} : C^{\text{op}} \times C \rightarrow C$ is a functor with natural bijections of the form below,*

$$C(X \otimes A, B) \cong \mathcal{V}(X, \underline{C}(A, B))$$

then \underline{C} is the hom functor of a \mathcal{V} -enriched category \underline{C} whose underlying ordinary category is isomorphic to C .

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(ii) If $\circ-$ is a right \mathcal{V}^{op} -action on C and $\underline{C} : C^{\text{op}} \times C \rightarrow C$ is a functor with natural bijections of the form below,

$$C(A, B \circ- X) \cong \mathcal{V}(X, \underline{C}(A, B))$$

then \underline{C} is the hom functor of a \mathcal{V} -enriched category \underline{C} whose underlying ordinary category is isomorphic to C .

Proof. (i). The natural isomorphism $A \cong I \otimes A$ induces a family of bijections

$$C(A, B) \cong \mathcal{V}(I, \underline{C}(A, B))$$

natural in A and B , so we have a morphism $e_A : I \rightarrow \underline{C}(A, A)$ in \mathcal{V} for every object A in C corresponding to $\text{id}_A : A \rightarrow A$ in C . Let $\text{ev}_{A,B} : \underline{C}(A, B) \otimes A \rightarrow B$ be the component at B of the counit of the adjunction $(-) \otimes A \dashv \underline{C}(A, -)$, and define $c_{A,B,C} : \underline{C}(B, C) \otimes \underline{C}(A, B) \rightarrow \underline{C}(A, C)$ to be the right adjoint transpose of the following morphism in C :

$$\text{ev}_{B,C} \circ (\text{id}_{\underline{C}(B,C)} \otimes \text{ev}_{A,B}) \circ (\mu_{\underline{C}(B,C), \underline{C}(A,B)})_A^{-1} : (\underline{C}(B, C) \otimes \underline{C}(A, B)) \otimes A \rightarrow C$$

By definition, the left adjoint transpose of e_B is η_B^{-1} , so the left and right unit axioms are satisfied:

$$\begin{aligned} c_{A,B,B} \circ (e_B \otimes \text{id}_{\underline{C}(A,B)}) &= \lambda_{\underline{C}(A,B)} \\ c_{B,B,C} \circ (\text{id}_{\underline{C}(B,C)} \otimes e_B) &= \rho_{\underline{C}(B,C)} \end{aligned}$$

One may similarly verify the associativity axiom:

$$c_{A,B,D} \circ (c_{B,C,D} \otimes \text{id}_{\underline{C}(A,B)}) = c_{A,C,D} \circ (\text{id}_{\underline{C}(C,D)} \otimes c_{A,B,C}) \circ \alpha_{\underline{C}(C,D), \underline{C}(B,C), \underline{C}(A,B)}$$

(ii). By duality and symmetry, $\circ-$ induces a left \mathcal{V}^{rev} -action on C^{op} , so we may construct a \mathcal{V}^{rev} -enriched category $\underline{C}^{\text{op}}$ using claim (i) and thence a \mathcal{V} -enriched category $\underline{C} = (\underline{C}^{\text{op}})^{\text{op}}$. \blacksquare

Definition B.2.10. Let \mathcal{V} be a monoidal category, and let C and D be categories with left \mathcal{V} -actions. A \mathcal{V} -**strength** for a functor $F : C \rightarrow D$ is a natural transformation $\sigma : (-) \otimes F(-) \Rightarrow F(- \otimes -)$ making these diagrams commute:

$$\begin{array}{ccc} & FA & \\ \eta_{FA} \swarrow & & \searrow F\eta_A \\ I \otimes FA & \xrightarrow{\sigma_{I,A}} & F(I \otimes A) \end{array}$$

$$\begin{array}{ccc}
 & X \otimes (Y \otimes FA) & \\
 \text{id}_X \otimes \sigma_{Y,A} \swarrow & & \searrow (\mu_{X,Y})_A \\
 X \otimes F(Y \otimes A) & & (X \otimes Y) \otimes FA \\
 \sigma_{X,Y \otimes A} \downarrow & & \downarrow \sigma_{X \otimes Y, A} \\
 F(X \otimes (Y \otimes A)) & \xrightarrow{F(\mu_{X,Y})_A} & F((X \otimes Y) \otimes A)
 \end{array}$$

A \mathcal{V} -strong functor is a functor equipped with a \mathcal{V} -strength.

Definition B.2.II. Let \mathcal{V} be a monoidal category, let C and D be categories with left \mathcal{V} -actions, and let $F, F' : C \rightarrow D$ be functors with \mathcal{V} -strengths σ and σ' respectively. A \mathcal{V} -strong natural transformation $\varphi : F \Rightarrow F'$ is a natural transformation making the following diagram commute:

$$\begin{array}{ccc}
 X \otimes FA & \xrightarrow{\sigma_{X,A}} & F(X \otimes A) \\
 \text{id}_X \otimes \varphi_A \downarrow & & \downarrow \varphi_{X \otimes A} \\
 X \otimes F'A & \xrightarrow{\sigma'_{X,A}} & F'(X \otimes A)
 \end{array}$$

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